CSE-250 Recitation

Sept 11-Sept 12: PA1, Inequalities, Logarithms, and Bounds



Questions?

- Java?
- PA1?
- Asymptotic Notation?
- Runtime Analysis?

PA1: Getting Started

- **PA1** will revolve around linked lists and how to implement them
- We will start **PA1** by writing tests
- Why Test Driven Development?
 - Deepens your understanding of the problem
 - Enables you to test your code without submitting to Autolab
 - Writing code before thinking about the problem will lead to disaster

PA1: Getting Started

- Remember, understanding the expected behavior of each method is more important than how to make your implementation when writing tests
- Some of the best tests are going to be written by asking "What situations could break my code"
- For our first exercise, we will try to come up with some good linked lists for testing
 - Side note: how can we make these lists without relying on methods like insert

Inequalities Cheat Sheet

 $f(n) \ge g(n) \text{ is true if } f(n)/a \ge g(n)/a \text{ (for any } a > 0)$ $f(n) \ge g(n) \text{ is true if } f(n)*a \ge g(n)*a \text{ (for any } a > 0)$ $x + a \ge y + b \text{ is true if } x \ge y \text{ and } a \ge b \text{ (for any } a, b)$ $x \ge y \text{ is true if } x \ge a \text{ and } a \ge y \text{ (for any } a)$

Show that there is some **c** for which...

 $12 \log(10 \times 2^n) \le c n$

... for all **n** greater than some **n**₀

Show that there is some **c** for which...

 $n^2 + n \log(n) \le c 2^n$

... for all **n** greater than some **n**₀

Let **f**(**x**) be a function defined as follows:

- If **x** is odd, **f(x) = 10x**
- If x is even, **f**(x) = 100x²

Show that there is some $\mathbf{c}_{\mathsf{low}}$ and $\mathbf{c}_{\mathsf{high}}$ for which...

$$f(n) \ge c_{low} n$$
 and $f(n) \le c_{high} n^2$

... for all **n** greater than some $\mathbf{n_0}$

Let **f**(**x**) be a function defined as follows:

- If x is odd, f(x) = 10x
- If **x** is even, **f**(**x**) = 100**x**²

Is there a c_{low} and c_{high} for which... $f(n) \ge c_{low} n^2$ and $f(n) \le c_{high} n$... for all **n** greater than some n_0

For each of the functions to the right:

- 1. Compute the closed form of the summation
- Provide a big-⊖ bound for the function.
- 3. Arrange them in order of complexity class.

$$egin{aligned} f_1(n) &= \sum_{i=1}^{5n} (n^2 2^i) \ f_2(n) &= \sum_{i=\frac{3}{4}n}^n i \ f_3(n) &= \sum_{i=1}^n \sum_{j=i}^n 2 \ f_4(n) &= \sum_{i=1}^{\log(n)} (5 \cdot 2^i + n) \ f_5(n) &= \sum_{i=5}^{\log(n)+5} (3 \cdot i \cdot \log(n)) \ f_6(n) &= \sum_{i=2}^4 \log(i) \end{aligned}$$

- Let a, b, c, n > 0
- Exponent Rule: $\log(n^a) = a \log(n)$
- Product Rule: log(an) = log(a) + log(n)
- Division Rule: $\log(\frac{n}{a}) = \log(n) \log(a)$
- Change of Base from *b* to $c: \log_b(n) = \frac{\log_c(n)}{\log_c(b)}$
- Log/Exponent are Inverses: $b^{\log_b(n)} = \log_b(b^n) = n$

Summation Cheat Sheet

$$1. \sum_{i=j}^{k} c = (k - j + 1)c$$

$$2. \sum_{i=j}^{k} (cf(i)) = c \sum_{i=j}^{k} f(i)$$

$$3. \sum_{i=j}^{k} (f(i) + g(i)) = \left(\sum_{i=j}^{k} f(i)\right) + \left(\sum_{i=j}^{k} g(i)\right)^{j}$$

$$4. \sum_{i=j}^{k} (f(i)) = \left(\sum_{i=\ell}^{k} (f(i))\right) - \left(\sum_{i=\ell}^{j-1} (f(i))\right) \text{ (for any } \ell < j\text{)}$$

$$5. \sum_{i=j}^{k} f(i) = f(j) + f(j + 1) + \dots + f(k - 1) + f(k)$$

$$6. \sum_{i=j}^{k} f(i) = f(j) + \dots + f(\ell - 1) + \left(\sum_{i=\ell}^{k} f(i)\right) \text{ (for any } j < \ell \le k\text{)}$$

$$7. \sum_{i=j}^{k} f(i) = \left(\sum_{i=j}^{\ell} f(i)\right) + f(\ell + 1) + \dots + f(k) \text{ (for any } j \le \ell < k\text{)}$$

$$8. \sum_{i=1}^{k} i = \frac{k(k+1)}{2}$$

$$9. \sum_{i=0}^{k} 2^{i} = 2^{k+1} - 1$$