CSE443
Compilers

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Phases of a compiler

Syntactic structure

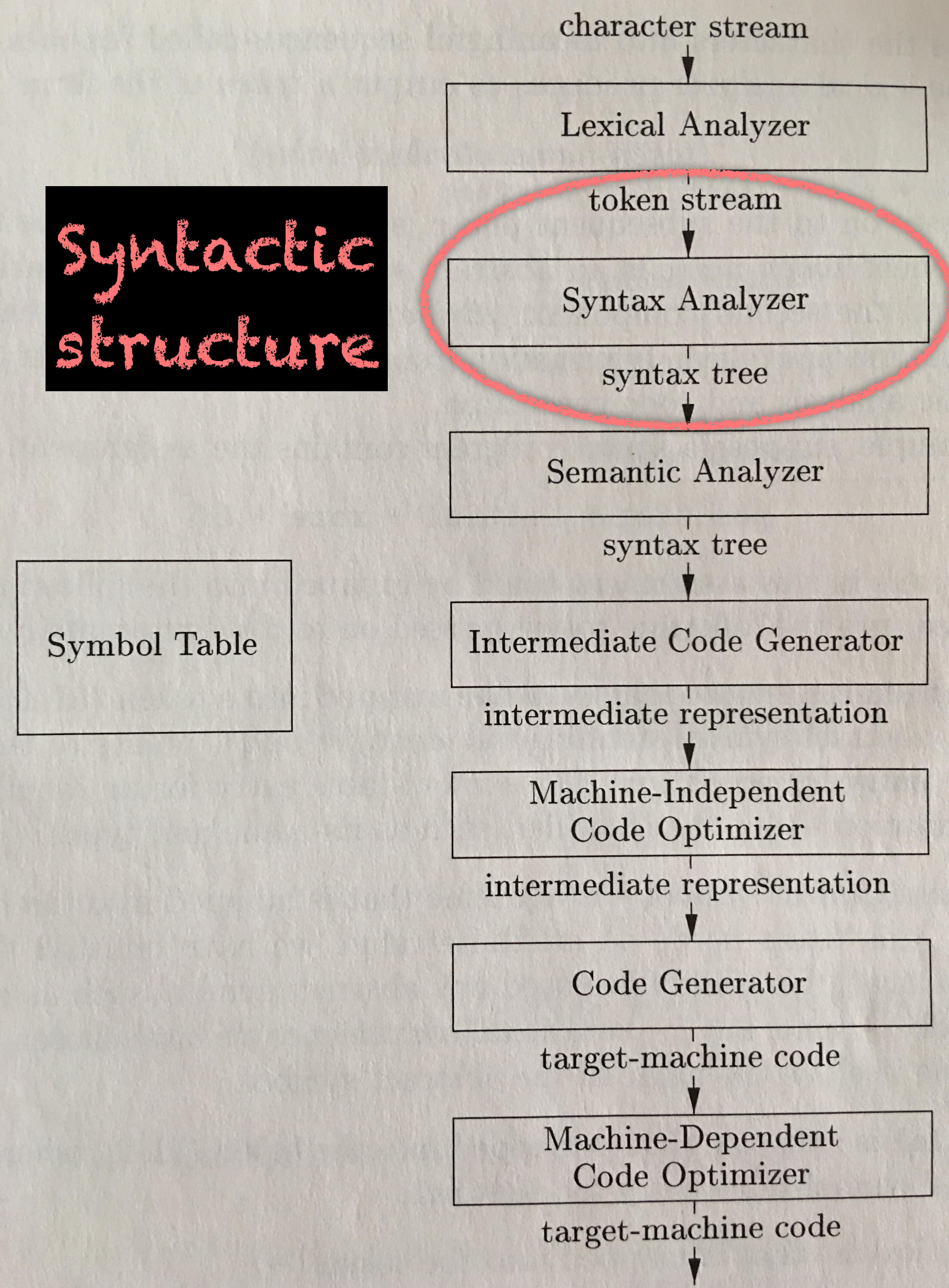

Figure 1.6, page 5 of text

FIRST(X)

- if $X \in T$ then FIRST(X) = { X }
- if $X \in N$ and $X \rightarrow Y_1\ Y_2\ \ldots\ Y_k \in P$ for $k \geq 1$, then
  - add $a \in T$ to FIRST(X) if $\exists i$ s.t. $a \in$ FIRST(Y_i) and $\varepsilon \in$ FIRST(Y_j) $\forall\ j < i$ (i.e. $Y_1\ Y_2\ \ldots\ Y_{i-1} \Rightarrow^* \varepsilon$)
  - if $\varepsilon \in$ FIRST(Y_j) $\forall\ j <= k$ add ε to FIRST(X)
  - if $X \rightarrow \varepsilon \in P$, then add ε to FIRST(X)

FIRST SETS

- FIRST(F) = ?
- FIRST(T) = ?
- FIRST(E) = ?
- FIRST(E') = ?
- FIRST(T') = ?

E -> T E'
E' -> + T E' | ε
T -> F T'
T' -> * F T' | ε
F -> (E) | id

- if $X \in T$ then $FIRST(X) = \{ X \}$
- if $X \in N$ and $X \rightarrow Y_1 Y_2 \ldots Y_k \in P$ for $k \geq 1$, then
  - add $a \in T$ to FIRST(X) if $\exists i$ s.t. $a \in FIRST(Y_i)$ and $\varepsilon \in FIRST(Y_j)$ $\forall j < i$ (i.e. $Y_1 Y_2 \ldots Y_j \Rightarrow^* \varepsilon$)
  - if $\varepsilon \in FIRST(Y_j)$ $\forall j < k$ add ε to FIRST(X)

FIRST SETS

E -> T E'
E' -> + T E' | ε
T -> F T'
T' -> * F T' | ε
F -> (E) | id

- FIRST(F) = { (, id }
- FIRST(T) = FIRST(F) = { (, id }
- FIRST(E) = FIRST(T) = FIRST(F) = { (, id }
- FIRST(E') = { + , ε }
- FIRST(T') = { * , ε }

- if $X \in T$ then FIRST(X) = { X }
- if $X \in N$ and $X \rightarrow Y_1 Y_2 \dots Y_k \in P$ for $k \geq 1$, then
  - add $a \in T$ to FIRST(X) if $\exists i$ s.t. $a \in$ FIRST(Y_i) and $\varepsilon \in$ FIRST(Y_j) $\forall j < i$ (i.e. $Y_1 Y_2 \dots Y_j \Rightarrow^* \varepsilon$)
  - if $\varepsilon \in$ FIRST(Y_j) $\forall j < k$ add ε to FIRST(X)

FOLLOW(X)

- Place \$ in FOLLOW(S), where S is the start symbol (\$ is an end marker)
- if $A \rightarrow \alpha B \beta \in P$, then $\text{FIRST}(\beta) - \{\varepsilon\}$ is in FOLLOW(B)
- if $A \rightarrow \alpha B \in P$ or $A \rightarrow \alpha B \beta \in P$ where $\varepsilon \in \text{FIRST}(\beta)$, then everything in FOLLOW(A) is in FOLLOW(B)

FOLLOW SETS

E -> T E'
E' -> + T E' | ε
T -> F T'
T' -> * F T' | ε
F -> (E) | id

- FOLLOW(E) = ?
- FOLLOW(E') = ?
- FOLLOW(T) = ?
- FOLLOW(T') = ?
- FOLLOW(F) = ?

- Place $ in FOLLOW(S), where S is the start symbol ($ is an end marker)
- if $A \rightarrow \alpha B \beta \in P$, then FIRST($\beta$) - {ε} is in FOLLOW(B)
- if $A \rightarrow \alpha B \in P$ or $A \rightarrow \alpha B \beta \in P$ where $\varepsilon \in$ FIRST(β), then everything in FOLLOW(A) is in FOLLOW(B)

FOLLOW SETS

$E \rightarrow T\ E'$
$E' \rightarrow +\ T\ E' \mid \varepsilon$
$T \rightarrow F\ T'$
$T' \rightarrow *\ F\ T' \mid \varepsilon$
$F \rightarrow (\ E\ ) \mid id$

- FOLLOW(E) = {) , $ }
- FOLLOW(E') = FOLLOW(E) = {) , $ }
- FOLLOW(T) = { + ,) , $ }
- FOLLOW(T') = FOLLOW(T) = { + ,) , $ }
- FOLLOW(F) = { + , * ,) , $ }

- Place $ in FOLLOW(S), where S is the start symbol ($ is an end marker)
- if $A \rightarrow \alpha B \beta \in P$, then $FIRST(\beta) - \{\varepsilon\}$ is in FOLLOW(B)
- if $A \rightarrow \alpha B \in P$ or $A \rightarrow \alpha B \beta \in P$ where $\varepsilon \in FIRST(\beta)$, then everything in FOLLOW(A) is in FOLLOW(B)

Table-driven predictive parsing Algorithm 4.32 (p. 224)

- INPUT: Grammar G = (N,T,P,S)
- OUTPUT: Parsing table M
- For each production A -> α of G:

1. For each terminal a $\in$ FIRST(α), add A -> α to M[A,a]
2. If $\varepsilon \in$ FIRST(α), then for each terminal b in FOLLOW(A), add A -> α to M[A,b]
3. If $\varepsilon \in$ FIRST(α) and \$ $\in$ FOLLOW(A), add A -> α to M[A,\$]

FIRST SETS

E -> T E'
E' -> + T E' | ε
T -> F T'
T' -> * F T' | ε
F -> (E) | id

- FIRST(F) = { (, id }
- FIRST(T) = FIRST(F) = { (, id }
- FIRST(E) = FIRST(T) = FIRST(F) = { (, id }
- FIRST(E') = { + , ε }
- FIRST(T') = { * , ε }

- if $X \in T$ then FIRST(X) = { X }
- if $X \in N$ and $X \rightarrow Y_1 Y_2 \ldots Y_k \in P$ for $k \geq 1$, then
  - add $a \in T$ to FIRST(X) if $\exists i$ s.t. $a \in$ FIRST(Y_i) and $\varepsilon \in$ FIRST(Y_j) $\forall j < i$ (i.e. $Y_1 Y_2 \ldots Y_j \Rightarrow^* \varepsilon$)
  - if $\varepsilon \in$ FIRST(Y_j) $\forall j < k$ add ε to FIRST(X)

For each production A -> α of G:

- For each terminal $a \in$ FIRST(α), add A -> α to M[A,a]
- If $\varepsilon \in$ FIRST(α), then for each terminal b in FOLLOW(A), add A -> α to M[A,b]
- If $\varepsilon \in$ FIRST(α) and $\$ \in$ FOLLOW(A), add A -> α to M[A,$]

FOLLOW SETS

E -> T E'
E' -> + T E' | ε
T -> F T'
T' -> * F T' | ε
F -> (E) | id

- FOLLOW(E) = {) , $ }
- FOLLOW(E') = FOLLOW(E) = {) , $ }
- FOLLOW(T) = { + ,) , $ }
- FOLLOW(T') = FOLLOW(T) = { + ,) , $ }
- FOLLOW(F) = { + , * ,) , $ }

- Place $ in FOLLOW(S), where S is the start symbol ($ is an end marker)
- if A -> $\alpha B \beta \in P$, then FIRST(β) – {ε} is in FOLLOW(B)
- if A -> $\alpha B \in P$ or A -> $\alpha B \beta \in P$ where $\varepsilon \in$ FIRST(β), then everything in FOLLOW(A) is in FOLLOW(B)

For each production A -> α of G:

- For each terminal a $\in$ FIRST(α), add A -> α to M[A,a]
- If $\varepsilon \in$ FIRST(α), then for each terminal b in FOLLOW(A), add A -> α to M[A,b]
- If $\varepsilon \in$ FIRST(α) and $ $\in$ FOLLOW(A), add A -> α to M[A,$]

EXERCISE:
fill in the parse table

(see next slide)

Parse-table M

NON TERMINALS	id	+	*	(	)	$
E						
E'						
T						
T'						
F						

FIRST(E) = FIRST(T) = FIRST(F) = { (, id }

FIRST(E') = { + , ε }

FIRST(T') = { * , ε }

FOLLOW(E') = FOLLOW(E) = {) , $ }

FOLLOW(T') = FOLLOW(T) = { + ,) , $ }

FOLLOW(F) = { + , * ,) , $ }

E -> T E'

E' -> + T E' | ε

T -> F T'

T' -> * F T' | ε

F -> (E) | id

For each production A -> α of G:

- For each terminal a ∈ FIRST(α), add A -> α to M[A,a]
- If ε ∈ FIRST(α), then for each terminal b in FOLLOW(A), add A -> α to M[A,b]
- If ε ∈ FIRST(α) and $ ∈ FOLLOW(A), add A -> α to M[A,$]

Parse-table M

NON TERMINALS	id	+	*	(	)	$
E	E -> T E'			E -> T E'		
E'						
T						
T'						
F						

FIRST(E) = FIRST(T) = FIRST(F) = { (, id }

FIRST(E') = { + , ε }

FIRST(T') = { * , ε }

FOLLOW(E') = FOLLOW(E) = {) , $ }

FOLLOW(T') = FOLLOW(T) = { + ,) , $ }

FOLLOW(F) = { + , * ,) , $ }

E -> T E'

E' -> + T E' | ε

T -> F T'

T' -> * F T' | ε

F -> (E) | id

For each production A -> α of G:

- For each terminal a ∈ FIRST(α), add A -> α to M[A,a]
- If ε ∈ FIRST(α), then for each terminal b in FOLLOW(A), add A -> α to M[A,b]
- If ε ∈ FIRST(α) and $ ∈ FOLLOW(A), add A -> α to M[A,$]

Parse-table M

NON TERMINALS	id	+	*	(	)	$
E	E -> T E'			E -> T E'		
E'		E' -> + T E'				
T						
T'						
F						

FIRST(E) = FIRST(T) = FIRST(F) = { (, id }

FIRST(E') = { + , ε }

FIRST(T') = { * , ε }

FOLLOW(E') = FOLLOW(E) = {) , $ }

FOLLOW(T') = FOLLOW(T) = { + ,) , $ }

FOLLOW(F) = { + , * ,) , $ }

E -> T E'

E' -> + T E' | ε

T -> F T'

T' -> * F T' | ε

F -> (E) | id

For each production A -> α of G:

- For each terminal a ∈ FIRST(α), add A -> α to M[A,a]
- If ε ∈ FIRST(α), then for each terminal b in FOLLOW(A), add A -> α to M[A,b]
- If ε ∈ FIRST(α) and $ ∈ FOLLOW(A), add A -> α to M[A,$]

Parse-table M

NON TERMINALS	id	+	*	(	)	$
E	E -> T E'			E -> T E'		
E'		E' -> + T E'			E' -> ε	E' -> ε
T						
T'						
F						

FIRST(E) = FIRST(T) = FIRST(F) = { (, id }

FIRST(E') = { + , ε }

FIRST(T') = { * , ε }

FOLLOW(E') = FOLLOW(E) = {) , $ }

FOLLOW(T') = FOLLOW(T) = { + ,) , $ }

FOLLOW(F) = { + , * ,) , $ }

E -> T E'

E' -> + T E' | ε

T -> F T'

T' -> * F T' | ε

F -> (E) | id

For each production A -> α of G:

- For each terminal a ∈ FIRST(α), add A -> α to M[A,a]
- If ε ∈ FIRST(α), then for each terminal b in FOLLOW(A), add A -> α to M[A,b]
- If ε ∈ FIRST(α) and $ ∈ FOLLOW(A), add A -> α to M[A,$]

EXERCISE:

to be continued!