

CSE 250: ADTs, Sequences, and Arrays

Lecture 8

Sept 15, 2023

Reminders

- PA1 Tests due Sun, Sept 17 at 11:59 PM
 - Recitations will cover writing good test cases.
- PA1 Implementation due Sun, Sept 24 at 11:59 PM
 - Implement a Sorted Linked List
- Autolab etiquette: Test locally *first*.
 - If you (repeatedly) submit code that doesn't compile to autolab, we reserve the right to mock you mercilessly.

Sequences

- 1, 2, 3, 4, 5
- 4, 7, 2, 13, 8, 12, 14, 20, 12, 1
- 17, 14, 20, 14, 17, 12, 8, 9, 8, 6
- 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, ... (Fibonacci Sequence)
- 'H', 'e', 'l', 'l', 'o', ' ', 'W', 'o', 'r', 'l', 'd'
- The lines of a file

What makes a sequence

What is common to these?

- A collection of elements of some type (e.g., numbers, characters, strings).
 - Two elements can have the same value.
- Each element has a unique position.
 - No two elements have the same position
- The positions of all elements are contiguous.
 - For an N element sequence, every position from $[0, N)$ is occupied.

What can we do with a sequence?

What are the simplest things we can do with a sequence?

- Get the 'i'th element.
- Modify the 'i'th element.
- Enumerate the elements in order.
- Get the number of elements

Abstract Data Types

Set the 'i'th element

Get the 'i'th element

Count the elements

Enumerate the elements

Abstract Data Types

Abstract Data Type defines...

- Domain: What kind of data is stored? (e.g., elements, key/value pairs)
- Constraints: How are items related? (e.g., ordered keys)
- Operations: How can the data be accessed/modified (e.g., 'i'th item)

Like a Java interface¹

¹The term `interface` is not quite the same as ADT; The interface only formalizes the permitted operations.

The Sequence ADT

```
1 public interface Sequence<E>
2 {
3     public E get(int idx);
4     public void set(int idx, E value);
5     public int size();
6     public Iterator<E> iterator();
7 }
```

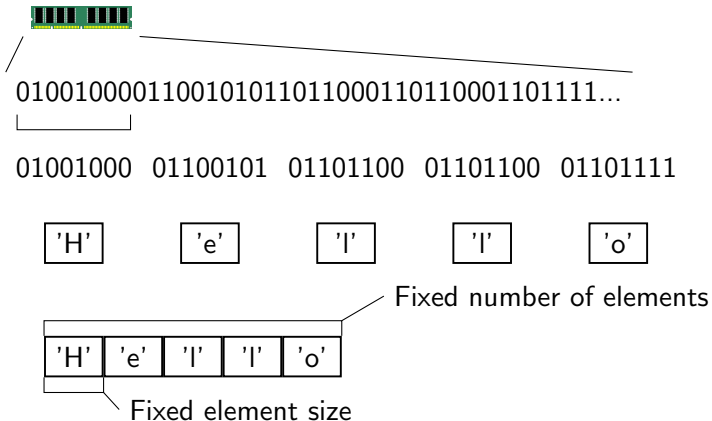
E is the type of thing in the Sequence.

The Sequence ADT



"The Wizard of Oz"; ©1939 Metro-Goldwyn-Mayer

CSE 220 Crossover



RAM

new T finds some unused part of memory big enough to fit a T, marks it used, and returns the *address*².

```
1 int[] data = new int[50];
```

... allocates $4 \cdot 50 = 200$ bytes³.

If data is at address a : where do you look for data[19]?

$$a + 4 \cdot 19 = a + 76$$

How long does it take to locate data[19]?

$$\theta(1)$$

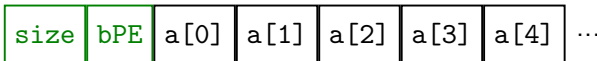
²You'll get *very* familiar with memory allocation in 220.

³An int in java is 4 bytes (32 bits).

Arrays

What information goes into an `T[]` array?

- `size`: 4 bytes for the number of elements⁴.
- `bytesPerElement`: 4 bytes for `sizeof(T)`⁵.
- `data`: `size × bytesPerElement` bytes.



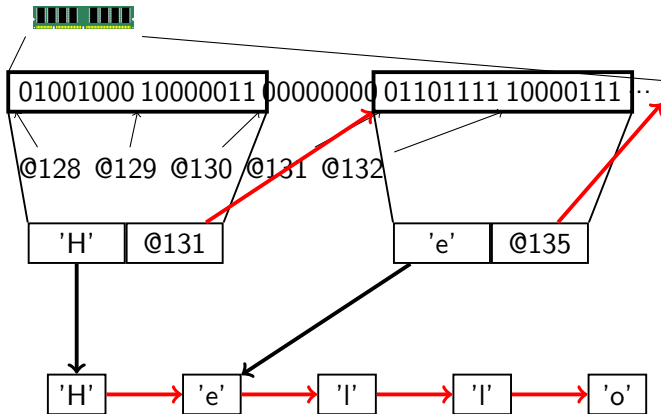
⁴Some languages (e.g., C) skip this, relying on the programmer to track it.

⁵Some languages (e.g., C, C++) skip this, since it's fixed at compile time.

How do we implement...

- `public E get(int idx)`
 - Return bytes $\text{bPE} \times \text{idx}$ to $\text{bPE} \times (\text{idx} + 1) - 1$
 - $\theta(1)$ (if we treat `bPE` as a constant)
- `public void set(int idx, E value)`
 - Update bytes $\text{bPE} \times \text{idx}$ to $\text{bPE} \times (\text{idx} + 1) - 1$
 - $\theta(1)$ (if we treat `bPE` as a constant)
- `public int size()`
 - Return size
 - $\theta(1)$

CSE 220 Crossover 2: List Harder



OpenClipArt: <https://freesvg.org/random-access-computer-memory-ram-vector-image>

Linked Lists

```
1 public class LinkedListNode<T>
2 {
3     T value;
4     LinkedListNode<T> next = null;
5 }
```

```
1 public class LinkedList<T> implements List<T>
2 {
3     LinkedListNode<T> head = null;
4     /* ... */
5 }
```

How do we implement...

- `public E get(int idx)`
 - Start at head, and move to the next element `idx` times.
Return the element's value.
 - $\theta(idx), O(N)$
- `public void set(int idx, E value)`
 - Start at head, and move to the next element `idx` times.
Update the element's value.
 - $\theta(idx), O(N)$
- `public int size()`
 - Start at head, and move to the next element until you reach the end. Return the number of steps taken.
 - $\theta(N)$

Linked Lists' size

Can we do better?

Store size

```
1 public class LinkedList<T> implements List<T>
2 {
3     LinkedListNode<T> head = null;
4     int size = 0;
5     /* ... */
6 }
```

- How expensive is `public int size()` now?
($\theta(1)$)
- How expensive is it to *maintain* size?
(Extra $\theta(1)$ work on insert/remove).

Storing redundant information can reduce complexity.

Optional

```
1 public class LinkedListNode<T>
2 {
3     T value;
4     LinkedListNode<T> next = null;
5 }
```

```
1 public class LinkedList<T> implements List<T>
2 {
3     LinkedListNode<T> head = null;
4     /* ... */
5 }
```

null is risky

```
1 public T get(int idx)
2 {
3     LinkedListNode<T> curr = head;
4     for(int j = 0; j < i; j++)
5     {
6         curr = curr.next;
7     }
8     return curr.value;
9 }
```

NullPointerException

null is risky

```
1 public T get(int idx)
2 {
3     LinkedListNode<T> curr = head;
4     for(int j = 0; j < i; j++)
5     {
6         if(curr.next == null){ /* do something useful */ }
7         curr = curr.next;
8     }
9     return curr.value;
10 }
```

It can be hard to remember when a value might be null.

Optional

```
1 public class LinkedListNode<T>
2 {
3     T value;
4     Optional<LinkedListNode<T>> next = Optional.empty();
5 }
```

```
1 public class LinkedList<T> implements List<T>
2 {
3     Optional<LinkedListNode<T>> head = Optional.empty();
4     /* ... */
5 }
```

Optional

```
1  public T get(int idx)
2  {
3      Optional<LinkedListNode<T>> curr = head;
4      for(int j = 0; j < i; j++)
5      {
6          curr = curr.next; ←
7      }
8      return curr.value; ←
9  }
```

Compiler Error! (No such method)

Optional

```
1 public T get(int idx)
2 {
3     Optional<LinkedListNode<T>> curr = head;
4     for(int j = 0; j < i; j++)
5     {
6         curr = curr.get().next; ←
7     }
8     return curr.get().value; ←
9 }
```

Compiler Error! (Unhandled Exception)

Optional

```
1 public T get(int idx)
2 {
3     Optional<LinkedListNode<T>> curr = head;
4     for(int j = 0; j < i; j++)
5     {
6         try {
7             curr.get().next;
8         } catch(NoSuchElementException e) {
9             /* do something useful */
10        }
11    }
12    try {
13        return curr.get().value;
14    } catch(NoSuchElementException e) {
15        /* do something useful */
16    }
17 }
```

Optional

Creating an Optional

- `Optional.empty()`: Like `null`.
- `Optional.of(x)`: Get an `Optional` with value `x`.
- `Optional.ofNullable(x)`: Like `.of(x)`, but return `.empty()` if `x == null`;

Using an Optional

- `.isPresent()`: Return `true` if a value is present.
- `.get()`: Return the value, or throw exception if not present.
- `.orElse(y)`: Return the value if present, or `y` if not.

Referential Access

What's the complexity of getting the value of the i 'th element of a `LinkedList`?

If you have a pointer to the i 'th `LinkedListNode` of a `LinkedList`, what's the complexity of getting its value?

If you have a pointer to the i 'th `LinkedListNode` of a `LinkedList`, what's the complexity of getting the value of the $i + 1$ 'th element?

If you have a pointer to the i 'th `LinkedListNode` of a `LinkedList`, what's the complexity of getting the value of the $i - 1$ 'th element?

Doubly Linked Lists

```
1 public class LinkedListNode<T>
2 {
3     T value;
4     Optional<LinkedListNode<T>> next = Optional.empty();
5     Optional<LinkedListNode<T>> prev = Optional.empty();
6 }
```

```
1 public class LinkedList<T> implements List<T>
2 {
3     Optional<LinkedListNode<T>> head = Optional.empty();
4     Optional<LinkedListNode<T>> tail = Optional.empty();
5     /* ... */
6 }
```

Linked Lists

Next time...

- Referential Access
- The full `List` type (insertions/deletions).
- More on Doubly-Linked Lists