CSE 250 Data Structures

Dr. Eric Mikida epmikida@buffalo.edu 208 Capen Hall

Lec 03: Math Refresher

Announcements and Feedback

- Make sure you are on Piazza and AutoLab
- Academic Integrity Quiz due 9/8 @ 11:59PM (MUST GET 100%)
- PA0 due 9/8 @ 11:59PM (MUST GET 100%)
- WA1 is out now, due 9/8 @ 11:59PM

Today's Topics

- Summations
- Logarithms

Summations

$\sum_{i=j}^{k} f(i) = f(j) + f(j+1) + \dots + f(k)$

If **c** is a constant (with respect to **i**)



If **c** is a constant (with respect to **i**)

$$\sum_{i=j}^{k} c = (c + c + ... + c)$$

$$= (k - j + 1) \cdot c$$

If **c** is a constant and **f**(**i**) is a function of **i**:

$$\sum_{i=j}^{k} cf(i) = (cf(j) + cf(j+1) + \dots + cf(k))$$

If **c** is a constant and **f**(**i**) is a function of **i**:

$$\sum_{i=j}^{k} cf(i) = (cf(j) + cf(j+1) + \dots + cf(k))$$

$$=c(f(j)+f(j+1)+\ldots+f(k))$$

If **c** is a constant and **f**(**i**) is a function of **i**:

$$\sum_{i=j}^{k} cf(i) = (cf(j) + cf(j+1) + \dots + cf(k))$$

$$= c(f(j) + f(j+1) + ... + f(k))$$

= $c \sum_{i=j}^{k} f(i)$

If *f*(*i*) and *g*(*i*) are functions of *i*:

$\sum_{i=j}^{k} f(i) + g(i) = (f(j) + g(j)) + (f(j+1) + g(j+1)) + \dots + (f(k) + g(k))$

If *f(i)* and *g(i)* are functions of *i*:

$\sum_{i=j}^{k} f(i) + g(i) = (f(j) + g(j)) + (f(j+1) + g(j+1)) + \dots + (f(k) + g(k))$

 $= (f(j) + f(j+1) + \ldots + f(k)) + (g(j) + g(j+1) + \ldots + g(k))$

If *f*(*i*) and *g*(*i*) are functions of *i*:

$$\sum_{i=j}^{k} f(i) + g(i) = (f(j) + g(j)) + (f(j+1) + g(j+1)) + \dots + (f(k) + g(k))$$

$$= (f(j) + f(j+1) + \dots + f(k)) + (g(j) + g(j+1) + \dots + g(k))$$

$$= \left(\sum_{i=j}^k f(i)\right) + \left(\sum_{i=j}^k g(i)\right)$$

lf*j < l < k*:

$$\sum_{i=j}^k f(i) = f(j) + \ldots + f(k)$$

lf*j < l < k*:

$$\sum_{i=j}^k f(i) = f(j) + \ldots + f(k)$$

$$= f(j) + \dots + f(l-1) + f(l) + \dots + f(k)$$

lf *j < l < k*:

$$\sum_{i=j}^{k} f(i) = f(j) + \dots + f(k)$$

= $f(j) + \dots + f(l-1) + f(l) + \dots + f(k)$
= $\left(\sum_{i=j}^{l-1} f(i)\right) + \left(\sum_{i=l}^{k} f(i)\right)$

lf*j < l < k*:

 $\left(\sum_{i=i}^{k} f(i)\right) = \left(\sum_{i=i}^{l-1} f(i)\right) + \left(\sum_{i=l}^{k} f(i)\right)$

lf*j < l < k*:



Subtract to other side

lf*j < l < k*:



Series

Some common closed form solutions:

$$\sum_{i=1}^{k} i = \frac{k(k+1)}{2}$$
$$\sum_{i=0}^{k} 2^{i} = 2^{k+1} - 1$$

Summary

- The previous rules will always be provided on WAs and exams
- Usually the goal will be to reduce some complicated summation to a simpler form without a summation
 - Some of the rules get rid of summations
 - Some allow you to manipulate summations/bounds so that you can apply rules that get rid of summations
- Be cognizant of what variables are constant with respect to the summation variable and which one aren't

$$a \cdot n = \underbrace{a + a + \ldots + a}_{\checkmark}$$

a added together **n** times

$$a \cdot n = \underbrace{a + a + \dots + a}_{\textbf{a} \text{ added together } \textbf{n} \text{ times}}$$

$$a^n = \underbrace{a \cdot a \cdot \ldots \cdot a}_{}$$

a multiplied together **n** times





 $\log_a(b)$ = the number of times you multiply a together to get b

$$\log_2(32) = 5$$

 $\log_3(27) = 3$
 $\log_2(\frac{1}{8}) = -3$
 $\log_2(2^{10}) = 10$

Logarithm is the inverse exponent

$$b^{\log_b(n)} = n = \log_b(b^n)$$

Let's say $n = a \cdot b$

Let's say $n = a \cdot b$

$$a = 2 \cdot \ldots \cdot 2$$
 $b = 2 \cdot \ldots \cdot 2$

Let's say $n = a \cdot b$

$$a = \underbrace{2 \cdot \ldots \cdot 2}_{\log_2(a) \text{ times}}$$

$$b = \underbrace{2 \cdot \ldots \cdot 2}_{\log_2(b) \text{ times}}$$

Let's say $n = a \cdot b$

$$a = 2 \cdot \ldots \cdot 2 \qquad \qquad b = 2 \cdot \ldots \cdot 2$$

$$n = \underbrace{2 \cdot \ldots \cdot 2}_{\log_2(n) \text{ times}} = \underbrace{2 \cdot \ldots \cdot 2}_{\log_2(a) \text{ times}} \cdot \underbrace{2 \cdot \ldots \cdot 2}_{\log_2(b) \text{ times}}$$

Let's say $n = a \cdot b$

How are $\log_2(n)$, $\log_2(a)$, and $\log_2(a)$ related?

$$a = 2 \cdot \ldots \cdot 2 \qquad \qquad b = 2 \cdot \ldots \cdot 2$$

$$n = \underbrace{2 \cdot \ldots \cdot 2}_{\log_2(n) \text{ times}} = \underbrace{2 \cdot \ldots \cdot 2}_{\log_2(a) \text{ times}} \cdot \underbrace{2 \cdot \ldots \cdot 2}_{\log_2(a) \text{ times}}$$

 $\log_2(n) = \log_2(ab) = \log_2(a) + \log_2(b)$

$$\log_2(a^n) = \log_2(a \cdot \dots \cdot a)$$

 $\log_2(a^n) = \log_2(\overbrace{a \cdot \ldots \cdot a}^{n \text{ times}})$

$$\begin{aligned} \mathbf{n} \text{ times} \\ \log_2(a^n) &= \log_2(\overrightarrow{a \cdot \ldots \cdot a}) \\ &= \log_2(a) + \log_2(\overrightarrow{a \cdot \ldots \cdot a}) \end{aligned}$$

$$log_2(a^n) = log_2(\overrightarrow{a \cdot \ldots \cdot a}) \quad \textbf{n-1 times}$$
$$= log_2(a) + log_2(\overrightarrow{a \cdot \ldots \cdot a}) \quad \textbf{n-2 times}$$
$$= log_2(a) + log_2(a) + log_2(\overrightarrow{a \cdot \ldots \cdot a})$$



$$\log_2(a^n) = \log_2(a \cdot \dots \cdot a)$$

=
$$\log_2(a) + \log_2(a \cdot \dots \cdot a)$$

=
$$\log_2(a) + \log_2(a) + \log_2(a \cdot \dots \cdot a)$$

=
$$\log_2(a) + \dots + \log_2(a)$$

=
$$n \cdot \log_2(a)$$

$$\log_2(\frac{a}{b}) = \log_2(a \cdot \frac{1}{b})$$

$$\log_2(\frac{a}{b}) = \log_2(a \cdot \frac{1}{b})$$
$$= \log_2(a) + \log_2(\frac{1}{b})$$

]

$$\log_2(\frac{a}{b}) = \log_2(a \cdot \frac{1}{b})$$
$$= \log_2(a) + \log_2(\frac{1}{b})$$
$$= \log_2(a) + \log_2(b^{-1})$$

$$\log_2\left(\frac{a}{b}\right) = \log_2\left(a \cdot \frac{1}{b}\right)$$
$$= \log_2(a) + \log_2\left(\frac{1}{b}\right)$$
$$= \log_2(a) + \log_2(b^{-1})$$
$$= \log_2(a) - \log_2(b)$$

 $b^m = n$

$$b^m = n$$
$$\log_c(b^m) = \log_c(n)$$

$$b^{m} = n$$
$$\log_{c}(b^{m}) = \log_{c}(n)$$
$$m \cdot \log_{c}(b) = \log_{c}(n)$$

$$b^{m} = n$$
$$\log_{c}(b^{m}) = \log_{c}(n)$$
$$m \cdot \log_{c}(b) = \log_{c}(n)$$
$$m = \frac{\log_{c}(n)}{\log_{c}(b)}$$

$$b^{m} = n$$
$$\log_{c}(b^{m}) = \log_{c}(n)$$
$$m \cdot \log_{c}(b) = \log_{c}(n)$$
$$m = \frac{\log_{c}(n)}{\log_{c}(b)}$$
$$\log_{b}(n) = \frac{\log_{c}(n)}{\log_{c}(b)}$$

Summary

Exponent Rule $log(n^a) = a log(n)$ Product Rulelog(ab) = log(a) + log(b)Division Rule $log\left(\frac{a}{b}\right) = log(a) - log(b)$ Change of Base $log_b(n) = \frac{log_c(n)}{log_c(b)}$ Inverse $b^{log_b(n)} = log_b(b^n) = n$

* for this class, always assume base 2 unless otherwise stated *