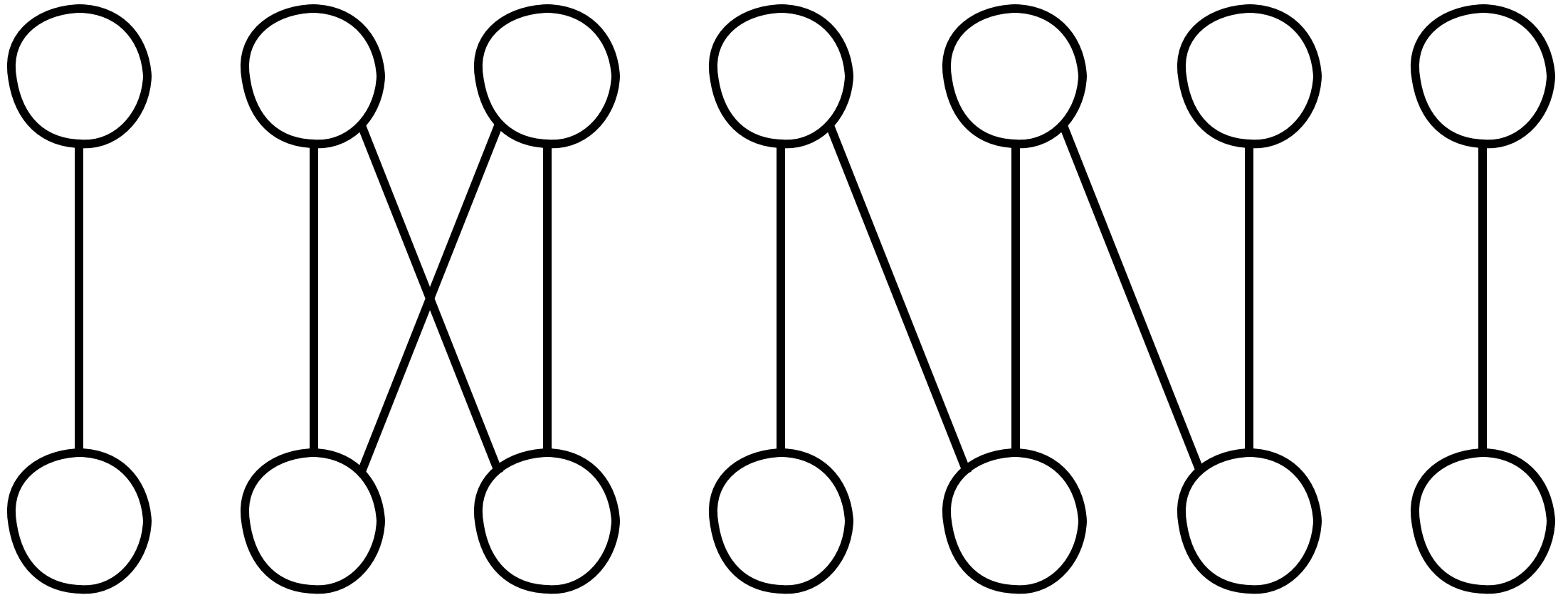


Q: How many perfect matchings in this graph?





CSE 331:

Algorithms & Complexity

“Gale-Shapley IV”

Prof. Charlie Anne Carlson (She/Her)

Lecture 7

Friday September 12th, 2025



University at Buffalo®



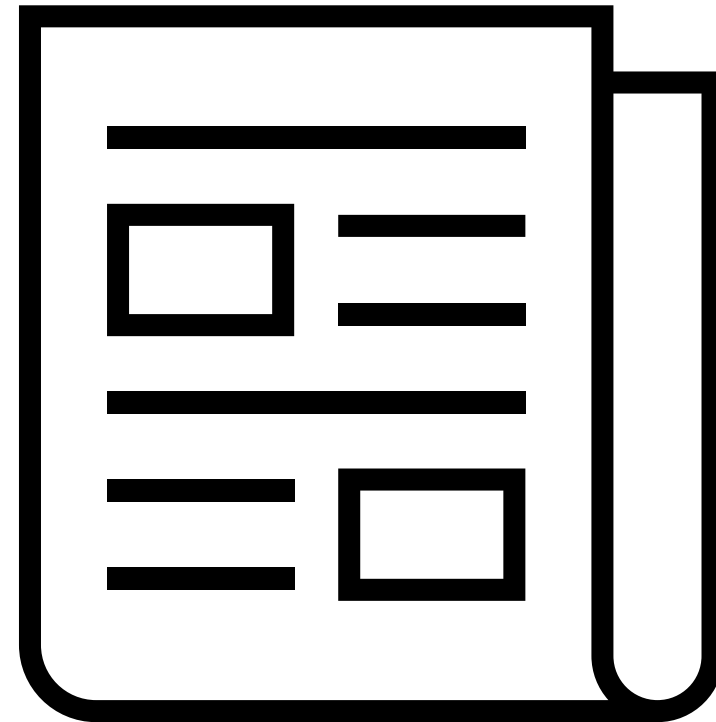
Schedule

1. Course Updates
2. Recap
3. Even More Gale-Shapely



Course Updates

- Complete Syllabus Quiz
- HW 1 Posted
 - Autolab Tomorrow
- HW 0 Being Graded
- Website Moved
- Project Signup before 19th



Q: What do we know about Gale-Shapley?

GALE-SHAPLEY (*preference lists for hospitals and students*)

INITIALIZE M to empty matching.

WHILE (some hospital h is unmatched and hasn't proposed to every student)

$s \leftarrow$ first student on h 's list to whom h has not yet proposed.

 IF (s is unmatched)

 Add $h-s$ to matching M .

 ELSE IF (s prefers h to current partner h')

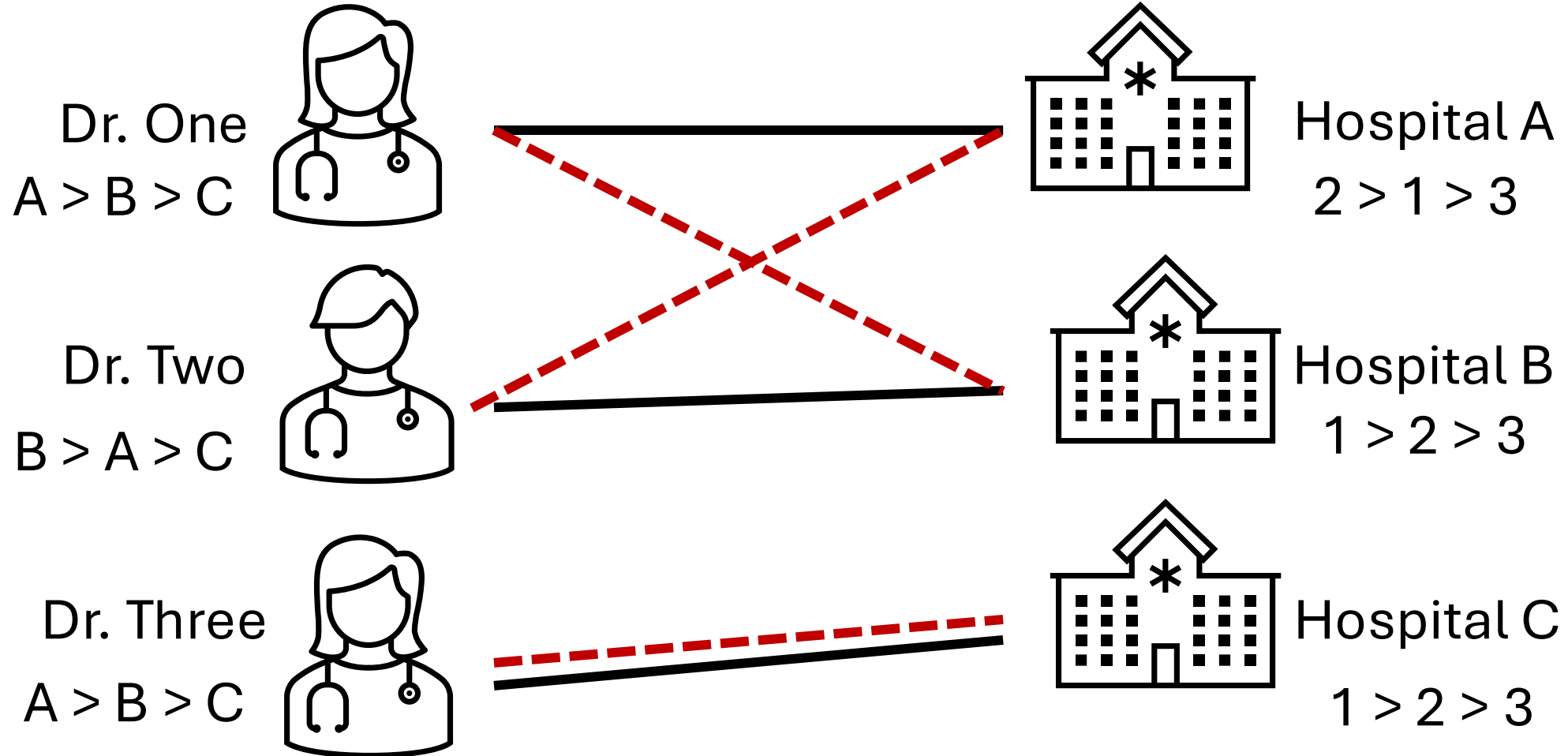
 Replace $h'-s$ with $h-s$ in matching M .

 ELSE

s rejects h .

RETURN stable matching M .

Matching I vs Matching II

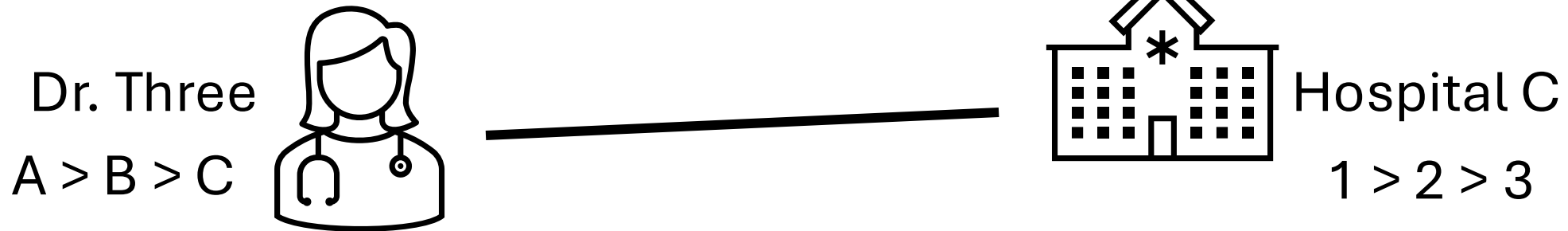


Valid Partners

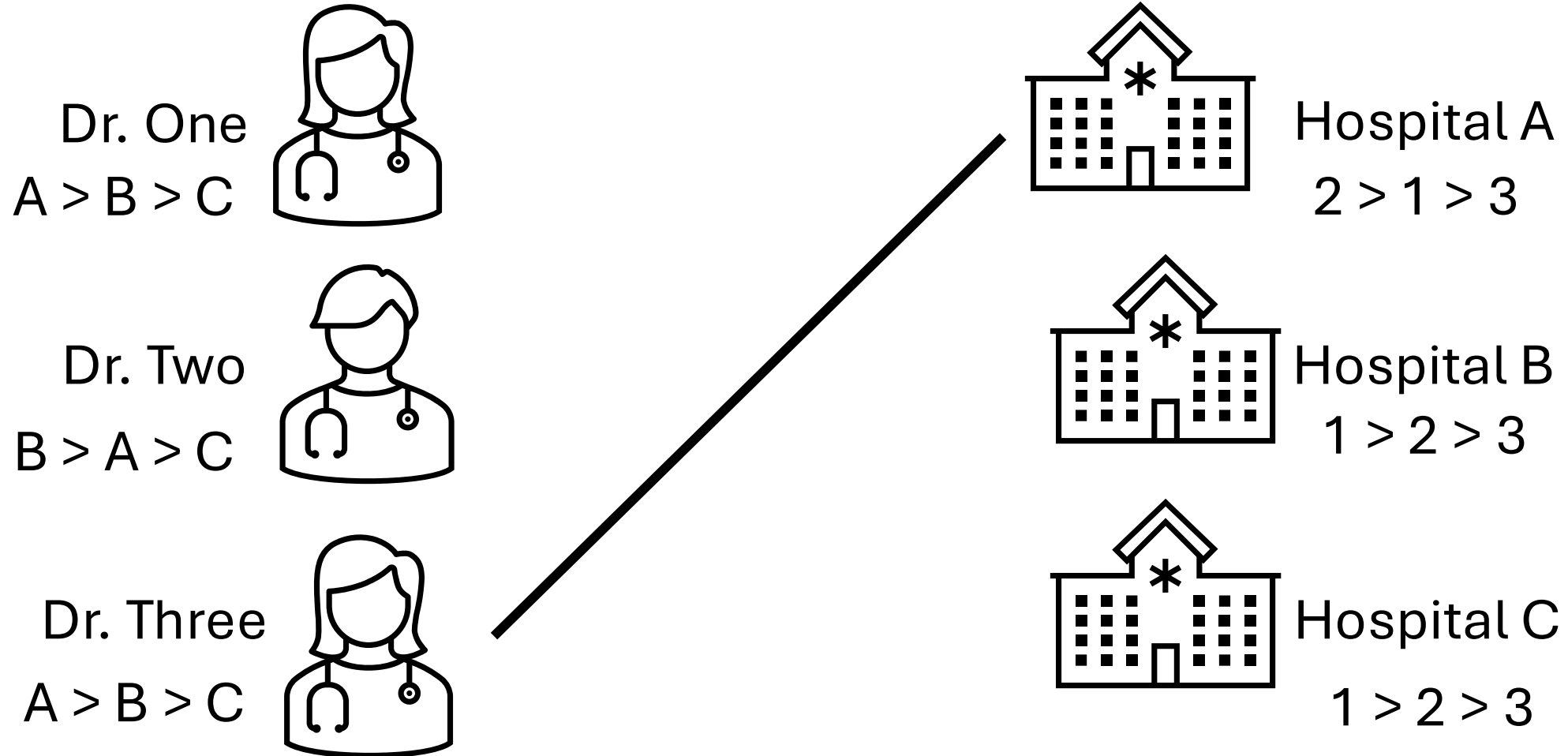
Def: We say student s is a *valid partner* of hospital h if there exists a stable matching in which s and h are matched.

Q: What would be a student's *best valid partner*?

E.g.:




Q: Is there a Stable Matching with this edge?




Q: Is there a Stable Matching with this edge?

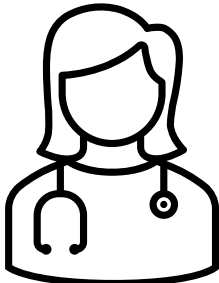
Dr. One
 $A > B > C$



Dr. Two
 $B > A > C$



Dr. Three
 $A > B > C$



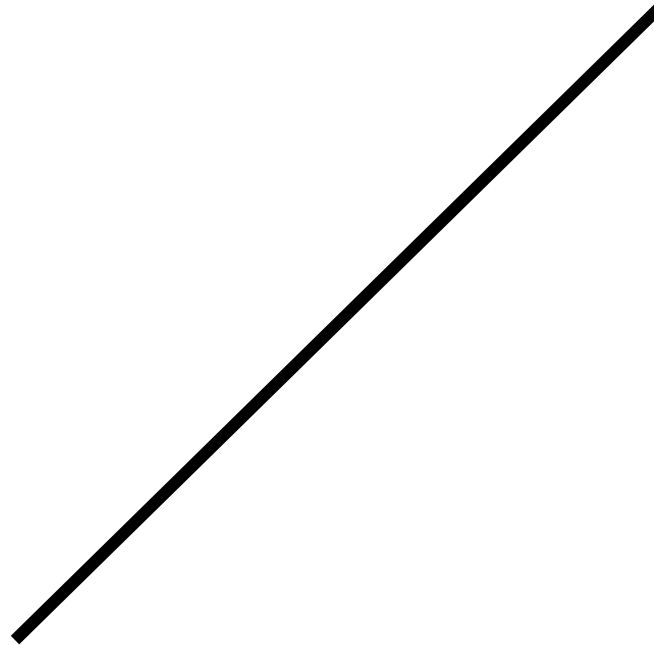
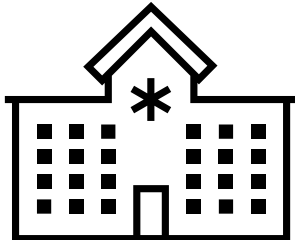
Hospital A
 $2 > 1 > 3$



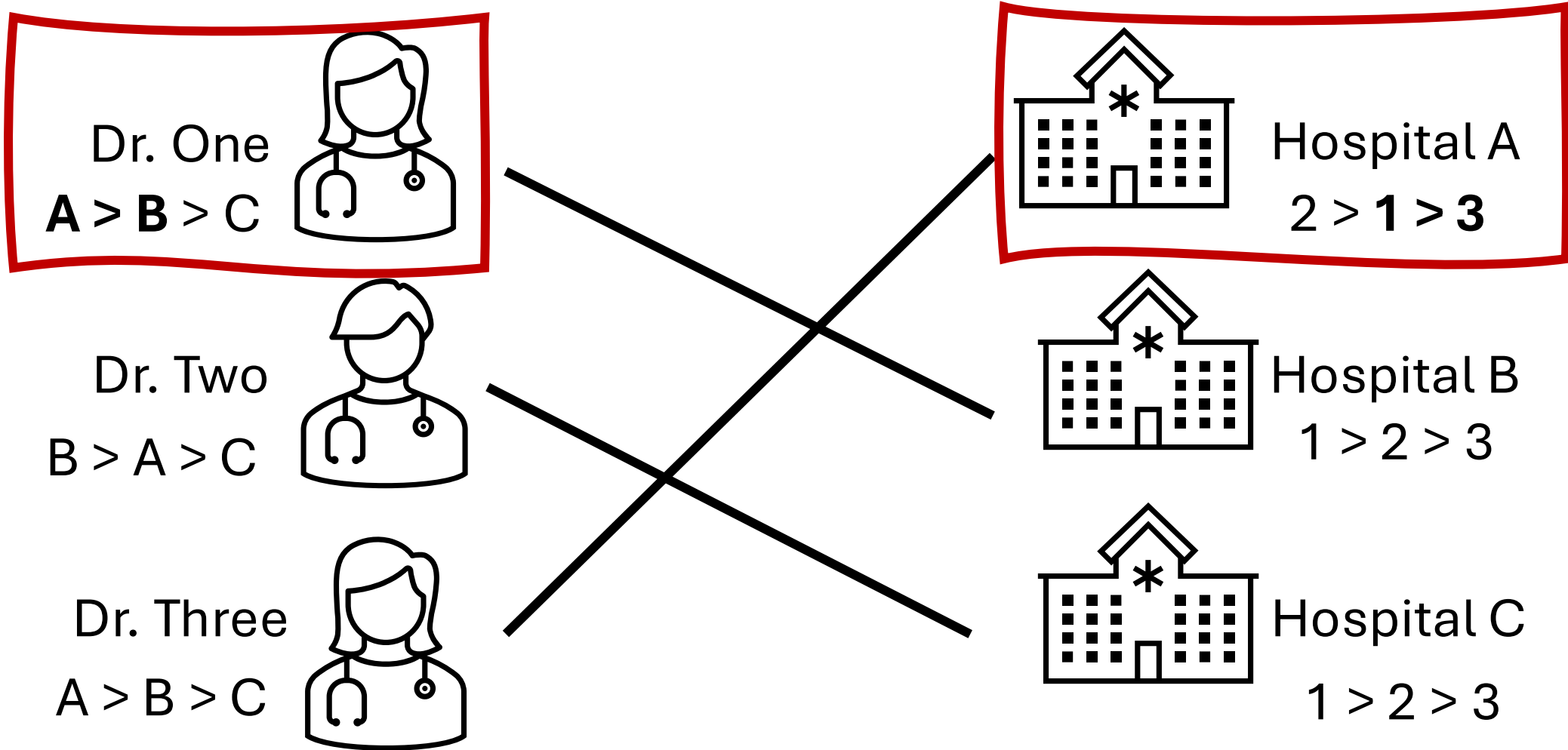
Hospital B
 $1 > 2 > 3$



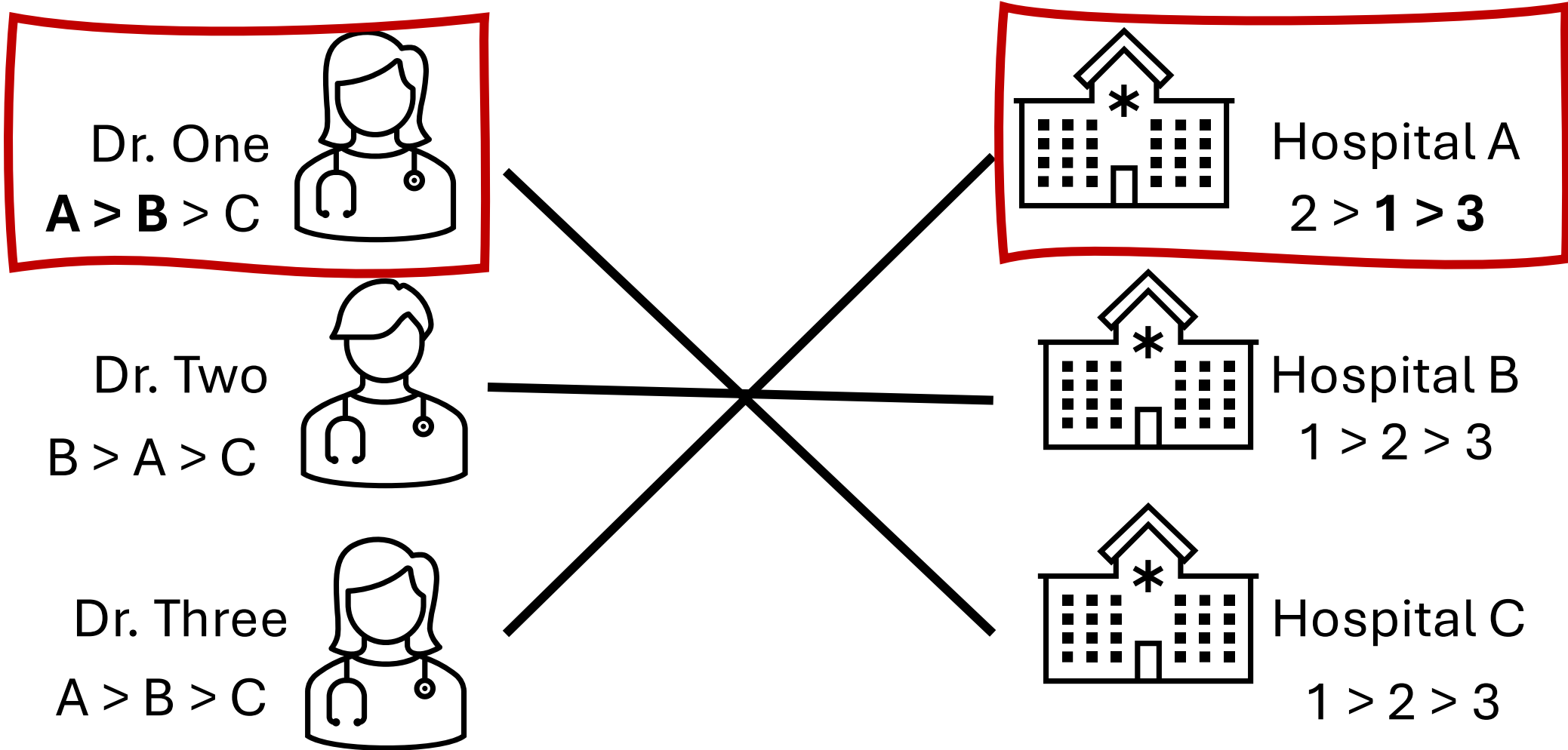
Hospital C
 $1 > 2 > 3$



Attempt One



Attempt Two



Valid Partners

Def: The *student optimal assignment* is the one in which all students are matched with their best valid partner. The *hospital optimal assignment* is the one in which all hospitals are matched with their best valid partner.

Q: Are these optimal assignments stable matchings?

Q: Can we find them?

Proof Idea Time!

Claim: The algorithm always outputs a stable matching that is hospital optimal.

Proof:

- Suppose to the contrary that the output is not hospital optimal. We will show that....

Proof Idea Time!

Proof:

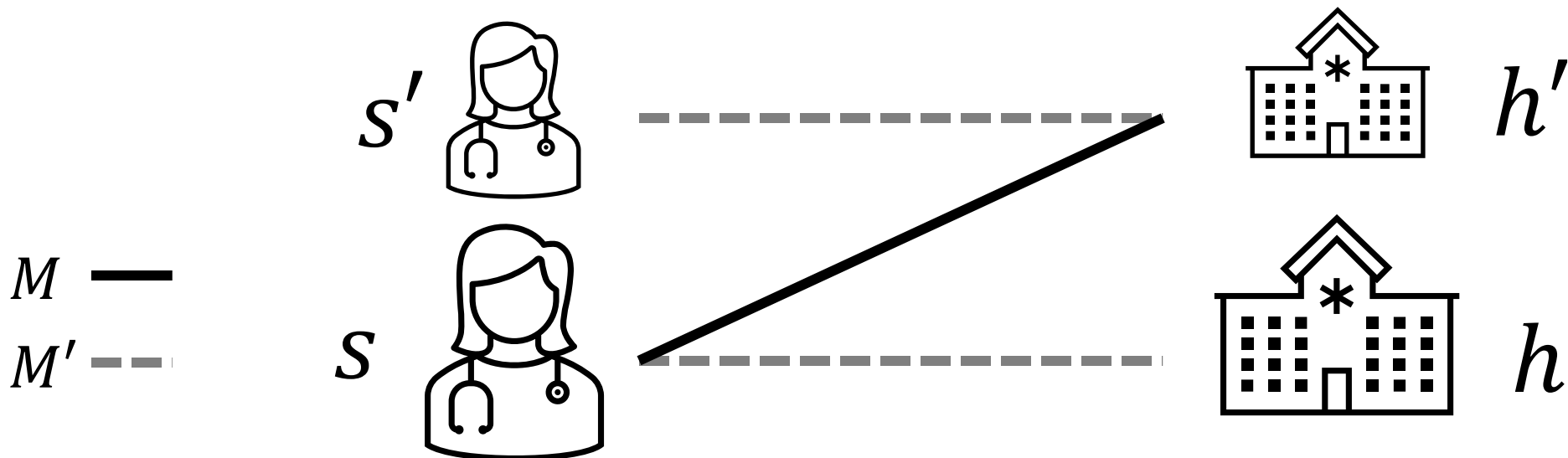
- Suppose to the contrary that the output M is not hospital-optimal.
- Then let h be the **first hospital** to be rejected by a valid partner s .
- Let M' be the stable matching with s and h paired together.



Proof Idea Time!

Proof: ...

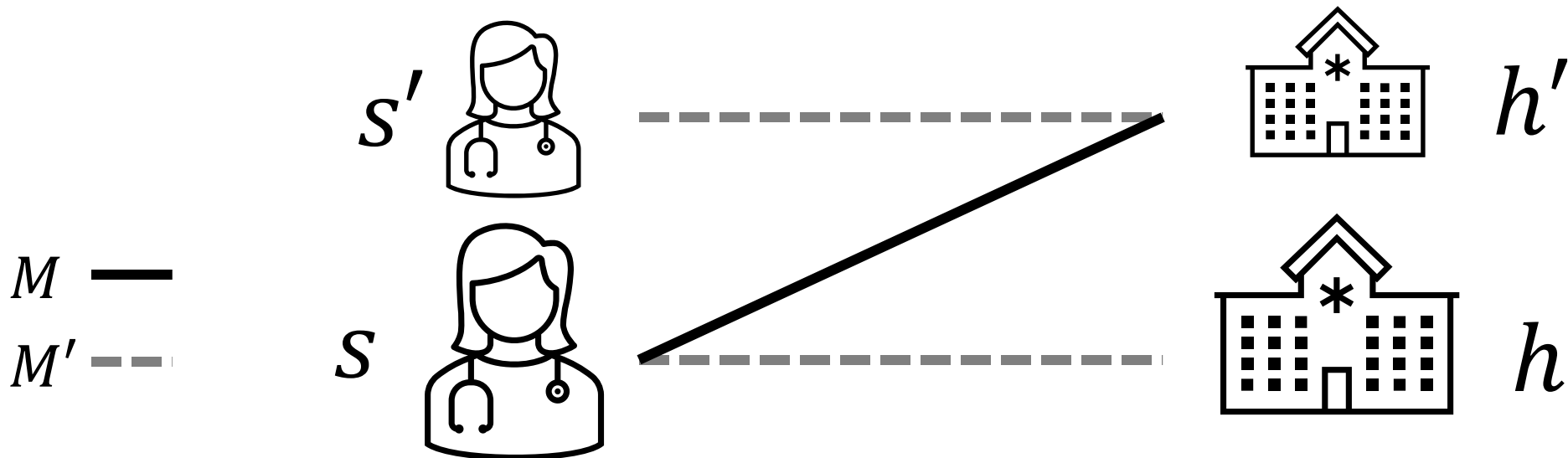
- Let h' be the hospital paired with s in the output of the algorithm. Let s' be the student paired with h' in M .



Proof Idea Time!

Proof: ...

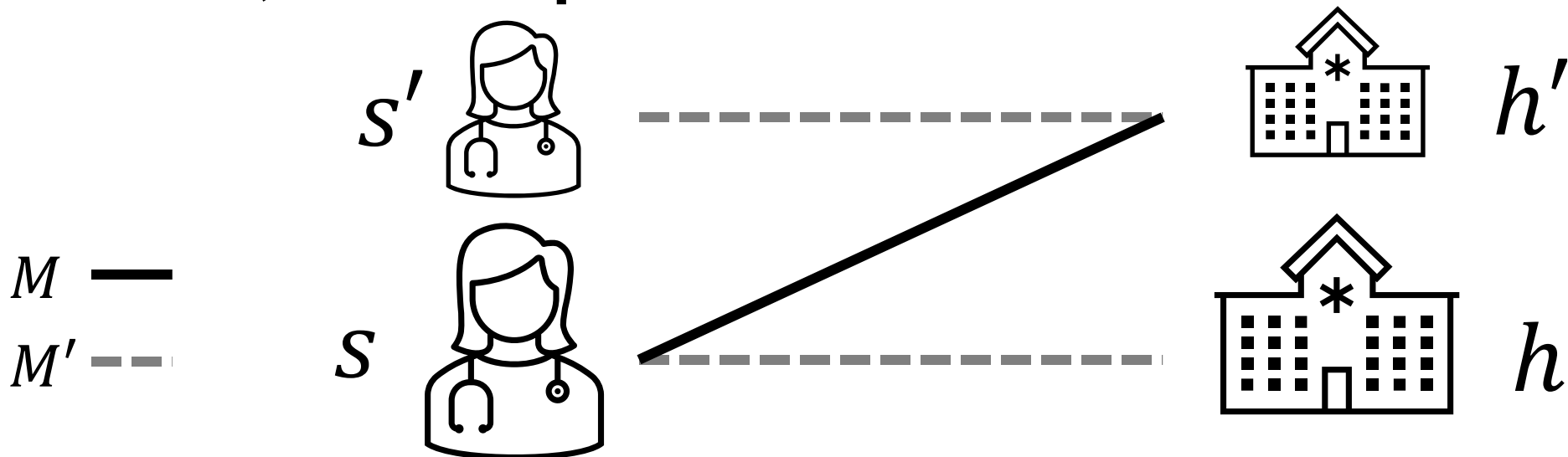
- Since h was rejected by s , it means that s **prefers h' over h** .



Proof Idea Time!

Proof: ...

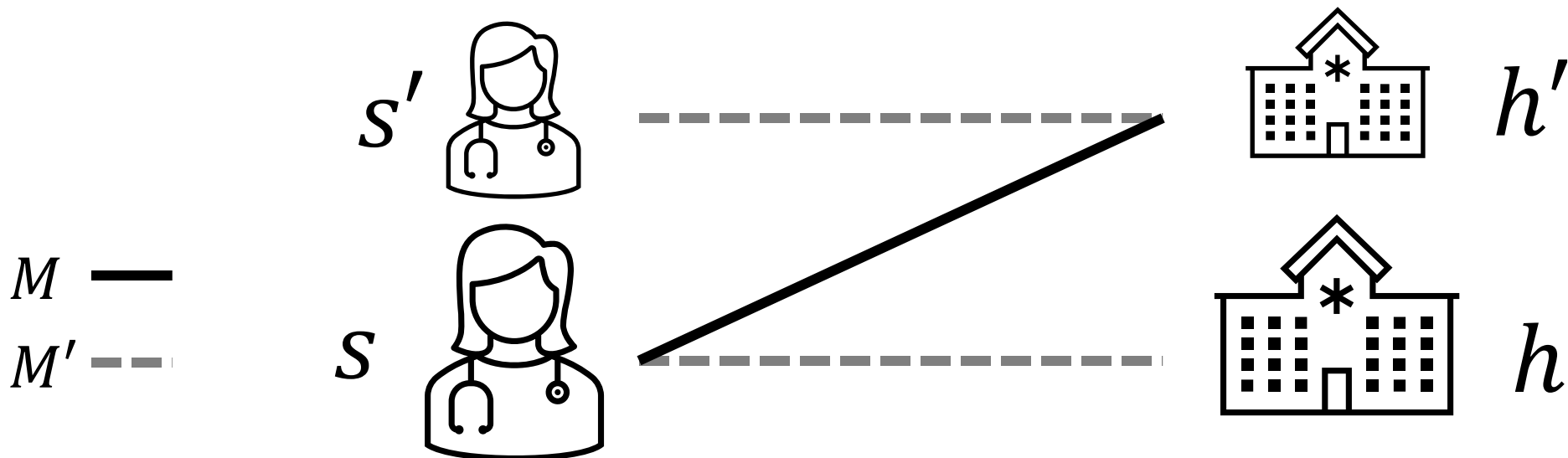
- Since h was rejected by s , it means that s prefers h' over h .
- Since h was the first rejected, h' did not get rejected and thus, **h' must prefer s to s'** .



Proof Idea Time!

Proof: ...

- If s prefers h' over h and h' must prefer s to s' then (h', s) is an unstable pair in M' . This is a contradiction!



Claim Time!

Claim: The algorithm always outputs the *student-pessimal* assignment.

WHILE (some hospital h is unmatched and hasn't proposed to every student)

$s \leftarrow$ first student on h 's list to whom h has not yet proposed.

IF (s is unmatched)

Add $h-s$ to matching M .

ELSE IF (s prefers h to current partner h')

Replace $h'-s$ with $h-s$ in matching M .

ELSE

s rejects h .

Claim Time!

Each student gets their worst valid partner.

Claim: The algorithm always outputs the *student-pessimal* assignment.

WHILE (some hospital h is unmatched and hasn't proposed to every student)

$s \leftarrow$ first student on h 's list to whom h has not yet proposed.

IF (s is unmatched)

Add $h-s$ to matching M .

ELSE IF (s prefers h to current partner h')

Replace $h'-s$ with $h-s$ in matching M .

ELSE

s rejects h .

Proof Idea Time!

Claim: The algorithm always outputs the *student-pessimal* assignment.

Proof:

- Suppose to the contrary that the output M is not **student-pessimal**.
 - Then there exists student s that is not matched with their worst valid partner h in M .

Proof Idea Time!

Claim: The algorithm always outputs the *student-pessimal* assignment.

Proof:

- Suppose to the contrary that the output M is not **student-pessimal**.
 - Then there exists student s that is not matched with their worst valid partner h in M .
 - Let M' be a stable matching such that s and h are matched.

Proof Idea Time!

Proof:

- Suppose to the contrary that the output M is not **student-pessimal**.
 - Then there exists student s that is not matched with their worst valid partner h in M .
 - Let M' be a stable matching such that s and h are matched.
- Let h' be the hospital s is matched to in M .

Proof Idea Time!

Proof:

- Suppose to the contrary that the output M is not **student-pessimal**.
 - Then there exists student s that is not matched with their worst valid partner h in M .
 - Let M' be a stable matching such that s and h are matched.
- Let h' be the hospital s is matched to in M .
 - Then s prefers h' to h .

Proof Idea Time!

Proof:

- Let h' be the hospital s is matched to in M .
 - Then s prefers h' to h otherwise h would not be worst valid pair as assumed.
- Let s' be the hospital h is matched to in M .

Proof Idea Time!

Proof:

- Let h' be the hospital s is matched to in M .
 - Then s prefers h' to h otherwise h would not be worst valid pair as assumed.
- Let s' be the hospital h is matched to in M .
 - Since M is hospital-optimal, h prefers s to s' .

Proof Idea Time!

Proof:

- Let h' be the hospital s is matched to in M .
 - Then s prefers h' to h otherwise h would not be worst valid pair as assumed.
- Let s' be the hospital h is matched to in M .
 - Since M is hospital-optimal, h prefers s to s' .
- Thus, (h, s) is an unstable pair in M . <- **Contradiction!**

Q: Does it Help to Lie?

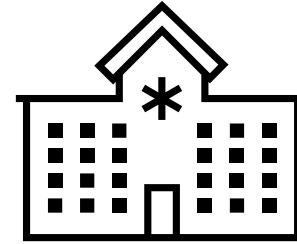
Dr. One
 $C > A > B$



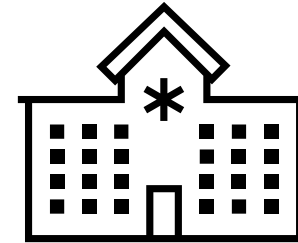
Dr. Two
 $A > C > B$



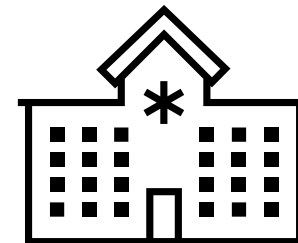
Dr. Three
 $A > B > C$



Hospital A
 $1 > 2 > 3$

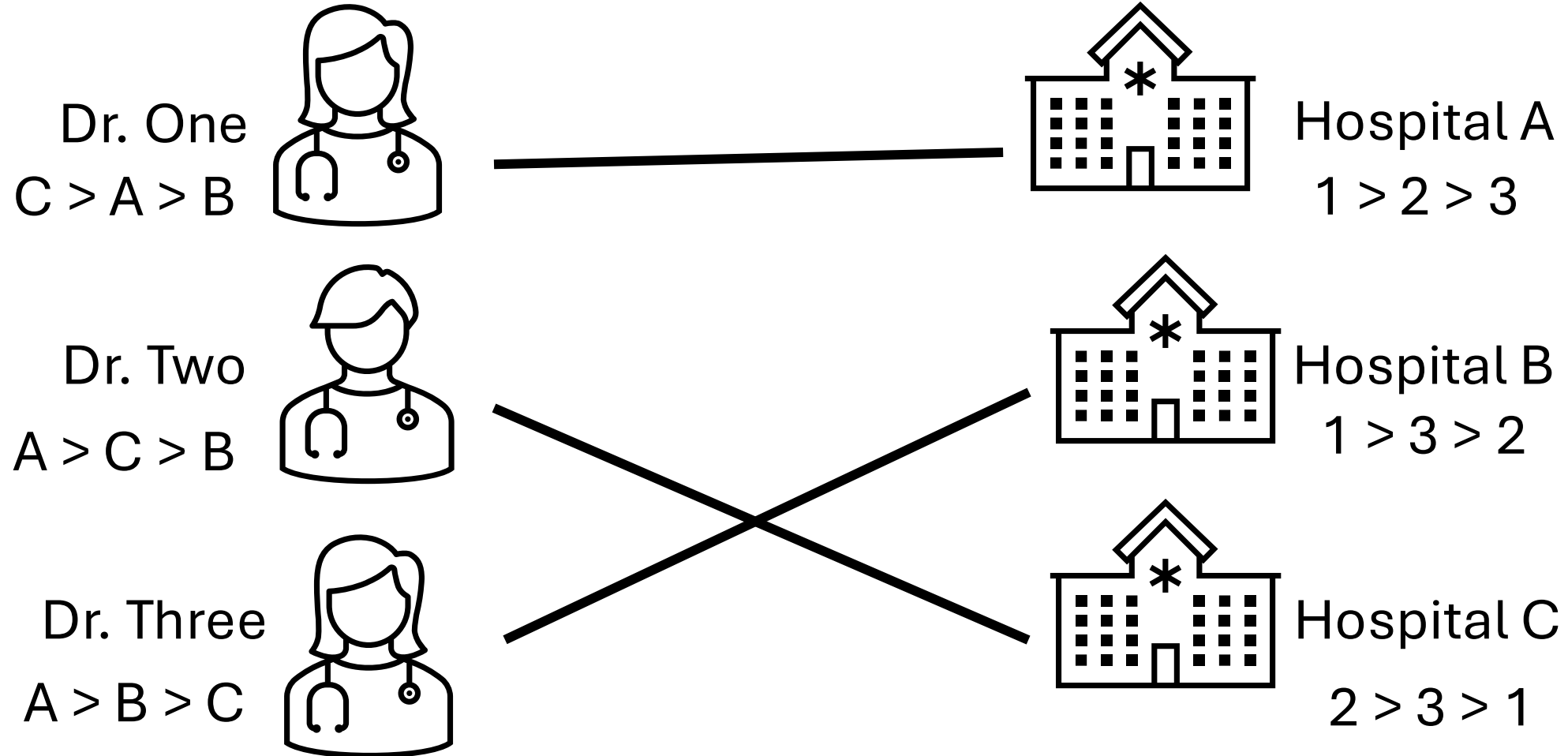


Hospital B
 $1 > 3 > 2$



Hospital C
 $2 > 3 > 1$

Q: Does it Help to Lie?



Q: Does it Help to Lie?

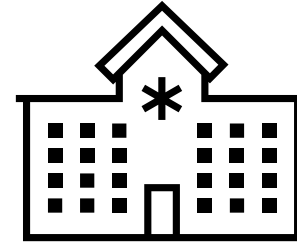
Dr. One
 $C > \textbf{B} > \textbf{A}$



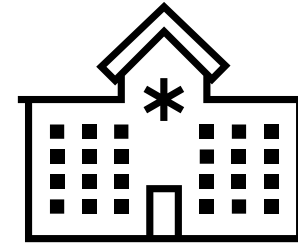
Dr. Two
 $A > C > B$



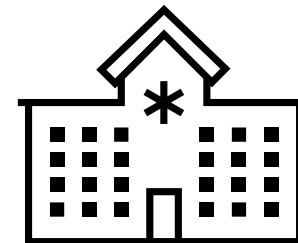
Dr. Three
 $A > B > C$



Hospital A
 $1 > 2 > 3$



Hospital B
 $1 > 3 > 2$



Hospital C
 $2 > 3 > 1$

Q: Does it Help to Lie?

