


# Smith Waterman Algorithm for Sequence Alignment

Karthika Jayaprakash

CSE 633 Parallel Algorithms (Dr. Russ Miller)  
April 21, 2020

 University at Buffalo  
Department of Computer Science  
and Engineering  
School of Engineering and Applied Sciences



# Smith Waterman Overview

Biological sequence alignment is a frequently performed task in bioinformatics.

The Smith-Waterman algorithm, based on dynamic programming, is one of the most fundamental algorithms used in local sequence alignment.

# Smith Waterman Overview

Given two sequences find the best local alignment.

	T	G	T	T	A	C	G	G
	0	0	0	0	0	0	0	0
G	0	0	3	1	0	0	0	3
G	0	0	3	1	0	0	0	3
T	0	3	1	6	4	2	0	1
T	0	3	1	4	9	7	5	3
G	0	1	6	4	7	6	4	8
A	0	0	4	3	5	10	8	6
C	0	0	2	1	3	8	13	11
T	0	3	1	5	4	6	11	10
A	0	1	0	3	2	7	9	8



	T	G	T	T	A	C	G	G
	0	0	0	0	0	0	0	0
G	0	0	3	1	0	0	0	3
G	0	0	3	1	0	0	0	3
T	0	3	1	6	4	2	0	1
T	0	3	1	4	9	7	5	3
G	0	1	6	4	7	6	4	8
A	0	0	4	3	5	10	8	6
C	0	0	2	1	3	8	13	11
T	0	3	1	5	4	6	11	10
A	0	1	0	3	2	7	9	8

3	6	9	7	10	13
G	T	T	-	A	C
G	T	T	G	A	C

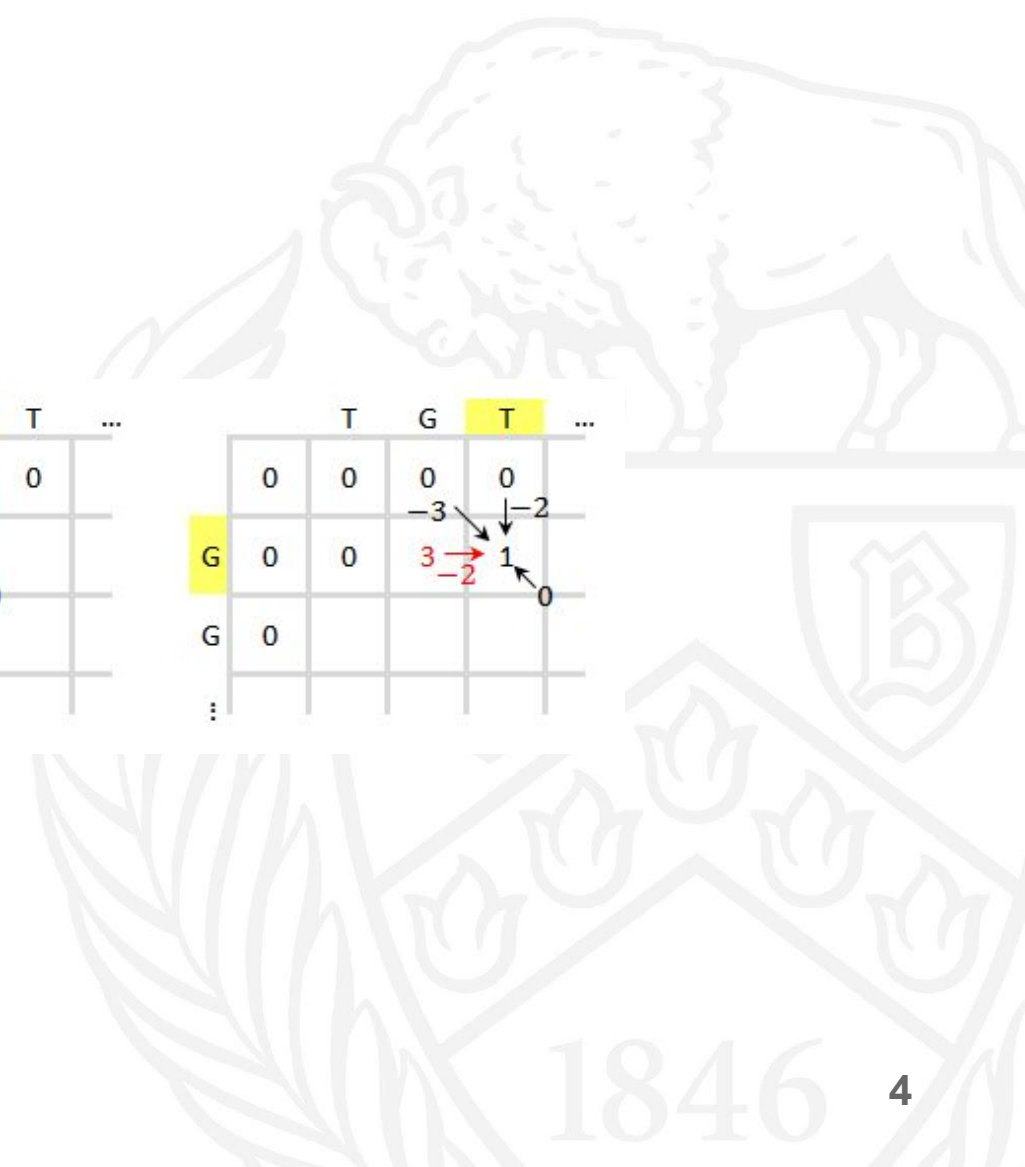
# Smith Waterman Overview

	T	G	T	...
	0	0	0	0
G	0			
G	0			
⋮				

	T	G	T	...
	0	0	0	0
G	-3	0	-2	
G	0	-2	0	0
G	0			
⋮				

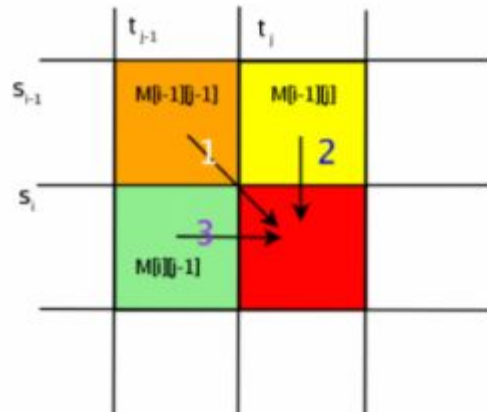
	T	G	T	...
	0	0	0	0
G	0	0	-2	
G	0	0	3	0
G	0			
⋮				

	T	G	T	...
	0	0	0	0
G	0	0	-3	-2
G	0	0	3	1
G	0			
⋮				



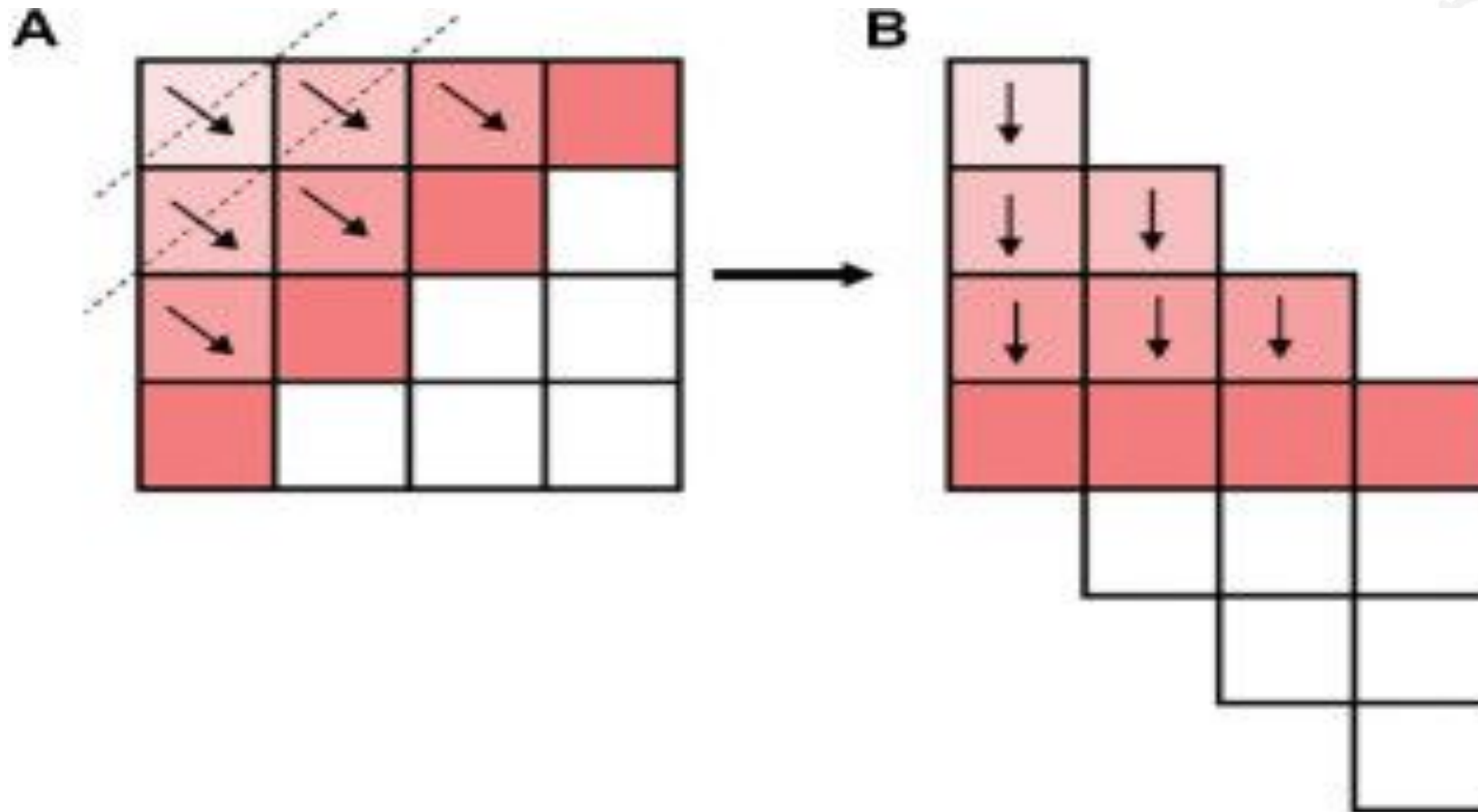
# Smith Waterman - Parallel Approach

- Diagonal Calculations are independent of each other.
- Hence Diagonal Computations are parallelized.

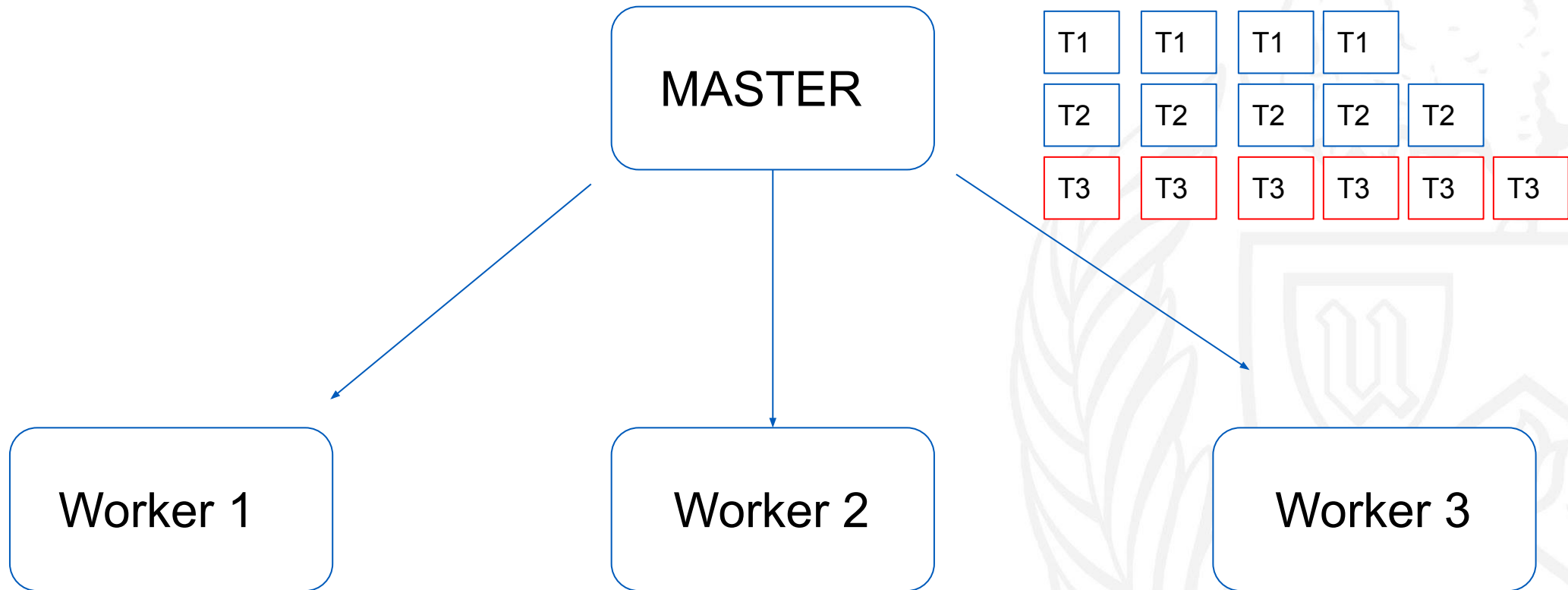


	-	C	G	G	G	T	A	T	C
-	0	0	0	0	0	0	0	0	0
C	0	T1	T2	T3	T4	T5	T6	T7	T8
C	0	T2	T3	T4	T5	T6	T7	T8	T9
C	0	T3	T4	T5	T6	T7	T8	T9	T10
T	0	T4	T5	T6	T7	T8	T9	T10	T11
A	0	T5	T6	T7	T8	T9	T10	T11	T12
G	0	T6	T7	T8	T9	T10	T11	T12	T13
G	0	T7	T8	T9	T10	T11	T12	T13	T14
T	0	T8	T9	T10	T11	T12	T13	T14	T15

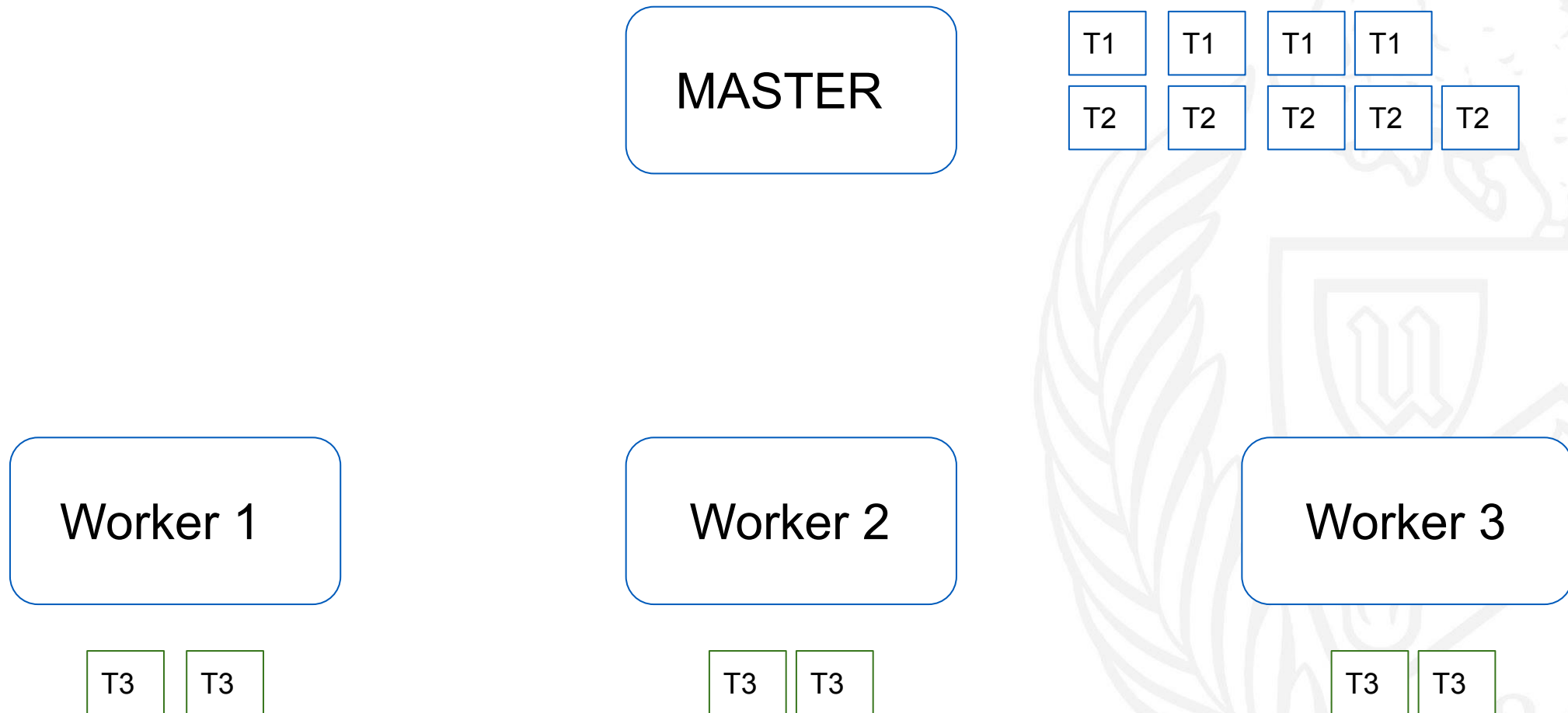
# Smith Waterman - Parallel Approach



# Smith Waterman - Parallel Approach

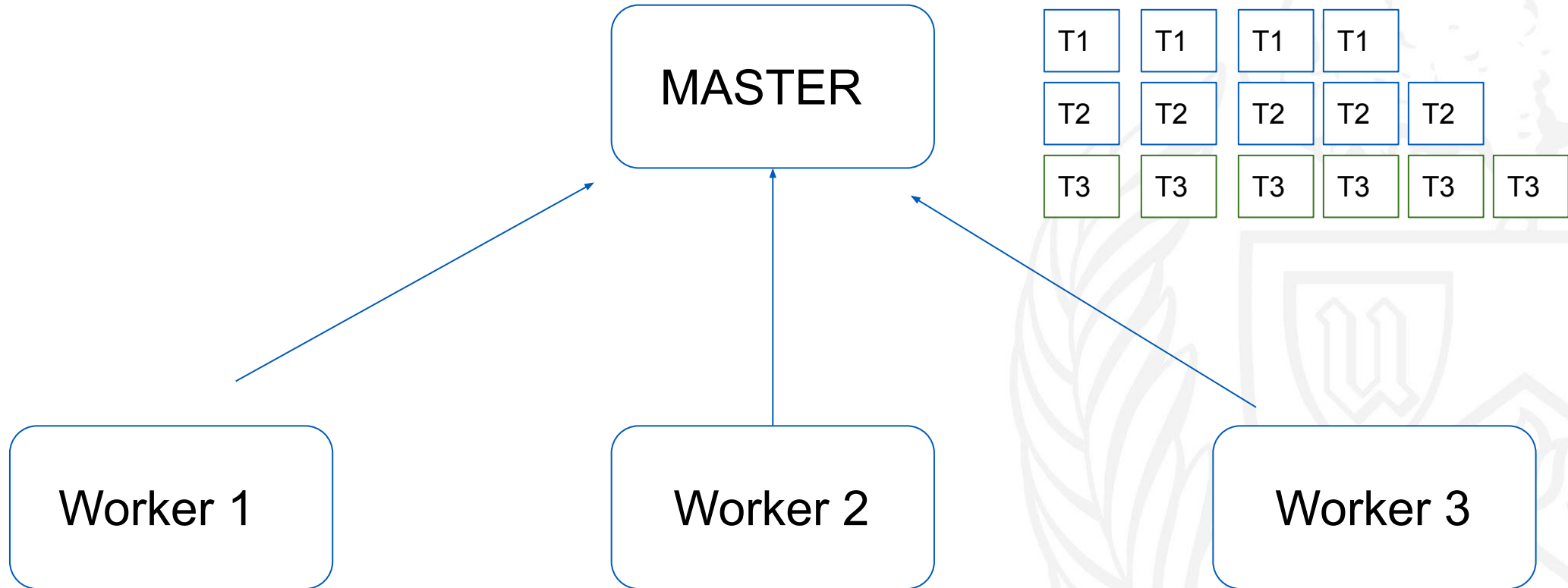


# Smith Waterman - Parallel Approach





# Smith Waterman - Parallel Approach





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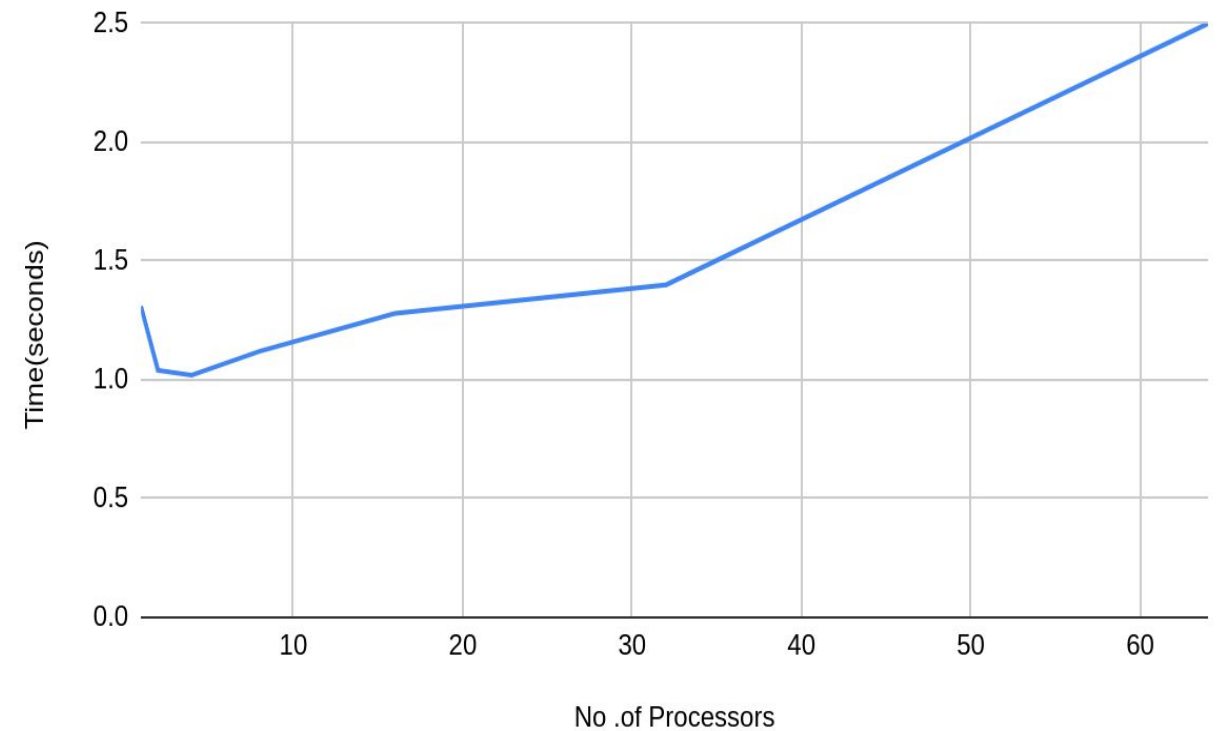
# Runtime Comparison



# 1000 x 1000 matrix

1	1.31
2	1.04
4	1.02
8	1.12
16	1.28
32	1.4
64	2.5

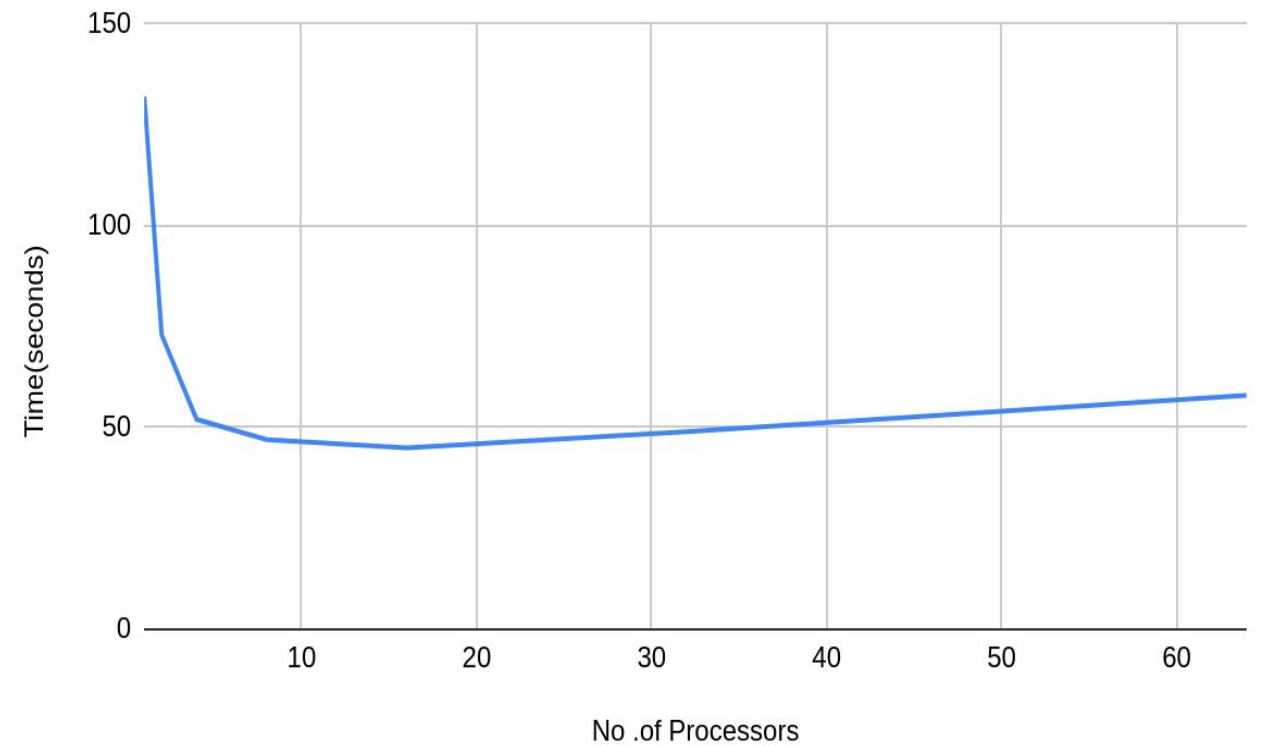
Time(seconds) vs. No .of Processors



# 10000 x 10000 matrix

1	132
2	73
4	52
8	47
16	45
32	49
64	58

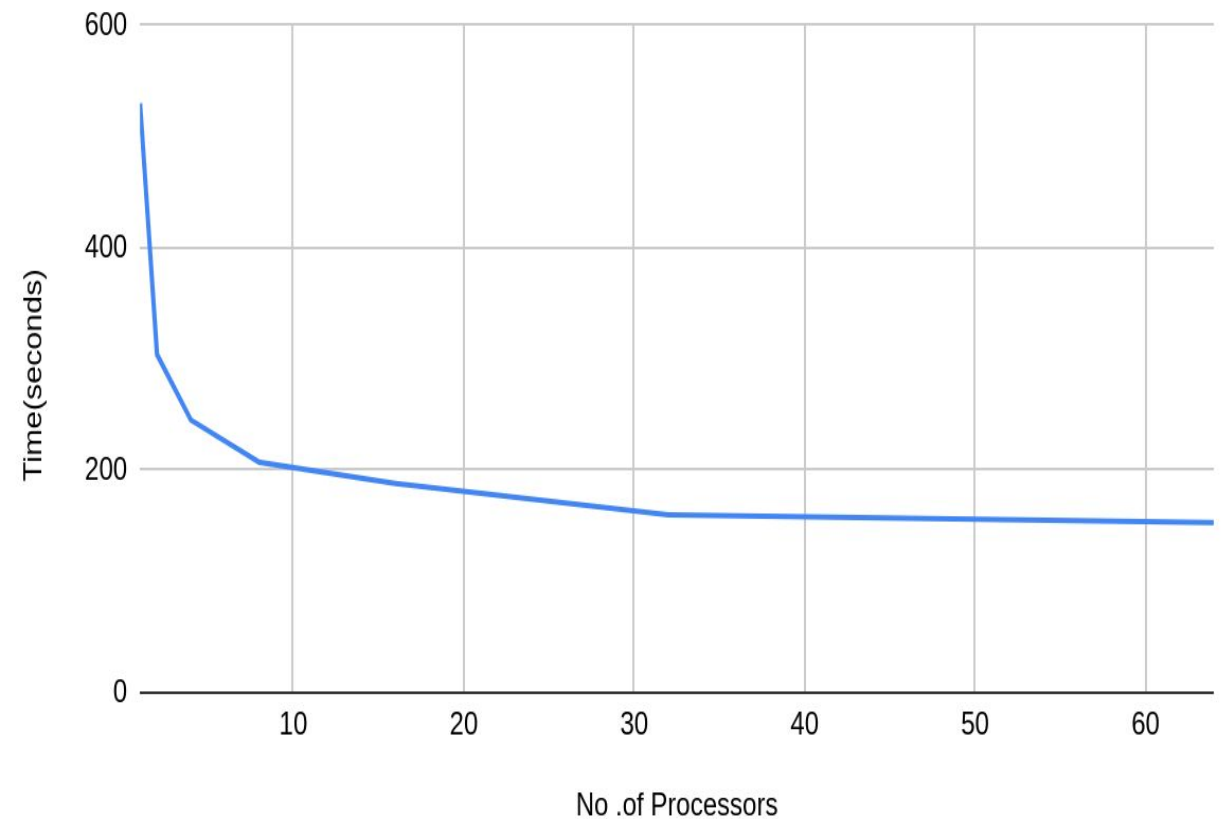
Time(seconds) vs. No .of Processors



# 20000 x 20000 matrix

1	530
2	304
4	245
8	207
16	188
32	160
64	153

Time(seconds) vs. No .of Processors



# Conclusion

The computations which can be executed parallelly are identified.

The diagonals are computed parallelly at each step.

The data in each diagonal are split among the processes to compute scores.

Parallelization has resulted in speedup of computations.

# References

Fa ZHANG<sup>1</sup> , Xiang-Zhen QIAO<sup>1</sup> , Zhi-Yong LIU “A Parallel Smith-Waterman Algorithm Based on Divide and Conquer” Proceedings of the Fifth International Conference on Algorithms and Architectures for Parallel Processing

MPI4PY documentation - <https://mpi4py.readthedocs.io/en/stable/tutorial.html>

CCR Resources - <https://ubccr.freshdesk.com/support/home>



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# Thank You

