

CSE 443  
Compilers

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# Announcements

- Project document posted on website
  - Only worry about page 1 - 5 for now
- Team formation
  - We can take a few minutes to do it now
  - Make private Piazza post with UBIT username and corresponding GitHub username.



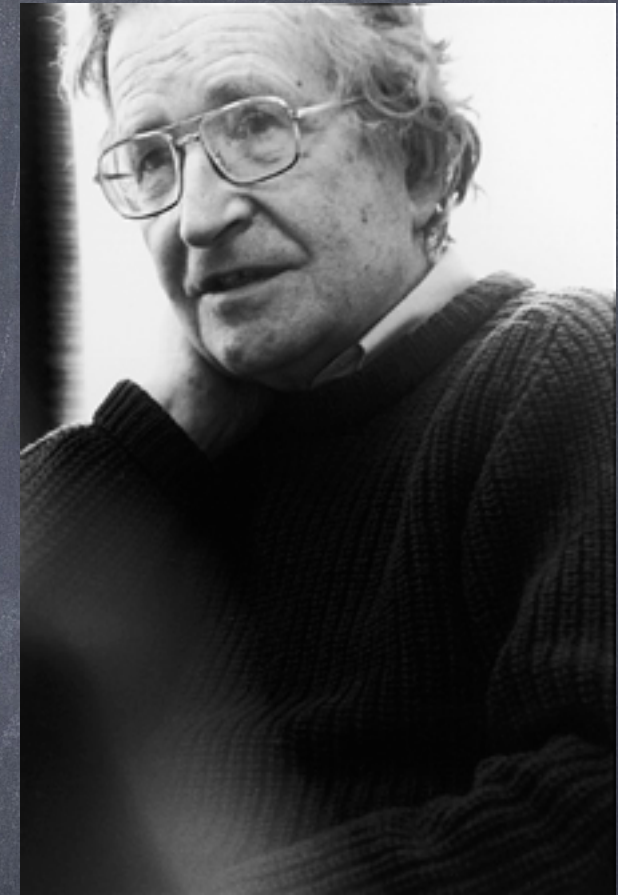
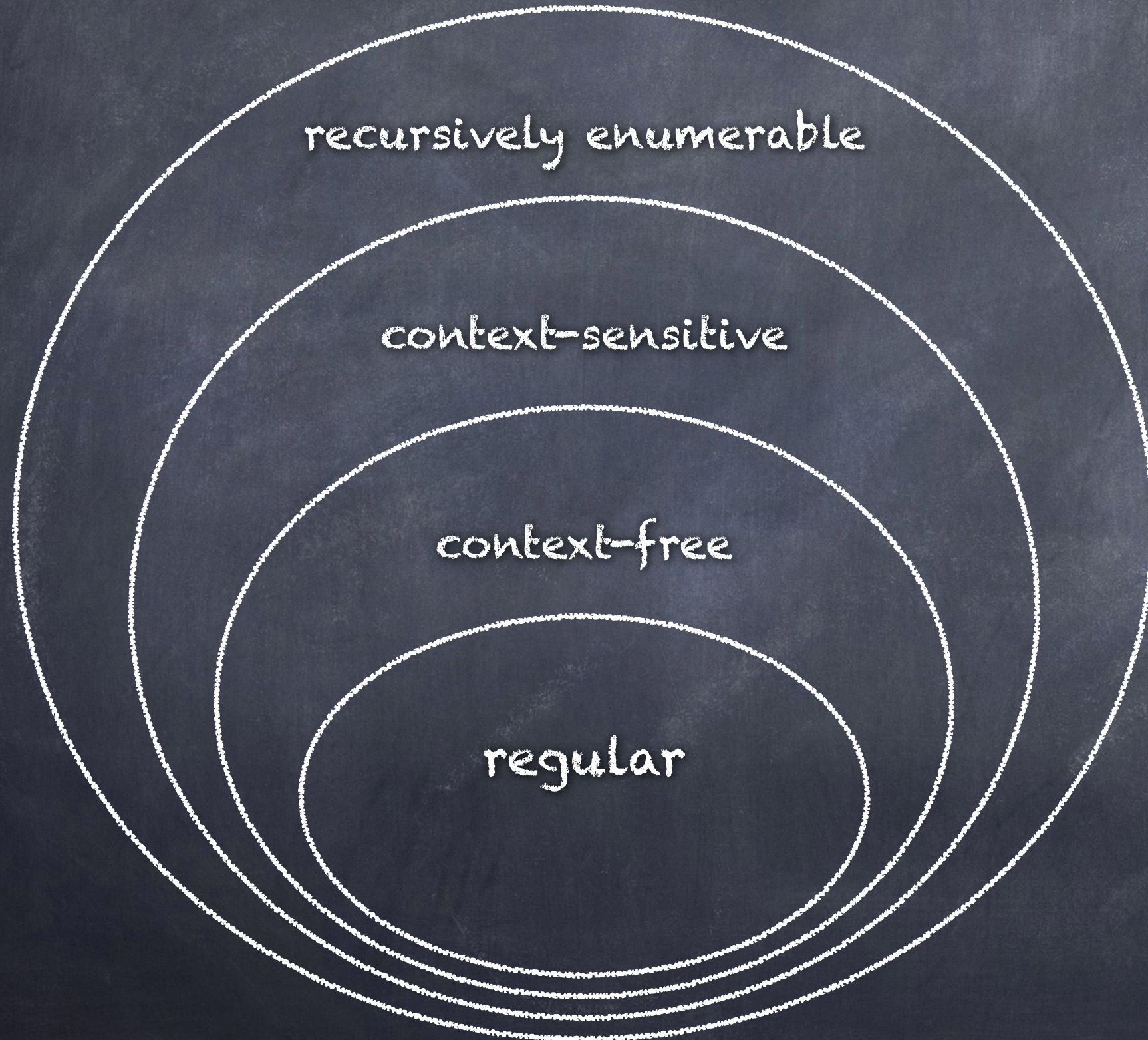
# Syntax and Semantics

- Syntax: program structure
- Semantics: program meaning
- Semantics are determined (in part) by program structure.

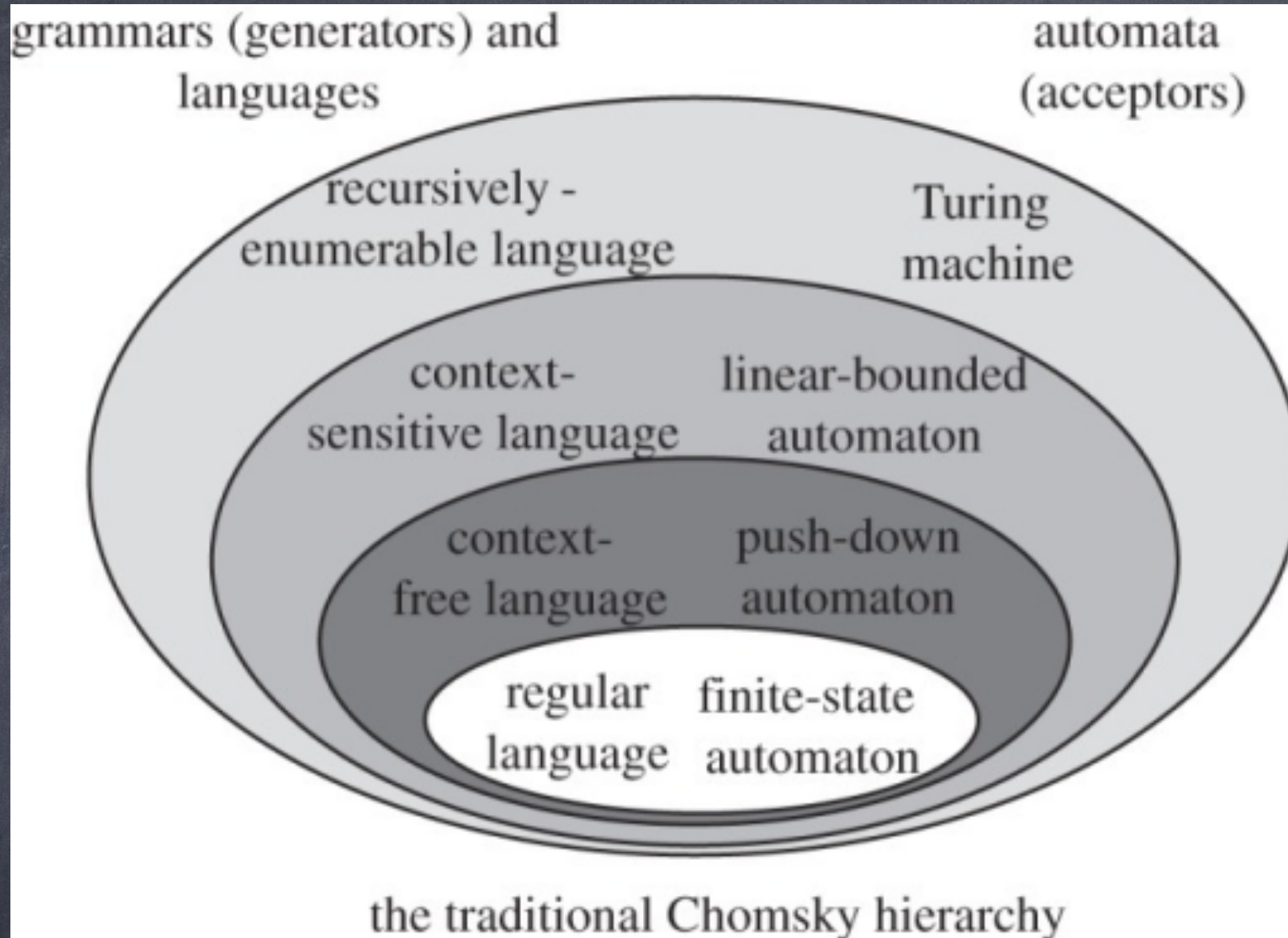


# Languages: the Chomsky hierarchy

"On Certain Formal Properties of Grammars" published 1959





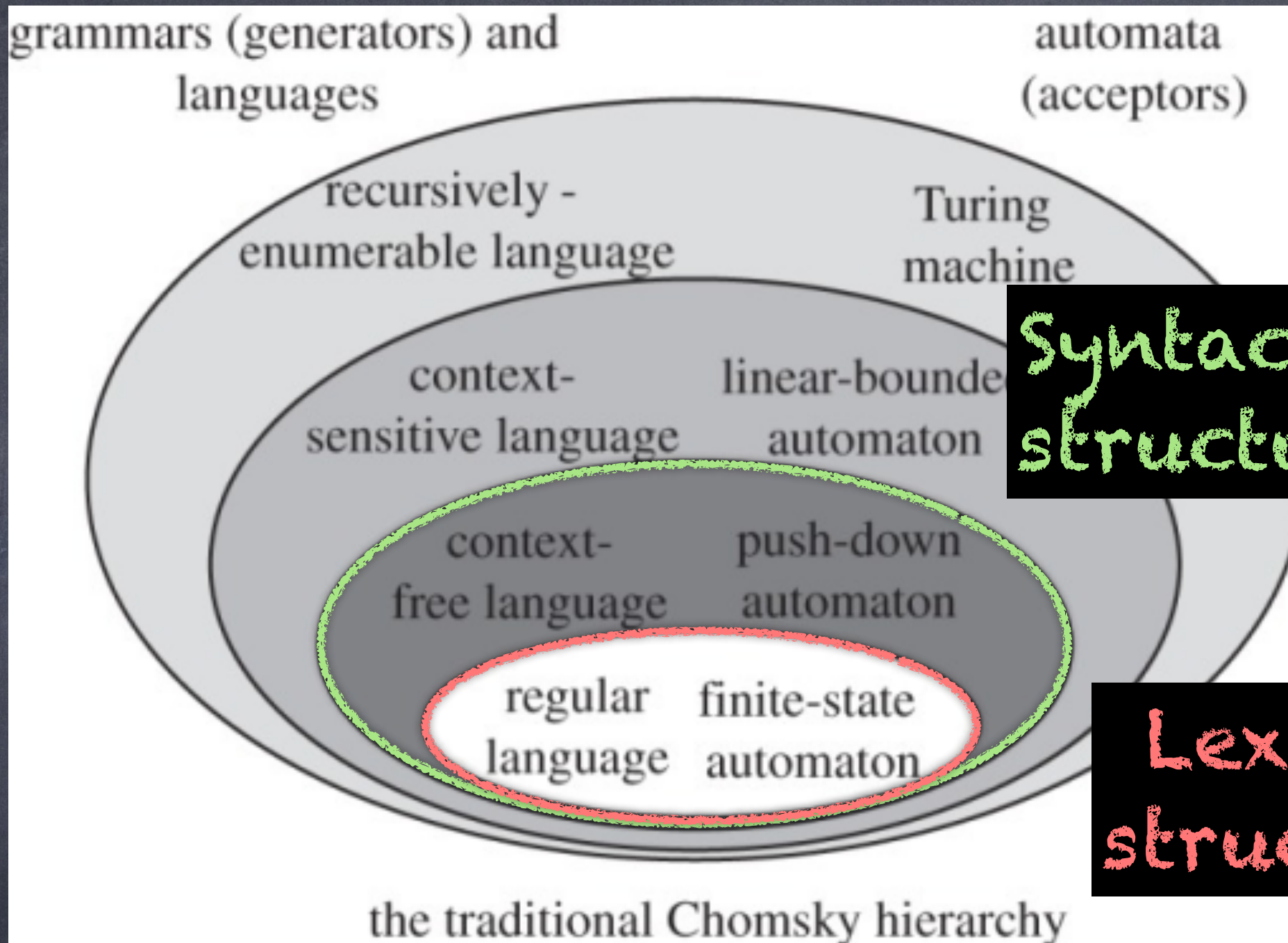


SOURCE: [https://openi.nlm.nih.gov/detailedresult.php?img=PMC3367694\\_rstb20120103-g2&req=4](https://openi.nlm.nih.gov/detailedresult.php?img=PMC3367694_rstb20120103-g2&req=4)

AUTHORS: Fitch WT, Friederici AD - Philos. Trans. R. Soc. Lond., B, Biol. Sci. (2012)

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# Phases of a compiler

Lexical  
structure

Syntactic  
structure

Symbol Table

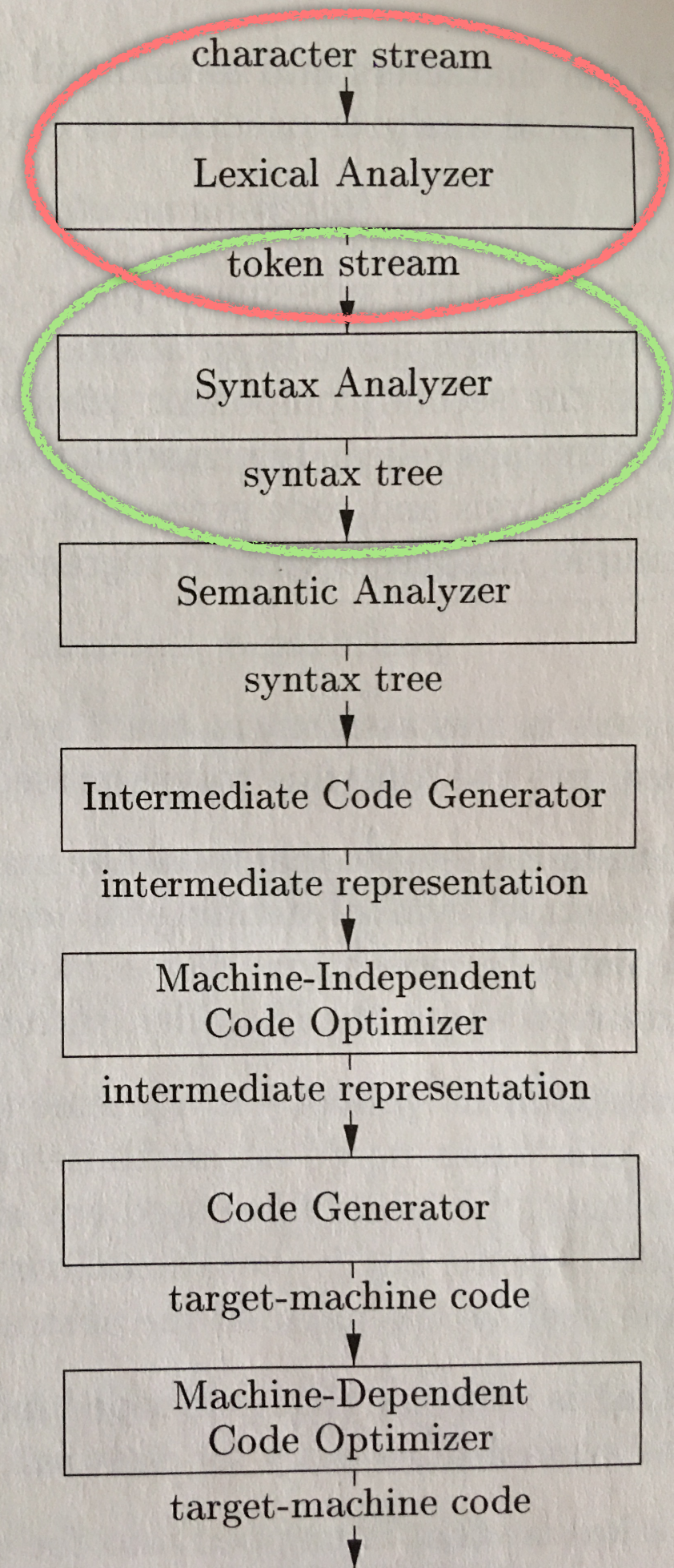


Figure 1.6,  
page 5 of text



# Phases of a compiler

## Lexical structure

Symbol Table

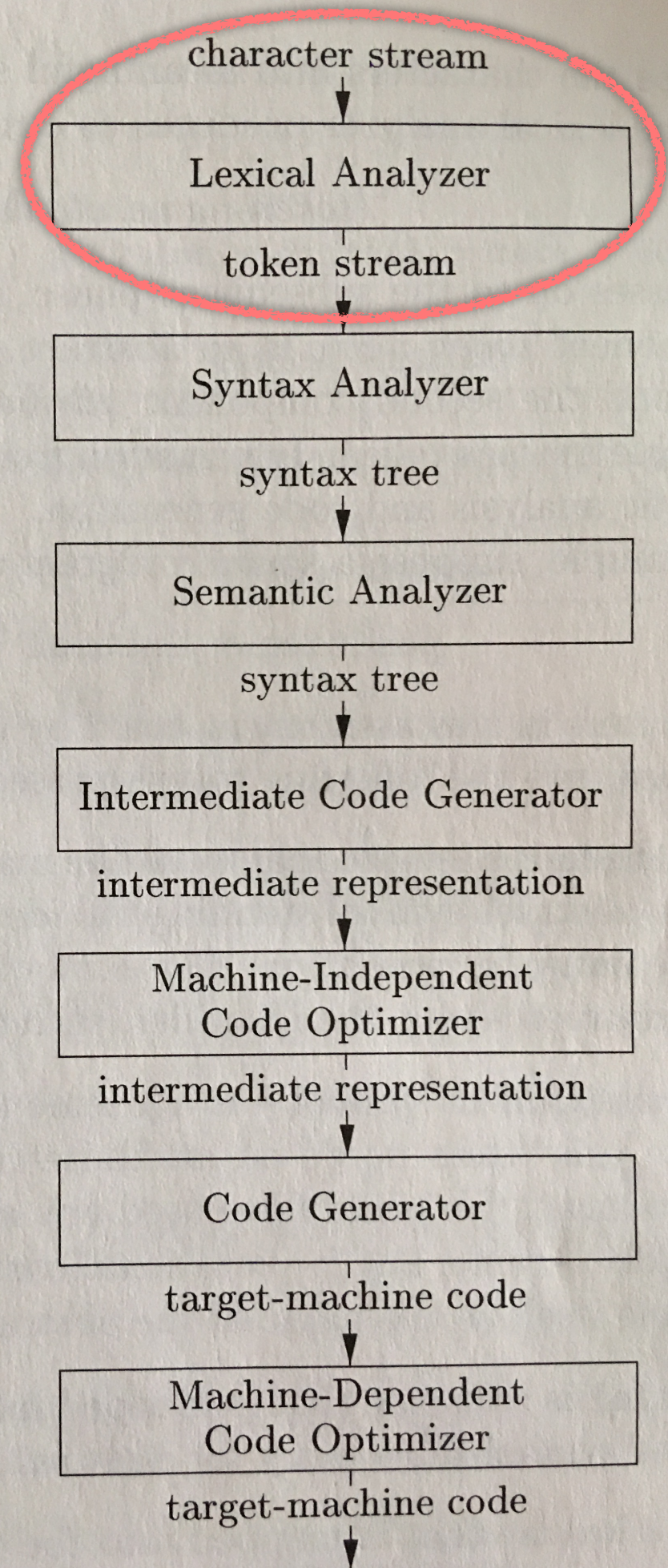


Figure 1.6,  
page 5 of text



# Lexical Structure

```
int main(){
```



# Lexical Structure

```
int main(){
```

character stream

```
int main() {
```



# Lexical Structure

```
int main(){
```

character stream → token stream

```
int main() {      id("int") id("main") LPAR RPAR LBRACE
```



# Lexical Structure

## tokens

- keywords (e.g. static, for, while, struct)
- operators (e.g. <, >, <=, =, ==, +, -, &, .)
- identifiers (e.g. foo, bar, sum, mystery)
- literals (e.g. -17, 34.52E-45, true, 'e', "Serenity")
- punctuation (e.g. {, }, (, ), ;)



# Describing lexical structure

- We need some formal way of describing the lexical structure of a language.



# meta vs object Language

- object language: the language we are describing
- meta language: the language we use to describe the object language



# meta vs object language

- How do we distinguish between the two?



# meta vs object Language

- use quotes (meta vs 'object')
- punctuation (e.g. '{', '}', '"', '\'', '`')
- use font or font property (meta vs **object**)
- punctuation (e.g. **{**, **}**, **(**, **)**, **;**)



# Languages & grammars

- Formally, a **language** is a set of strings over some alphabet
- Ex.  $\{00, 01, 10, 11\}$  is the set of all strings of length 2 over the alphabet  $\{0, 1\}$
- Ex.  $\{00, 11\}$  is the set of all even parity strings of length 2 over the alphabet  $\{0, 1\}$



# Languages & grammars

Formally, a grammar is defined by 4 items:

1.  $N$ , a set of non-terminals

2.  $\Sigma$ , a set of terminals

3.  $P$ , a set of productions

4.  $S$ , a start symbol

$G = (N, \Sigma, P, S)$



# Lexical analysis: a bird's eye view

{ for, while, x, factorial, ... }

language: a set of strings

finite automaton

a machine for language

C program

generated by FLEX

$G = (N, \Sigma, P, S)$

grammar: rules for generating language

regular expression

regex: a form of grammar



# Languages & grammars

Formally, a grammar  $G = (N, \Sigma, P, S)$  is defined by 4 items:

1.  $N$ , a set of non-terminals

$$N = \{ X, Y \}$$

2.  $\Sigma$ , a set of terminals (alphabet)

$$\Sigma = \{ a, b \} \quad \leftarrow \text{for example}$$

$$N \cap \Sigma = \{ \} \quad \leftarrow \text{general grammar constraints}$$

3.  $P$ , a set of productions of the form (right linear)

$$X \rightarrow aY$$

$$Y \rightarrow bX$$

$$Y \rightarrow a \quad \leftarrow \text{a right linear grammar describing a regular language}$$

$$X \rightarrow \varepsilon$$

$$X \in N, Y \in N, a \in \Sigma, b \in \Sigma, \varepsilon \text{ denotes the empty string}$$

4.  $S$ , a start symbol

$$S = Y$$

$$S \in N$$



# Languages & grammars

[...] a regular grammar is a grammar that is right-regular or left-regular:

- all production rules have at most one non-terminal symbol
- that symbol is either always at the end or always at the start of the rule's right hand side.

[https://en.wikipedia.org/wiki/Regular\\_grammar](https://en.wikipedia.org/wiki/Regular_grammar)



# Languages & grammars

Given a string  $\alpha A$ , where  $\alpha \in \Sigma^*$  and  $A \in N$ ,  
and a production

$$A \rightarrow \beta \in P$$

we write  $\alpha A \Rightarrow \alpha\beta$  to indicate that  $\alpha A$  derives  
 $\alpha\beta$  in one step.

$\Rightarrow^k$  and  $\Rightarrow^*$  can be used to indicate  $k$  or  
arbitrarily many derivation steps, respectively.