

CSE 250

Data Structures

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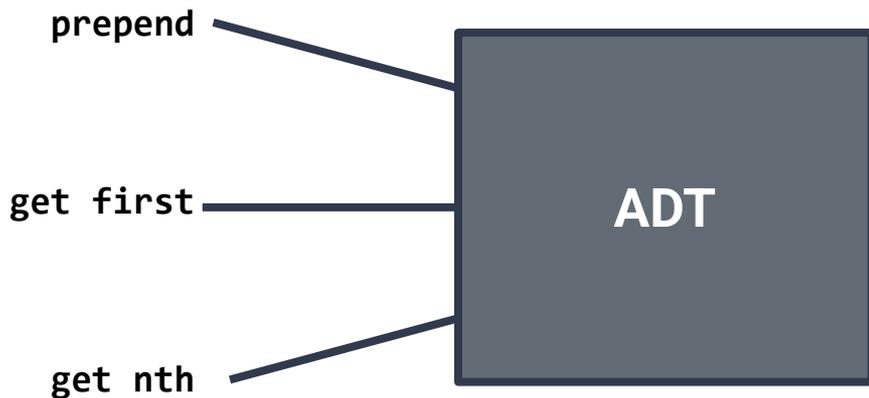
Lec 05: Intro to Complexity

Announcements and Feedback

- AI Quiz, PA0, WA1 due Sunday @ 11:59PM
- PA1 releases Monday

Thought Experiment

Often, many data structures can satisfy a given ADT...how do you choose?



Thought Experiment

Data Structure 1

- Very fast **prepend**, **get first**
- Very slow **get nth**

Data Structure 2

- Very fast **get nth**, **get first**
- Very slow **prepend**

Data Structure 3

- Very fast **get nth**, **get first**
- Occasionally slow **prepend**

Which is better?

Thought Experiment

Data Structure 1 (LinkedList)

- Very fast prepend, get first
- Very slow get nth

Data Structure 2 (Array)

- Very fast get nth, get first
- Very slow prepend

Data Structure 3 (ArrayList...in reverse)

- Very fast get nth, get first
- Occasionally slow prepend

Which is better?

IT DEPENDS!

Thought Experiment

Data Structure 1 (LinkedList)

- Very **fast** prepend, get first
- Very **slow** get nth

Data Structure 2 (ArrayList)

- Very **fast** get nth, get first
- Very **slow** prepend

Data Structure 3 (ArrayList...in reverse)

- Very **fast** get nth, get first
- Occasionally **slow** prepend

What is "fast"? "slow"?

Attempt #1: Wall-clock time?

- What is fast?
 - 10s? 100ms? 10ns?
 - ...it depends on the task
- Algorithm vs Implementation
 - Compare Grace Hopper's implementation to yours
- What machine are you running on?
 - Your old laptop? A lab machine? The newest, shiniest processor on the market?
- What bottlenecks exist? CPU vs IO vs Memory vs Network...

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Wall-clock time is not terribly useful... 8

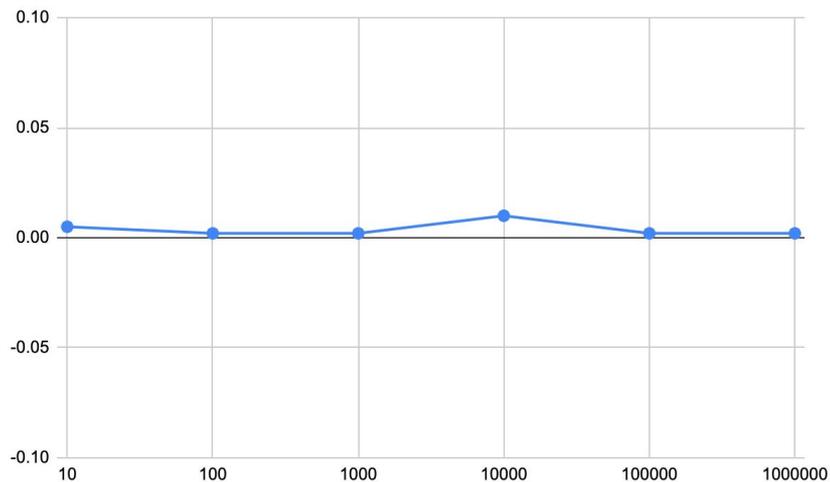
Analysis Checklist

1. Don't think in terms of wall-time, think in terms of “number of steps”

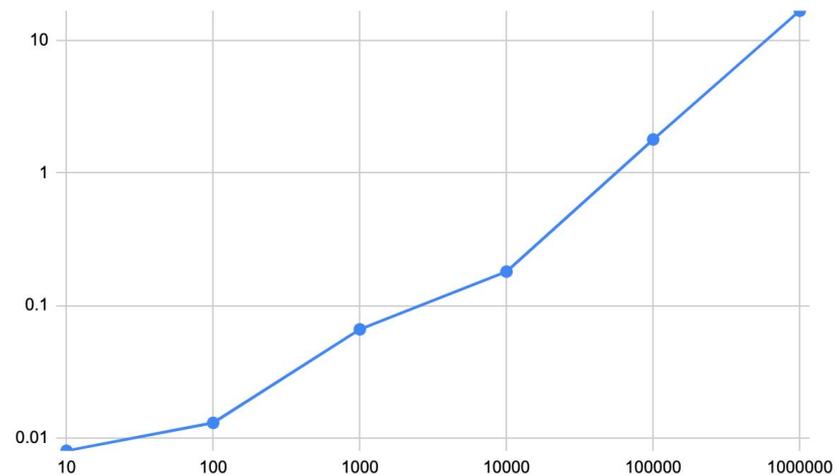
Let's do a quick demo...

Comparing Random Access for Array vs List

Array



List

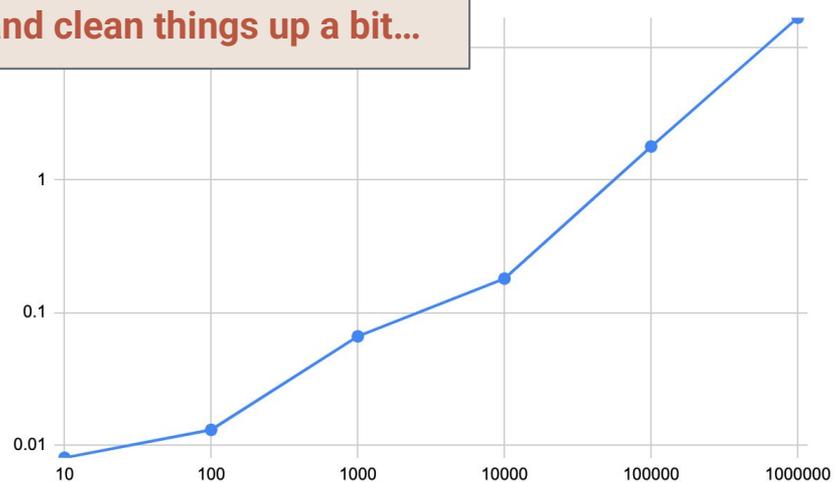


Comparing Random Access for Array vs List

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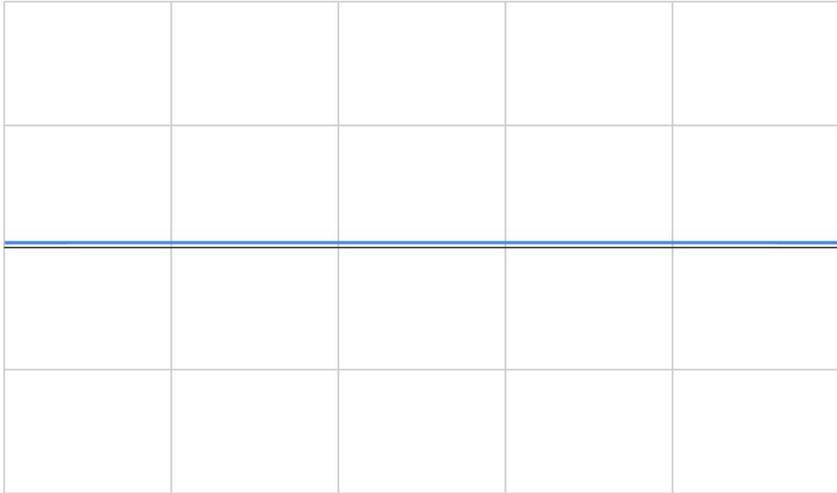
List



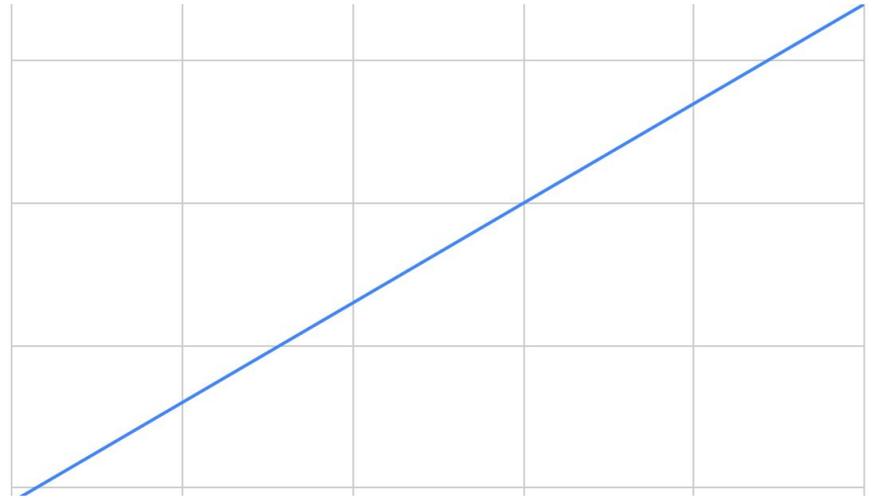
Let's ignore the specific numbers and clean things up a bit...

Comparing Random Access for Array vs List

Array



List

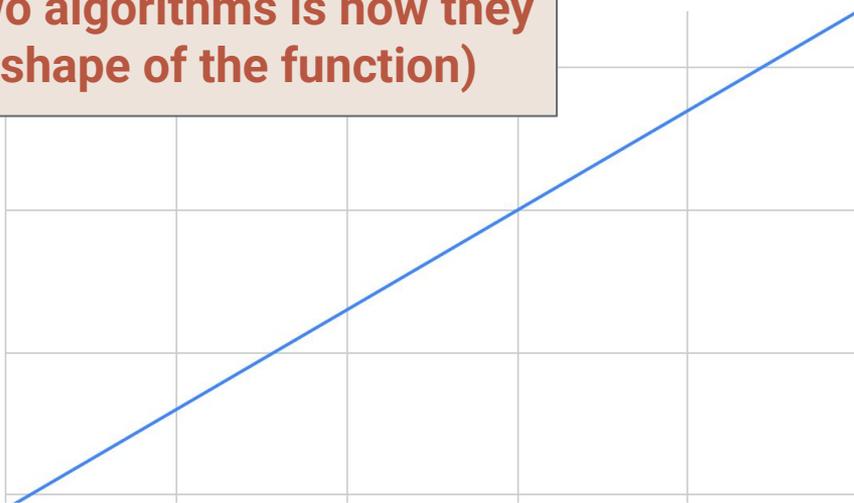
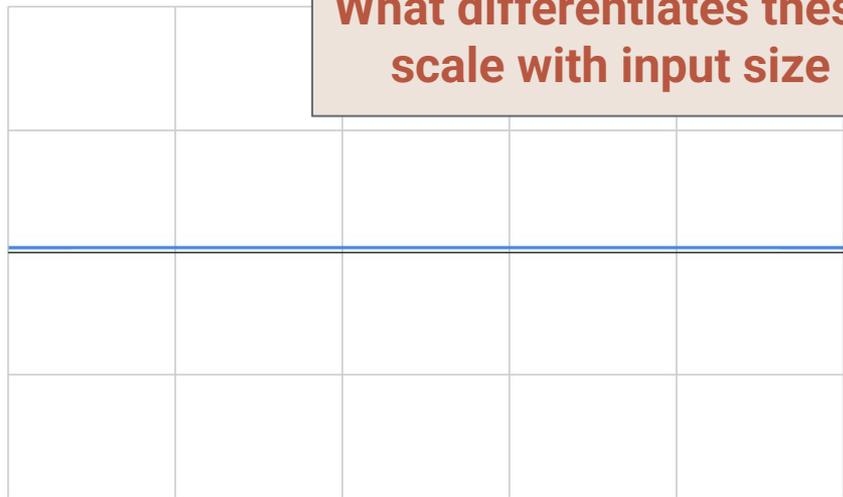


Comparing Random Access for Array vs List

Array

List

What differentiates these two algorithms is how they scale with input size (the shape of the function)



Analysis Checklist

1. Don't think in terms of wall-time, think in terms of “number of steps”
2. **To give a useful solution, we should take “scale” into account**
 - **How does the runtime change as we change the size of the input?**

Counting Steps

```
1 public void userFullName(User[] users, int id) {  
2     User user = users[id];  
3     String fullName = user.firstName + user.lastName;  
4     return fullName;  
5 }
```

Counting Steps

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1 public void userFullName(User[] users, int id) {  
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```

How many steps does this function take?

Counting Steps

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1 public void userFullName(User[] users, int id) {  
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3     String fullName = user.firstName + user.lastName;  
4     return fullName;  
5 }
```

7 steps...ish? Maybe? What the heck is a step?

Counting Steps

```
1 public void updateUsers(User[] users) {  
2     x = 1;  
3     for(User user : users) {  
4         user.id = x;  
5         x = x + 1;  
6     }  
7 }
```

Counting Steps

```
1 public void updateUsers(User[] users) {  
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1

Counting Steps

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$$1 + \sum_{user \in users}$$

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$$1 + \sum_{user \in users} 4$$

Counting Steps

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$$1 + \sum_{user \in users} 4 = 1 + 4 \cdot |users|$$

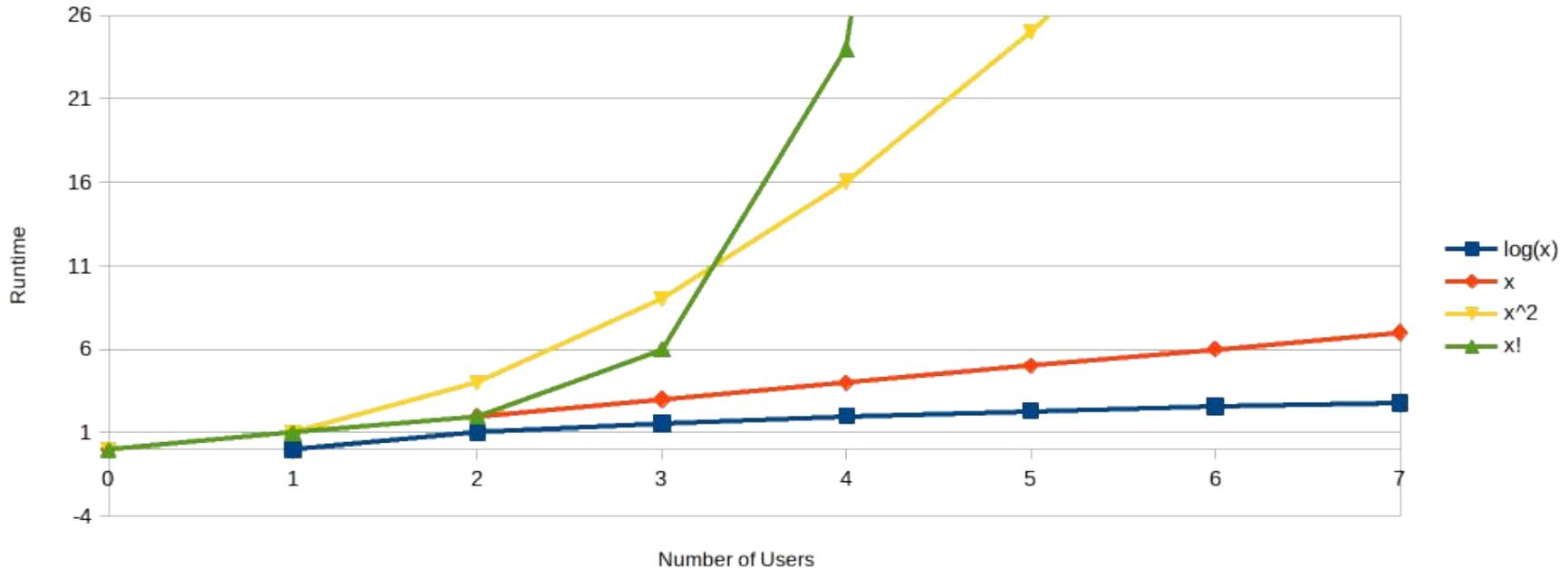
Steps to "Functions"

Now that we have number of steps* in terms of summations...
...which we can simplify (like in WA1) into mathematical functions...

We can start analyzing runtime as a function

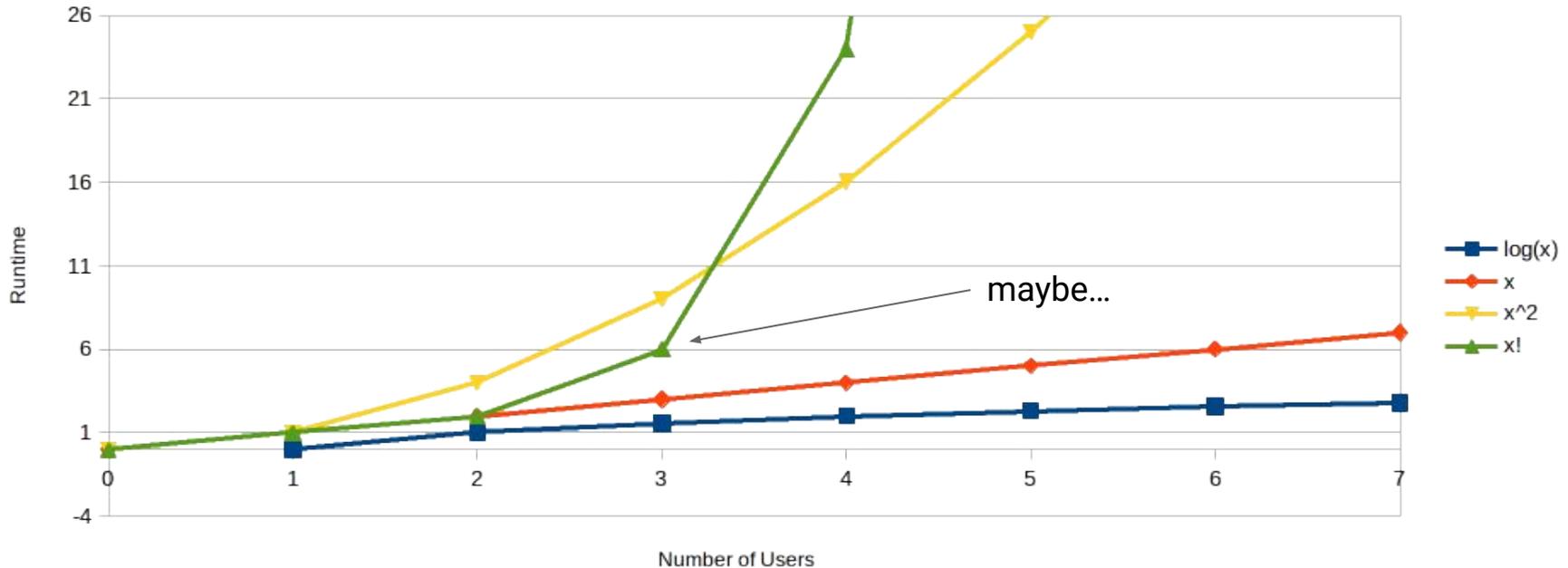
** we'll give a better definition of what a "step" is later*

Runtime as a Function



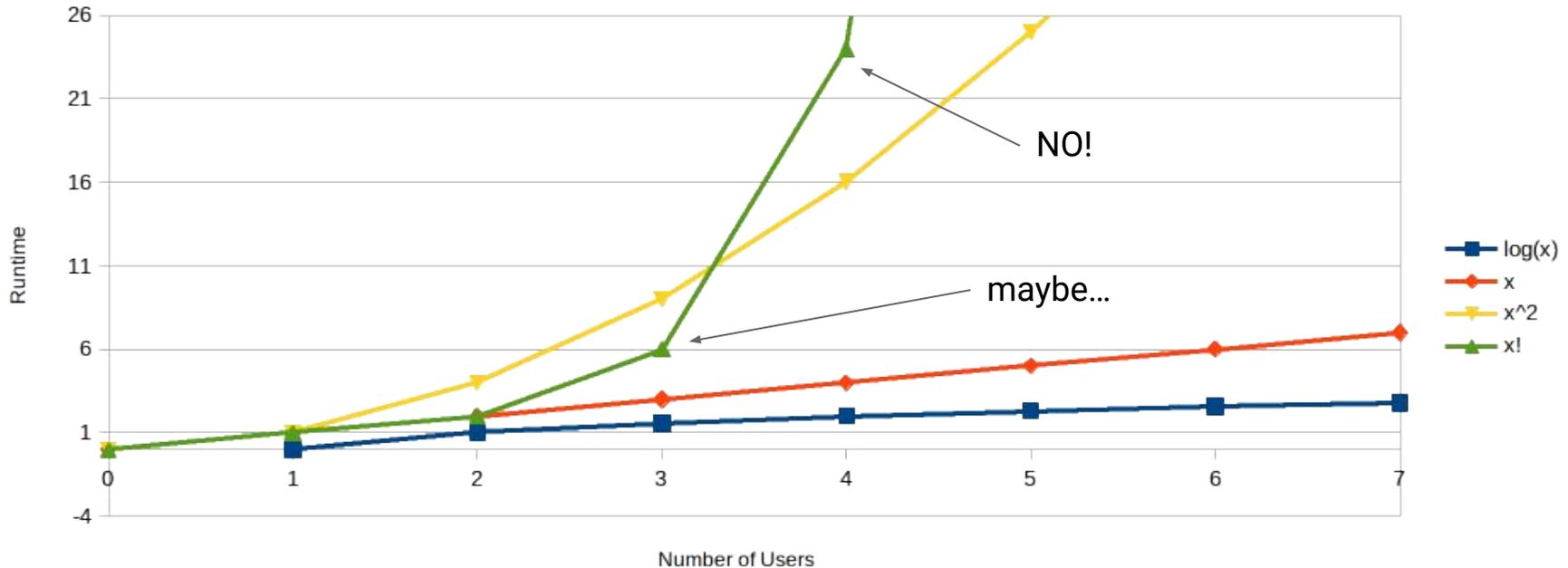
Would you consider an algorithm that takes $|\text{Users}|!$ number of steps?

Runtime as a Function



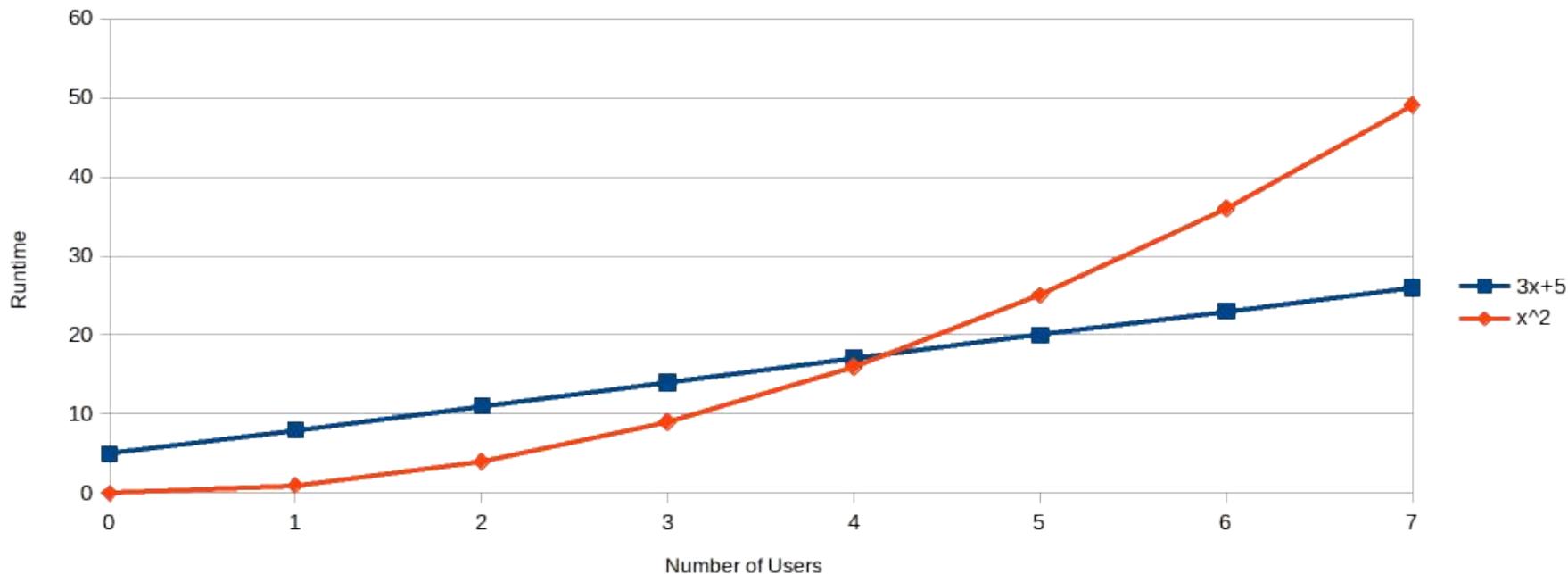
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Runtime as a Function



Would you consider an algorithm that takes $|\text{Users}|!$ number of steps?

Runtime as a Function



Which is better? $3x|\text{Users}|+5$ or $|\text{Users}|^2$

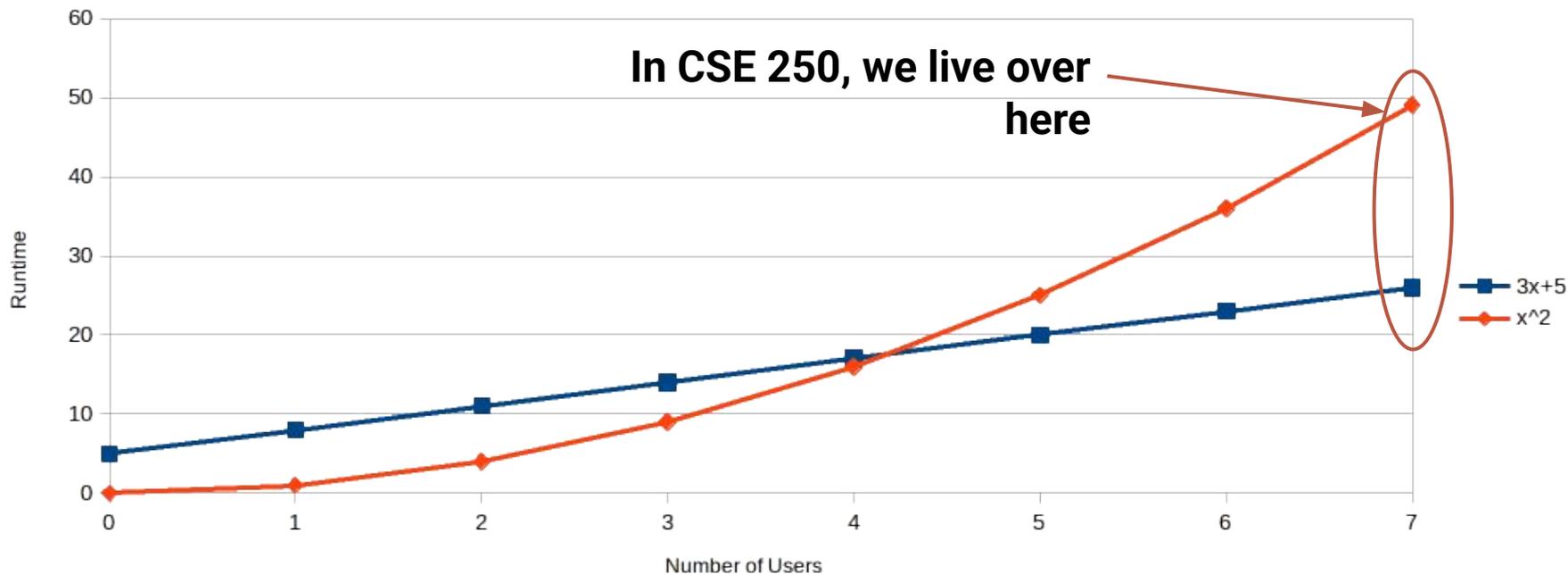
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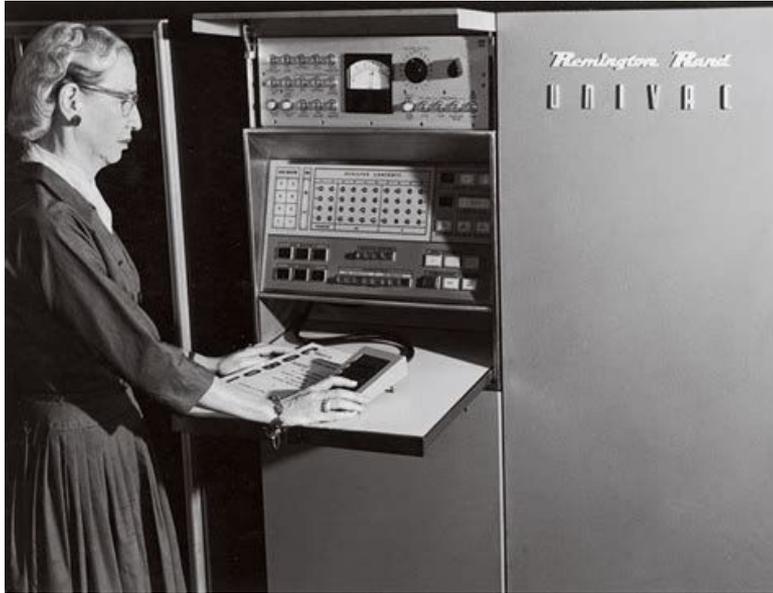
1. Don't think in terms of wall-time, think in terms of “number of steps”
2. To give a useful solution, we should take “scale” into account
 - How does the runtime change as we change the size of the input?
3. **Focus on “large” inputs**
 - **Rank functions based on how they behave at large scales**

Runtime as a Function



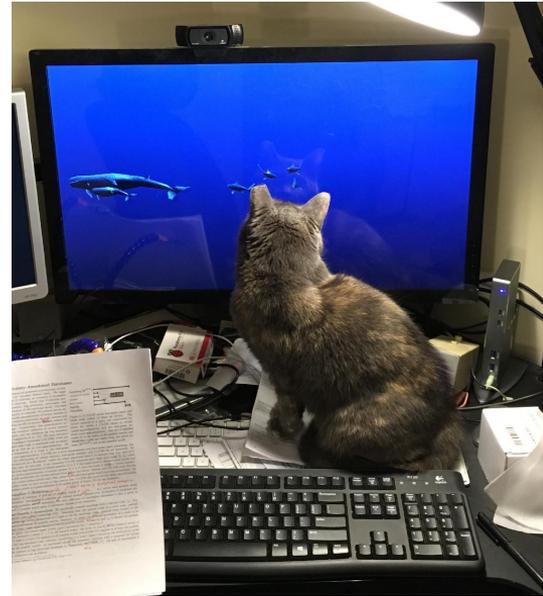
Which is better? $3x|\text{Users}|+5$ or $|\text{Users}|^2$

Goal: Ignore implementation details



Seasoned Pro Implementation

VS



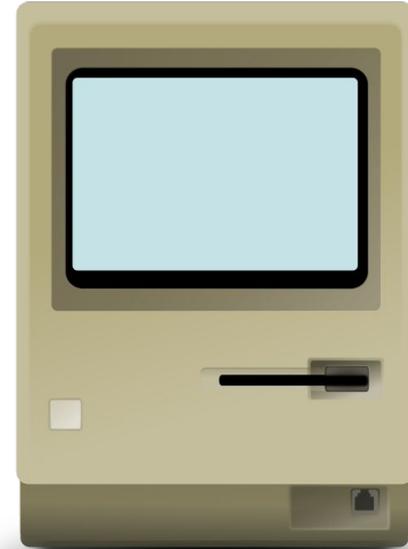
Error 23: Cat on Keyboard

Goal: Ignore execution environment



Intel i9

VS



Motorola 68000

Goal: Judge the Algorithm Itself

- How fast is a step? Don't care
 - Only count number of steps
- Can this be done in two steps instead of one?
 - “3 steps per user” vs “some number of steps per user”
 - Sometimes we don't care...sometimes we do

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 - Rank functions based on how they behave at large scales
4. **Decouple algorithm from infrastructure/implementation**
 - **Asymptotic notation...?**

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Wall-clock time is not terribly useful...³⁷

Attempt #2: Growth Functions

Not a function in code...but a mathematical function:

$$T(n)$$

n: The “size” of the input

ie: number of users, rows, pixels, etc

$T(n)$: The number of “steps” taken for input of size n

ie: 20 steps per user, where $n = |\text{Users}|$, is $20 \times n$

Some Basic Assumptions:

Problem sizes are non-negative integers

$$n \in \{0, 1, 2, 3, \dots\} = \{0\} \cup \mathbb{Z}^+$$

We can't reverse time...(obviously)

$$T(n) > 0$$

Smaller problems aren't harder than bigger problems

$$n_1 < n_2 \Rightarrow T(n_1) \leq T(n_2)$$

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$$T: \{0\} \cup \mathbb{Z}^+ \rightarrow \mathbb{R}^+$$

T is non-decreasing

First Problem...

We are still implementation dependent...

$$T_1(n) = 19n$$

$$T_2(n) = 20n$$

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$$T_1(n) = 19n$$

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Does 1 extra step per element really matter...?

Is this just an implementation detail?

First Problem...

We are still implementation dependent...

$$T_1(n) = 19n$$

$$T_2(n) = 20n$$

$$T_3(n) = 2n^2$$

T_1 and T_2 are much more “similar” to each other than they are to T_3

First Problem...

We are still implementation dependent...

$$T_1(n) = 19n$$

$$T_2(n) = 20n$$

$$T_3(n) = 2n^2$$

T_1 and T_2 are much more “similar” to each other than they are to T_3

How do we capture this idea formally?

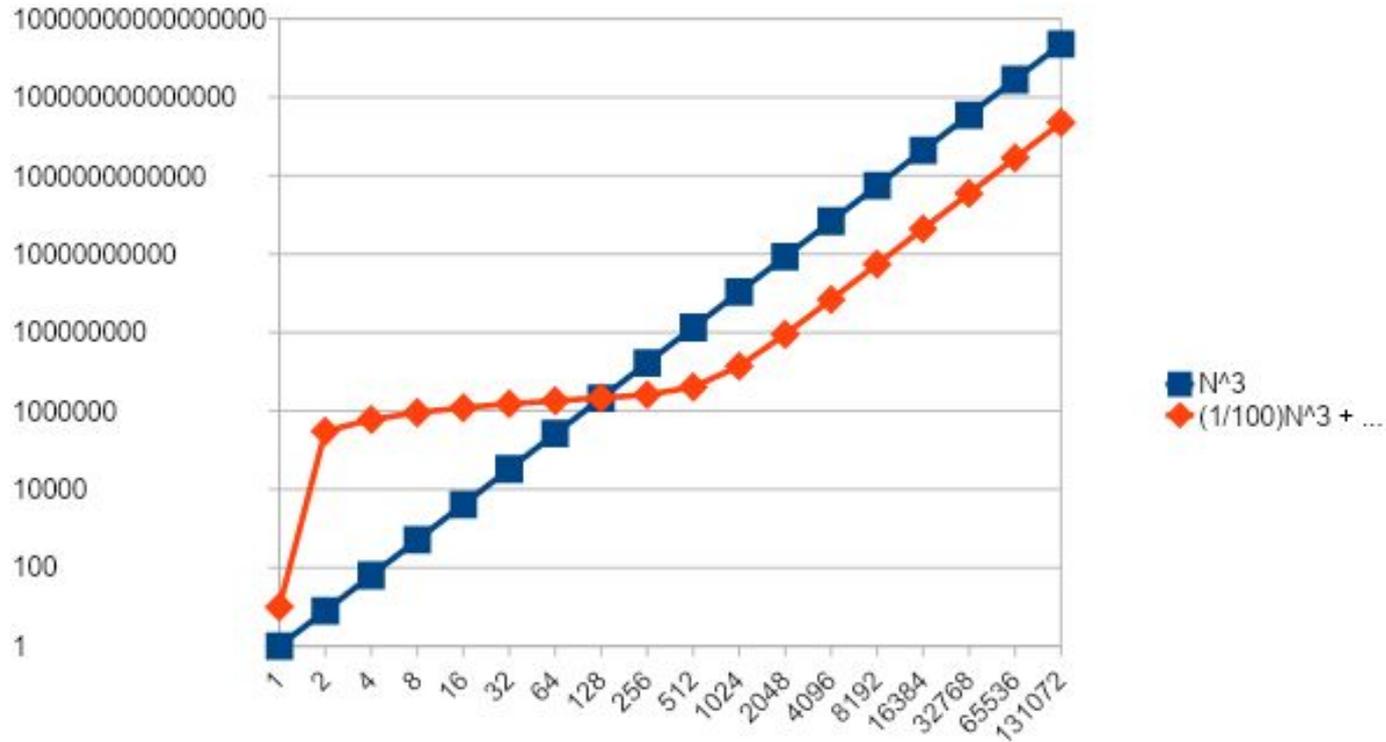
How Do We Capture Behavior at Scale?

Consider the following two functions:

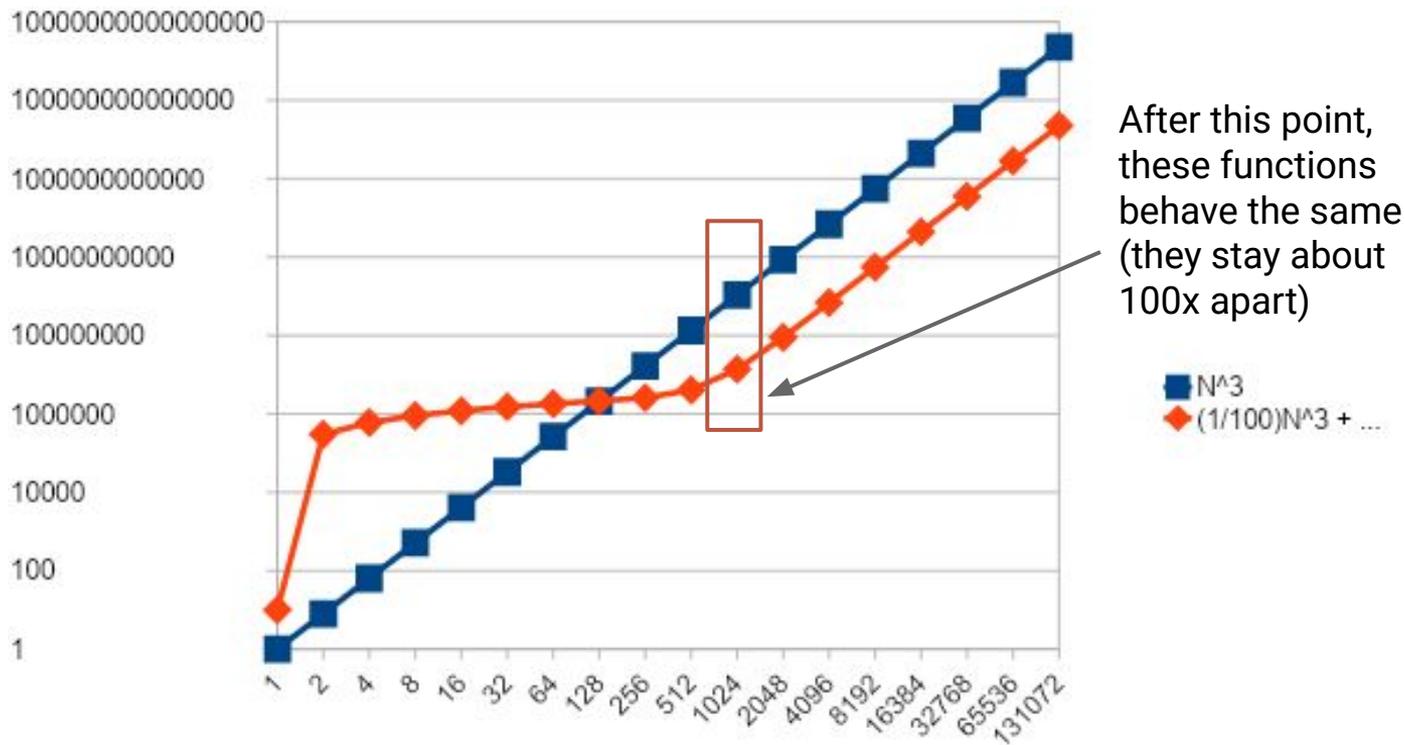
$$\frac{1}{100}n^3 + 10n + 10000000 \log(n)$$

$$n^3$$

How Do We Capture Behavior at Scale?



How Do We Capture Behavior at Scale?



Attempt #3: Asymptotic Analysis

We want to organize runtimes (growth functions) into different ***Complexity Classes***

Within the same complexity class, runtimes “behave the same”/“have the same shape” (at scale)