

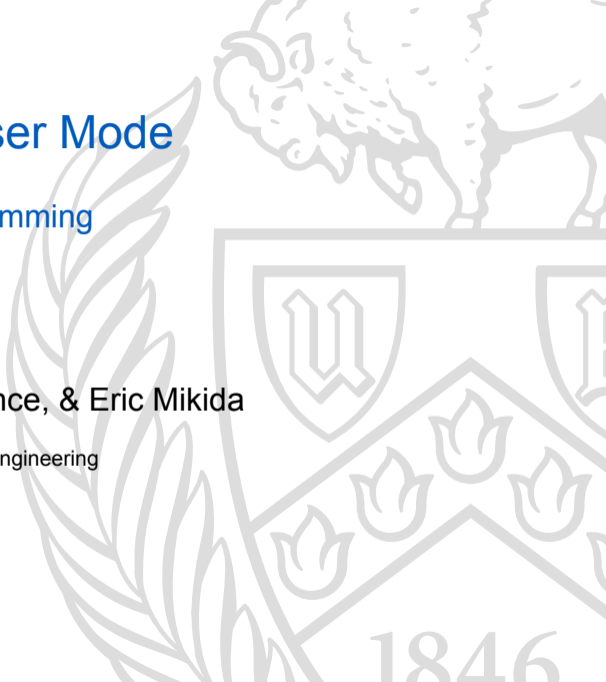
# The Kernel and User Mode

CSE 220: Systems Programming

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# The Operating System Kernel

We have talked about **the operating system** or **the kernel**.

The operating system **manages the hardware**.

On our systems, it also:

- Supports the **dedicated computer** model
- Provides **protection** against misbehaving programs

The kernel is **the code of the inner core of the OS**.

In some sense the OS and Kernel are **just programs**.

# User Mode Programs

Our programs run in **user mode**.

User mode programs **appear** to run on a **dedicated computer**.

This means that **shared resources** must be managed for them.

User mode programs **ask the kernel** for access to shared resources.

If the user mode program has a dedicated computer ... **where is the kernel?**

# Exceptions

Exceptions are another type of **control flow**.

Unlike *if*, *for*, *etc.*, they:

- Allow **non-local** (to another function or even program) transfer of control
- Can be **asynchronous** (triggered by an external event)

Exceptions may be caused by **hardware** or **software**.

The handling of exceptions requires both.

# System Calls

A **system call** is a special kind of exception.

It allows a program to:

- “break out” of its **dedicated computer**, and
- **contact the kernel**

System calls are **synchronous** but **non-local** transfer of control.

# Lecture Question

Ask a review question!

# The Kernel and Supervisor Mode

The kernel **does not have** a dedicated computer.<sup>¶</sup>

The kernel has the **real computer!**

It runs in a special mode (often called **supervisor mode**).

It can:

- Access hardware **directly**
- Manipulate virtual memory mappings
- Modify process memory
- ...

# Protection Domains

**Protection domains** represent the amount of privilege to access the “real computer” allowed to a process.

Supervisor mode is a special **protection domain**.

User mode is a **less-privileged** protection domain.

Protection domains are a **hardware capability**.

User programs run in **user mode**, the kernel in **supervisor mode**.

The hardware **enforces access restrictions** on user mode.

Some hardware has **more than two** protection domains.



# Changing Protection Domains

Changing protection domains is a **supervisor mode operation**.

This prevents programs from **breaking out** of user mode.

It also means there must be a **safe way** to switch domains!

We will see how **exceptions** provide a controlled mode change.

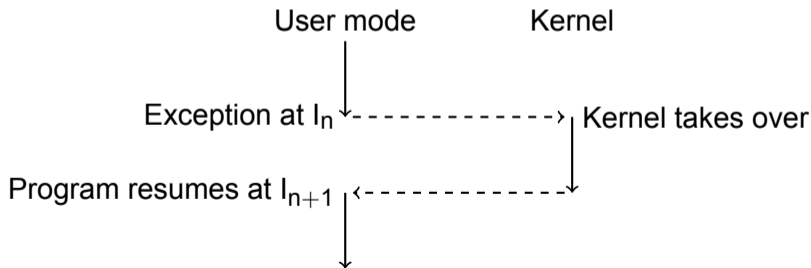
Changing protection domains can be **slow** and **expensive**.

# Exception Flow

When an exception occurs, **control passes to the kernel**.

If control is **already in the kernel**, it **changes location**.

If control is in a user mode program, it **switches contexts**.



# Types of Exceptions

There are four major types of exceptions:

- **Interrupts** are **asynchronous** notifications from hardware
- **Traps** are **synchronous** exceptions caused by software intentionally
- **Faults** are **synchronous** exceptions caused by software due to potentially recoverable errors
- **Aborts** are **synchronous** exceptions caused by unrecoverable errors outside of software control

We have only seen **faults** thus far (remember page faults?).

We are currently most interested in **traps**.

# Interrupts

Interrupts are a way for **hardware to signal the OS**.

Examples:

- A network packet has arrived
- A clock has “ticked”
- A disk has completed a read

Interrupts are handled **by the kernel**.

We will not discuss them more, although they are in the text.

# Faults

We have already seen faults!

Segmentation fault (core dumped)

Faults are **program errors** that **may be recoverable**.

When a fault occurs, the kernel may:

- try to fix it
- notify the program

It may also **terminate the program** or **shut down**.

# Fault Recovery

Some faults are not true errors:  
e.g., **page faults** to bring in new pages.

Other faults may be **recoverable by the program**:

- Divide by zero
- Segmentation fault
- Bus error
- ...

Each of these **is an error**, but might not be **fatal**.

For example, a calculation might usefully return **some concrete value** if it reaches a divide-by-zero.

# Aborts

Aborts are **relatively uninteresting** to us.

They represent some **unrecoverable error** that often ends in:

- Rebooting the computer
- Shutting down the computer
- Terminating some or all processes
- *etc.*

Aborts are handled **by the kernel**.

We will not discuss them more.

# Traps

Traps are **software-generated exceptions**.  
(They are sometimes called software interrupts.)

They are generated by **special instructions** run by a program.

Their critical feature is:

**Trap handlers** are run **by the kernel in supervisor mode**.

This means that a **user mode program** can **call into the kernel**.

This provides a **safe method** of changing protection domains.



# Lecture Question

Ask a lecture question!

# System Calls

System calls are:

- traps
- used by user-mode programs
- to invoke kernel functions

Many platforms have a dedicated hardware instruction for this:

- ARM: `svc`
- x86-64: `syscall`

# System Call Handling

When the system call instruction runs, the hardware:

- Switches to **supervisor mode**
- Invokes a specific kernel routine

When the kernel receives a system call, it:

- Identifies **what the program wants**
- Verifies the program **arguments**
- Authenticates the request
- Performs the operation (or indicates failure)

This allows the **kernel to decide** whether a program can **access something outside its “dedicated computer”**.

# The Implications of the Trap

User mode **cannot control** what code the kernel runs.

This is:

- **controlled** by the hardware
- **configured** by the kernel

This is how modern operating systems **protect themselves** from malicious or buggy programs.

# Invoking a System call

We have **invoked system calls!**

`open()`, `sbrk()`, `mmap()`, *etc.* are **system calls!**

**Anything** outside the dedicated computer needs a system call.

We never used the `syscall` instruction.

The **C library** makes system calls **look like a C function.**

*All functions in manual section 2 are system calls.*

# Overhead

System calls **are very slow**.

They can take tens to **hundreds of thousands** of clock cycles.

This is due to:

- Changing protection domains
- Validating arguments
- Adjusting memory mappings
- Cache effects
- ...

Programs should **make fewer system calls** when practical.

# Summary

- Exceptions are special control flow
- Protection domains control access to hardware resources
- Exception handlers run in supervisor mode in the kernel
- Special trap exceptions can be used to implement system calls
- System calls allow user mode programs to request access to the kernel

# References I

## Required Readings

- [2] Ian Weinand. *Computer Science from the Bottom Up*. Chapter 4: part 1; part 2 except 2.1 and its subsections; part 3 through 3.1. URL: <https://www.bottomupcs.com/index.html>.

## Optional Readings

- [1] Randal E. Bryant and David R. O'Hallaron. *Computer Science: A Programmer's Perspective*. Third Edition. Chapter 8: Intro, 8.1, 8.2. Pearson, 2016.
- [3] Ian Weinand. *Computer Science from the Bottom Up*. Chapter 4. URL: <https://www.bottomupcs.com/index.html>.



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