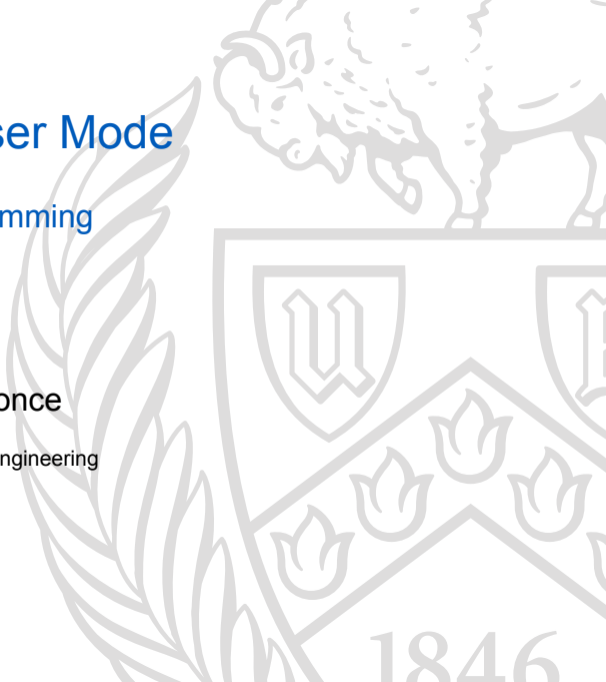


The Kernel and User Mode

CSE 220: Systems Programming

Ethan Blanton & Carl Alphonse

Department of Computer Science and Engineering
University at Buffalo



The Operating System Kernel

We have talked about **the operating system** or **the kernel**.

The operating system **manages the hardware**.

On our systems, it also:

- Supports the **dedicated computer** model
- Provides **protection** against misbehaving programs

The kernel is **the code of the inner core of the OS**.

In some sense the OS and Kernel are **just programs**.

User Mode Programs

Our programs run in **user mode**.

User mode programs **appear** to run on a **dedicated computer**.

This means that **shared resources** must be managed for them.

User mode programs **ask the kernel** for access to shared resources.

If the user mode program has a dedicated computer ... **where is the kernel?**

Exceptions

Exceptions are another type of **control flow**.

Unlike *if*, *for*, *etc.*, they:

- Allow **non-local** (to another function or even program) transfer of control
- Can be **asynchronous** (triggered by an external event)

Exceptions may be caused by **hardware** or **software**.

The handling of exceptions requires both.

System Calls

A **system call** is a special kind of exception.

It allows a program to:

- “break out” of its **dedicated computer**, and
- **contact the kernel**

System calls are **synchronous** but **non-local** transfer of control.

Lecture Question

Ask a review question!

The Kernel and Supervisor Mode

The kernel **does not have** a dedicated computer.¶

The kernel has the **real computer!**

It runs in a special mode (often called **supervisor mode**).

It can:

- Access hardware **directly**
- Manipulate virtual memory mappings
- Modify process memory
- ...

Protection Domains

Protection domains represent the amount of privilege to access the “real computer” allowed to a process.

Supervisor mode is a special **protection domain**.

User mode is a **less-privileged** protection domain.

Protection domains are a **hardware capability**.

User programs run in **user mode**, the kernel in **supervisor mode**.

The hardware **enforces access restrictions** on user mode.

Some hardware has **more than two** protection domains.

Changing Protection Domains

Changing protection domains is a **supervisor mode operation**.

This prevents programs from **breaking out** of user mode.

It also means there must be a **safe way** to switch domains!

We will see how **exceptions** provide a controlled mode change.

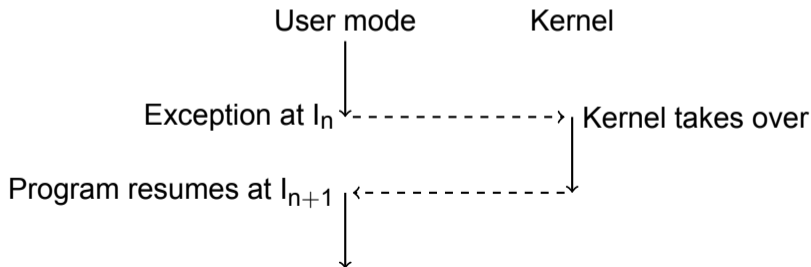
Changing protection domains can be **slow** and **expensive**.

Exception Flow

When an exception occurs, **control passes to the kernel**.

If control is **already in the kernel**, it **changes location**.

If control is in a user mode program, it **switches contexts**.



Types of Exceptions

There are four major types of exceptions:

- **Interrupts** are **asynchronous** notifications from hardware
- **Traps** are **synchronous** exceptions caused by software intentionally
- **Faults** are **synchronous** exceptions caused by software due to potentially recoverable errors
- **Aborts** are **synchronous** exceptions caused by unrecoverable errors outside of software control

We have only seen **faults** thus far (remember page faults?).

We are currently most interested in **traps**.

Interrupts

Interrupts are a way for **hardware to signal the OS**.

Examples:

- A network packet has arrived
- A clock has “ticked”
- A disk has completed a read

Interrupts are handled **by the kernel**.

We will not discuss them more, although they are in the text.

Faults

We have already seen faults!

Segmentation fault (core dumped)

Faults are **program errors** that **may be recoverable**.

When a fault occurs, the kernel may:

- try to fix it
- notify the program

It may also **terminate the program** or **shut down**.

Fault Recovery

Some faults are not true errors:
e.g., **page faults** to bring in new pages.

Other faults may be **recoverable by the program**:

- Divide by zero
- Segmentation fault
- Bus error
- ...

Each of these **is an error**, but might not be **fatal**.

For example, a calculation might usefully return **some concrete value** if it reaches a divide-by-zero.

Aborts

Aborts are **relatively uninteresting** to us.

They represent some **unrecoverable error** that often ends in:

- Rebooting the computer
- Shutting down the computer
- Terminating some or all processes
- *etc.*

Aborts are handled **by the kernel**.

We will not discuss them more.

Traps

Traps are **software-generated exceptions**.
(They are sometimes called software interrupts.)

They are generated by **special instructions** run by a program.

Their critical feature is:

Trap handlers are run **by the kernel in supervisor mode**.

This means that a **user mode program** can **call into the kernel**.

This provides a **safe method** of changing protection domains.

Lecture Question

Ask a lecture question!

System Calls

System calls are:

- traps
- used by user-mode programs
- to invoke kernel functions

Many platforms have a dedicated hardware instruction for this:

- ARM: `svc`
- x86-64: `syscall`

System Call Handling

When the system call instruction runs, the hardware:

- Switches to **supervisor mode**
- Invokes a specific kernel routine

When the kernel receives a system call, it:

- Identifies **what the program wants**
- Verifies the program **arguments**
- Authenticates the request
- Performs the operation (or indicates failure)

This allows the **kernel to decide** whether a program can **access something outside its “dedicated computer”**.

The Implications of the Trap

User mode **cannot control** what code the kernel runs.

This is:

- **controlled** by the hardware
- **configured** by the kernel

This is how modern operating systems **protect themselves** from malicious or buggy programs.

Invoking a System call

We have **invoked system calls!**

`open()`, `sbrk()`, `mmap()`, *etc.* are **system calls!**

Anything outside the dedicated computer needs a system call.

We never used the `syscall` instruction.

The **C library** makes system calls **look like a C function.**

All functions in manual section 2 are system calls.

Overhead

System calls **are very slow**.

They can take tens to **hundreds of thousands** of clock cycles.

This is due to:

- Changing protection domains
- Validating arguments
- Adjusting memory mappings
- Cache effects
- ...

Programs should **make fewer system calls** when practical.

Summary

- Exceptions are special control flow
- Protection domains control access to hardware resources
- Exception handlers run in supervisor mode in the kernel
- Special trap exceptions can be used to implement system calls
- System calls allow user mode programs to request access to the kernel

References I

Required Readings

- [2] Ian Weinand. *Computer Science from the Bottom Up*. Chapter 4: part 1; part 2 except 2.1 and its subsections; part 3 through 3.1. URL: <https://www.bottomupcs.com/index.html>.

Optional Readings

- [1] Randal E. Bryant and David R. O'Hallaron. *Computer Science: A Programmer's Perspective*. Third Edition. Chapter 8: Intro, 8.1, 8.2. Pearson, 2016.
- [3] Ian Weinand. *Computer Science from the Bottom Up*. Chapter 4. URL: <https://www.bottomupcs.com/index.html>.

License

Copyright 2019–2025 Ethan Blanton, All Rights Reserved.

Copyright 2024 Eric Mikida, All Rights Reserved.

Copyright 2022–2025 Carl Alphonse, All Rights Reserved.

Copyright 2019 Karthik Dantu, All Rights Reserved.

Reproduction of this material without written consent of the author is prohibited.

To retrieve a copy of this material, or related materials, see <https://www.cse.buffalo.edu/~eblanton/>.