

CSE 443

Compilers

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Phases of a compiler

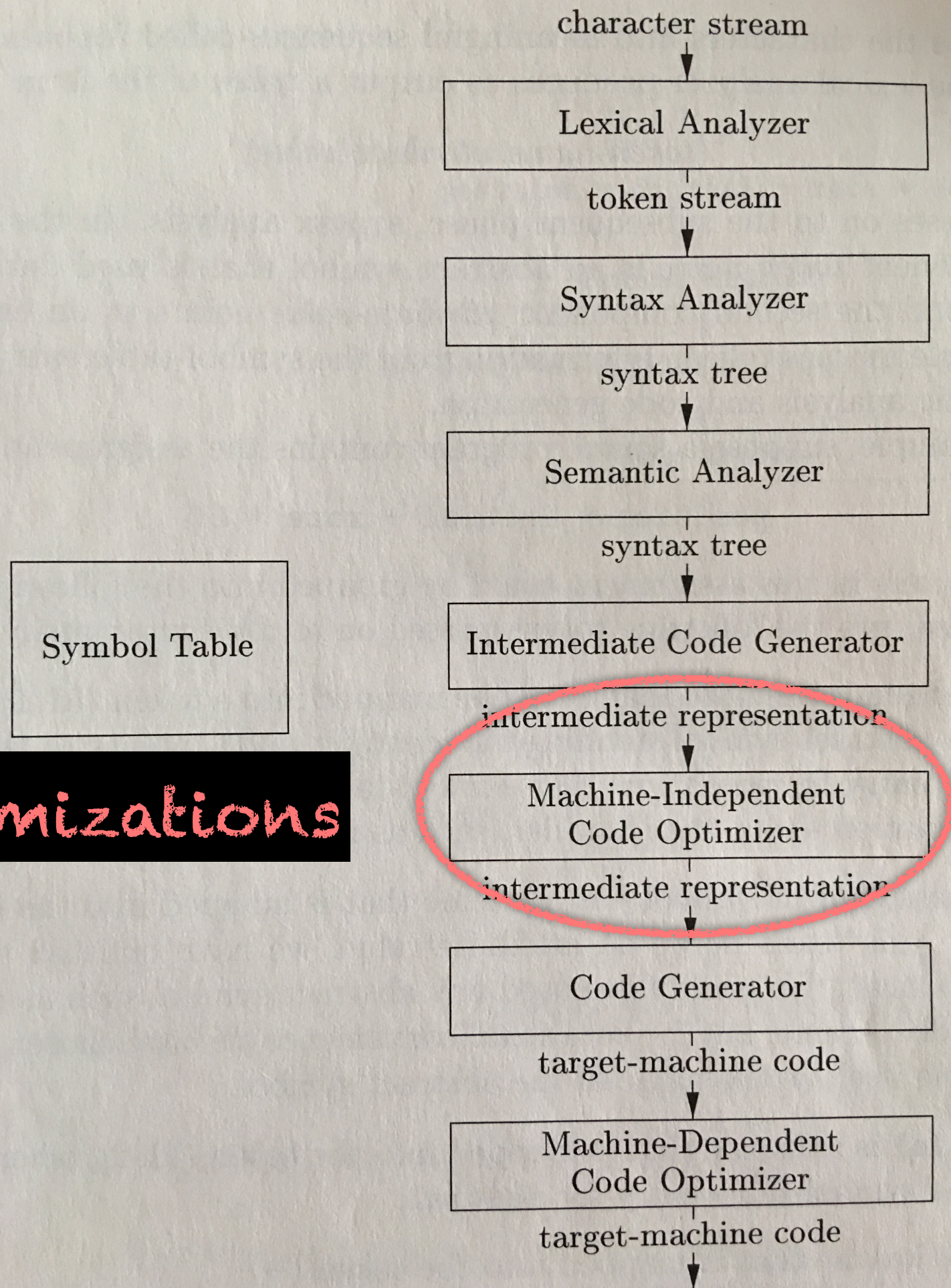


Figure 1.6,
page 5 of text

Machine-Independent Optimizations

Chapter 8: local optimization
(within a basic block)

Chapter 9: global optimization
(across basic blocks)

Optimization

- The semantics of a program **must** be preserved by optimizations.
- The compiler does not know a programmer's intent - it can only reason about the program as written.

Why is there redundancy?
Why is there redundancy?

- May be addressable in source code.

"a programmer may find it more direct and convenient to recalculate some result, leaving it the compiler to recognize that only one such calculation is necessary" [p. 584]

- Often cannot be resolved in HLL.

e.g. "...programmers have no choice but to refer to elements of an array or fields in a structure through accesses like $A[i][j]$..." [p.584]

PLAN

- Discuss code transformations in context of a small code example (Sedgewick's Quicksort implementation - see page 585)

Quicksort [p. 585]

```
void quicksort(int m, int n)
    /* recursively sorts a[m] through a[n] */
{
    int i, j;
    int v, x;
    if (n <= m) return;
    /* fragment begins here */
    i = m-1; j = n; v = a[n];
    while (1) {
        do i = i+1; while (a[i] < v);
        do j = j-1; while (a[j] > v);
        if (i >= j) break;
        x = a[i]; a[i] = a[j]; a[j] = x; /* swap a[i], a[j] */
    }
    /* fragment ends here */
    quicksort(m,j); quicksort(i+1,n);
}
```



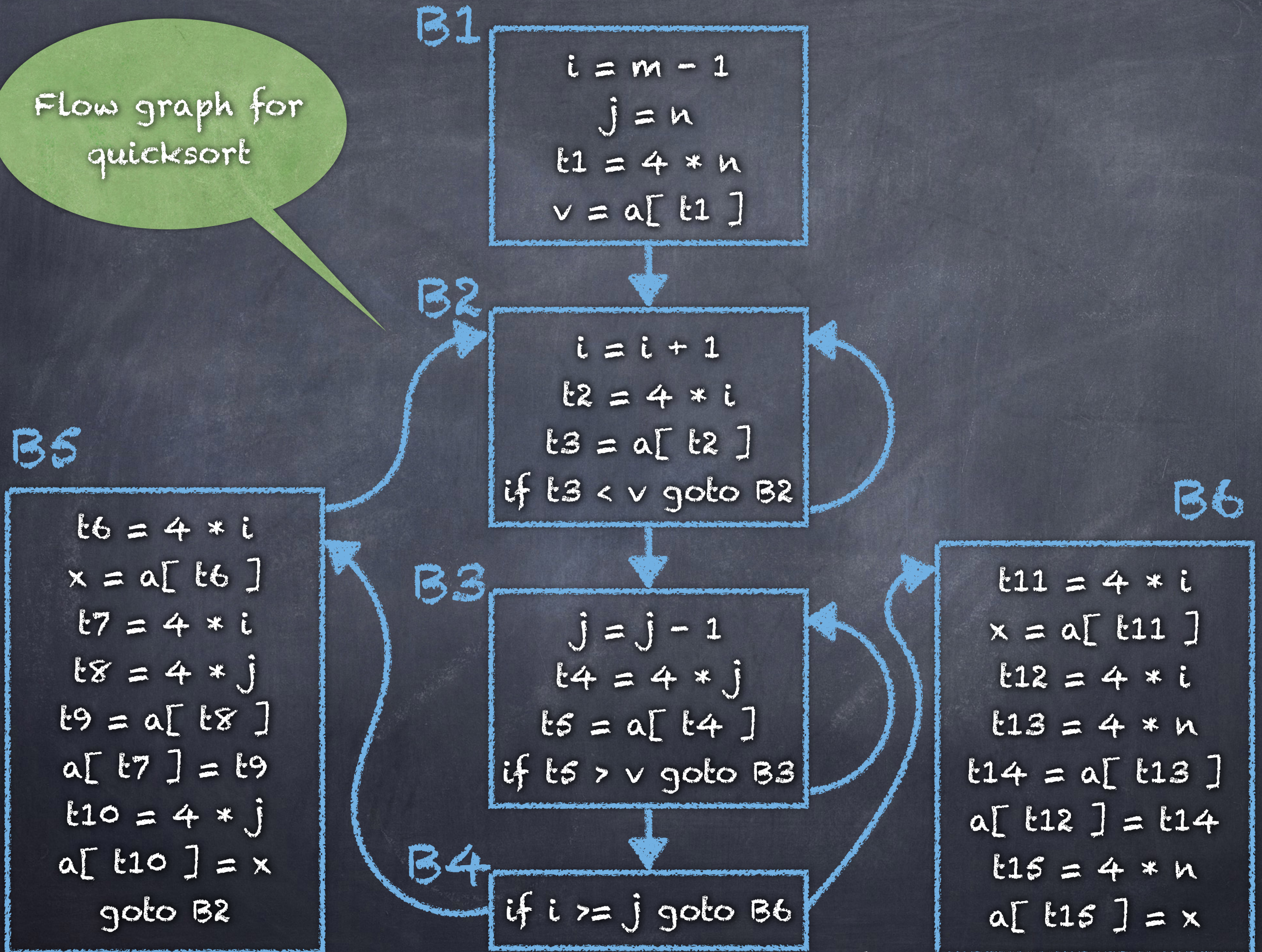
```
01) i = m - 1
02) j = n
03) t1 = 4 * n
04) v = a[ t1 ]
05) i = i + 1
06) t2 = 4 * i
07) t3 = a[ t2 ]
08) if t3 < v goto 5 (=B2)
09) j = j - 1
10) t4 = 4 * j
11) t5 = a[ t4 ]
12) if t5 > v goto 9 (=B3)
13) if i >= j goto 23 (=B6)
14) t6 = 4 * i
15) x = a[ t6 ]
16) t7 = 4 * i
17) t8 = 4 * j
18) t9 = a[ t8 ]
19) a[ t7 ] = t9
20) t10 = 4 * j
21) a[ t10 ] = x
22) goto 5 (=B2)
23) t11 = 4 * i
24) x = a[ t11 ]
25) t12 = 4 * i
26) t13 = 4 * n
27) t14 = a[ t13 ]
28) a[ t12 ] = t14
29) t15 = 4 * n
30) a[ t15 ] = x
```

IR code for
highlighted
fragment
of Quicksort

PLAN

- Original flow graph (9.3 - p. 587)
- Common Subexpression Elimination flow graph (9.5 - p. 589)
- Copy propagation flow graph (9.7 - p. 591)
- Dead code elimination (p. 591)
- Strength reduction flow graphs (9.8 - p. 593 and 9.9 - p. 594)

Flow graph for quicksort



B1

$i = m - 1$
 $j = n$
 $t1 = 4 * n$
 $v = a[t1]$

B2

$i = i + 1$
 $t2 = 4 * i$
 $t3 = a[t2]$
if $t3 < v$ goto B2

B3

$j = j - 1$
 $t4 = 4 * j$
 $t5 = a[t4]$
if $t5 > v$ goto B3

B4

if $i \geq j$ goto B6

B5

$t6 = 4 * i$
 $x = a[t6]$
 $t7 = 4 * i$
 $t8 = 4 * j$
 $t9 = a[t8]$
 $a[t7] = t9$
 $t10 = 4 * j$
 $a[t10] = x$
goto B2

B6

$t11 = 4 * i$
 $x = a[t11]$
 $t12 = 4 * i$
 $t13 = 4 * n$
 $t14 = a[t13]$
 $a[t12] = t14$
 $t15 = 4 * n$
 $a[t15] = x$

Local common subexpressions in B5 (similarly in B6)

B1

```
i = m - 1
j = n
t1 = 4 * n
v = a[t1]
```

B2

```
i = i + 1
t2 = 4 * i
t3 = a[t2]
if t3 < v goto B2
```

B3

```
j = j - 1
t4 = 4 * j
t5 = a[t4]
if t5 > v goto B3
```

B4

```
if i >= j goto B6
```

B5

```
t6 = 4 * i
x = a[t6]
t7 = 4 * i
t8 = 4 * j
t9 = a[t8]
a[t7] = t9
t10 = 4 * j
a[t10] = x
goto B2
```

B6

```
t11 = 4 * i
x = a[t11]
t12 = 4 * i
t13 = 4 * n
t14 = a[t13]
a[t12] = t14
t15 = 4 * n
a[t15] = x
```


Local common subexpressions in B5 (similarly in B6)

B1

```
i = m - 1
j = n
t1 = 4 * n
v = a[ t1 ]
```

B2

```
i = i + 1
t2 = 4 * i
t3 = a[ t2 ]
if t3 < v goto B2
```

B3

```
j = j - 1
t4 = 4 * j
t5 = a[ t4 ]
if t5 > v goto B3
```

B4

```
if i >= j goto B6
```

B5

```
t6 = 4 * i
x = a[ t6 ]

t8 = 4 * j
t9 = a[ t8 ]
a[ t6 ] = t9

a[ t8 ] = x
goto B2
```

B6

```
t11 = 4 * i
x = a[ t11 ]
t12 = 4 * i
t13 = 4 * n
t14 = a[ t13 ]
a[ t12 ] = t14
t15 = 4 * n
a[ t15 ] = x
```


Local common subexpressions in B5 (similarly in B6)

B1

```
i = m - 1
j = n
t1 = 4 * n
v = a[ t1 ]
```

Local common subexpressions in B6

B2

```
i = i + 1
t2 = 4 * i
t3 = a[ t2 ]
if t3 < v goto B2
```

B5

```
t6 = 4 * i
x = a[ t6 ]

t8 = 4 * j
t9 = a[ t8 ]
a[ t6 ] = t9

a[ t8 ] = x
goto B2
```

B6

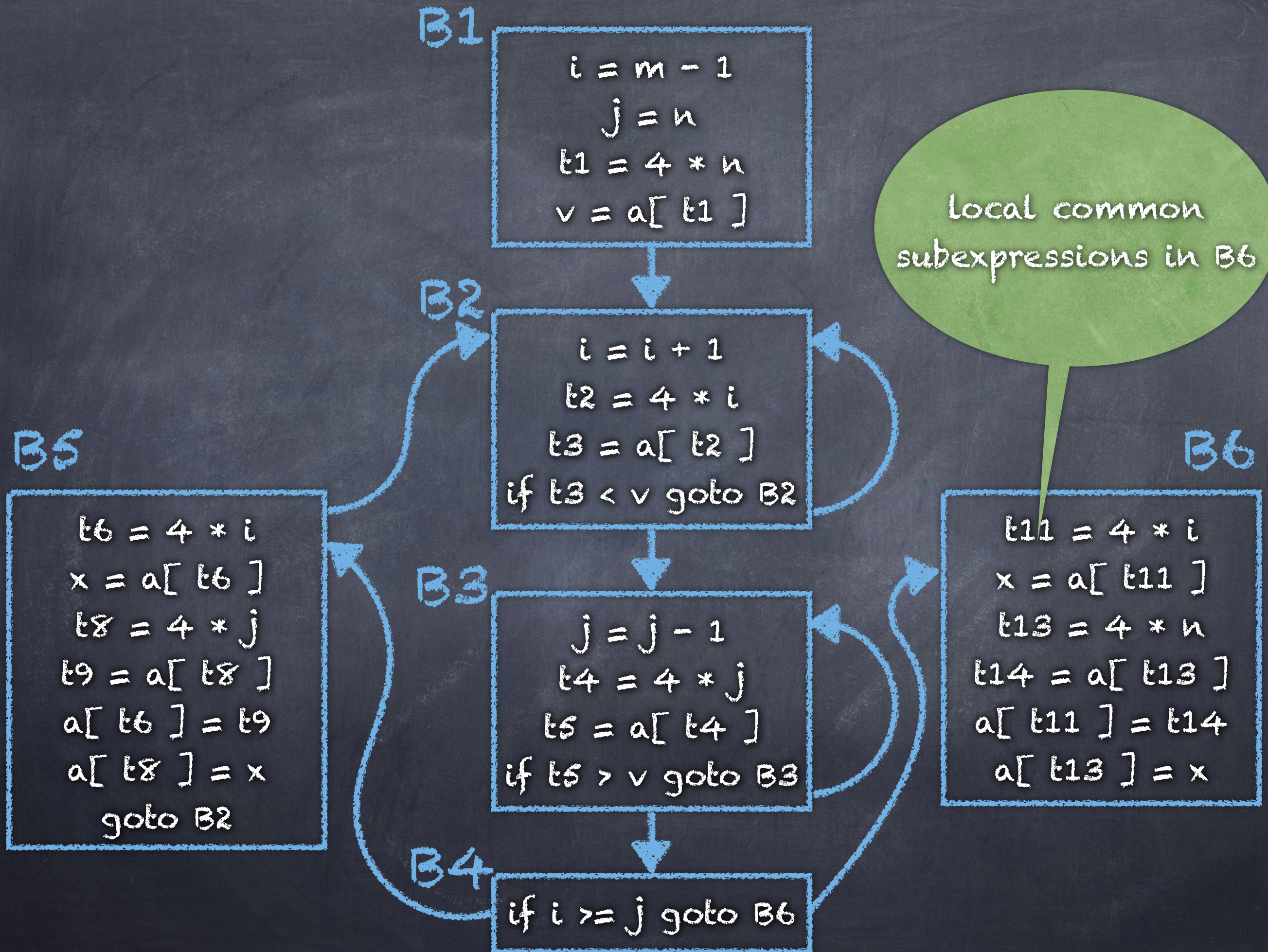
```
t11 = 4 * i
x = a[ t11 ]
t12 = 4 * i
t13 = 4 * n
t14 = a[ t13 ]
a[ t12 ] = t14
t15 = 4 * n
a[ t15 ] = x
```

B3

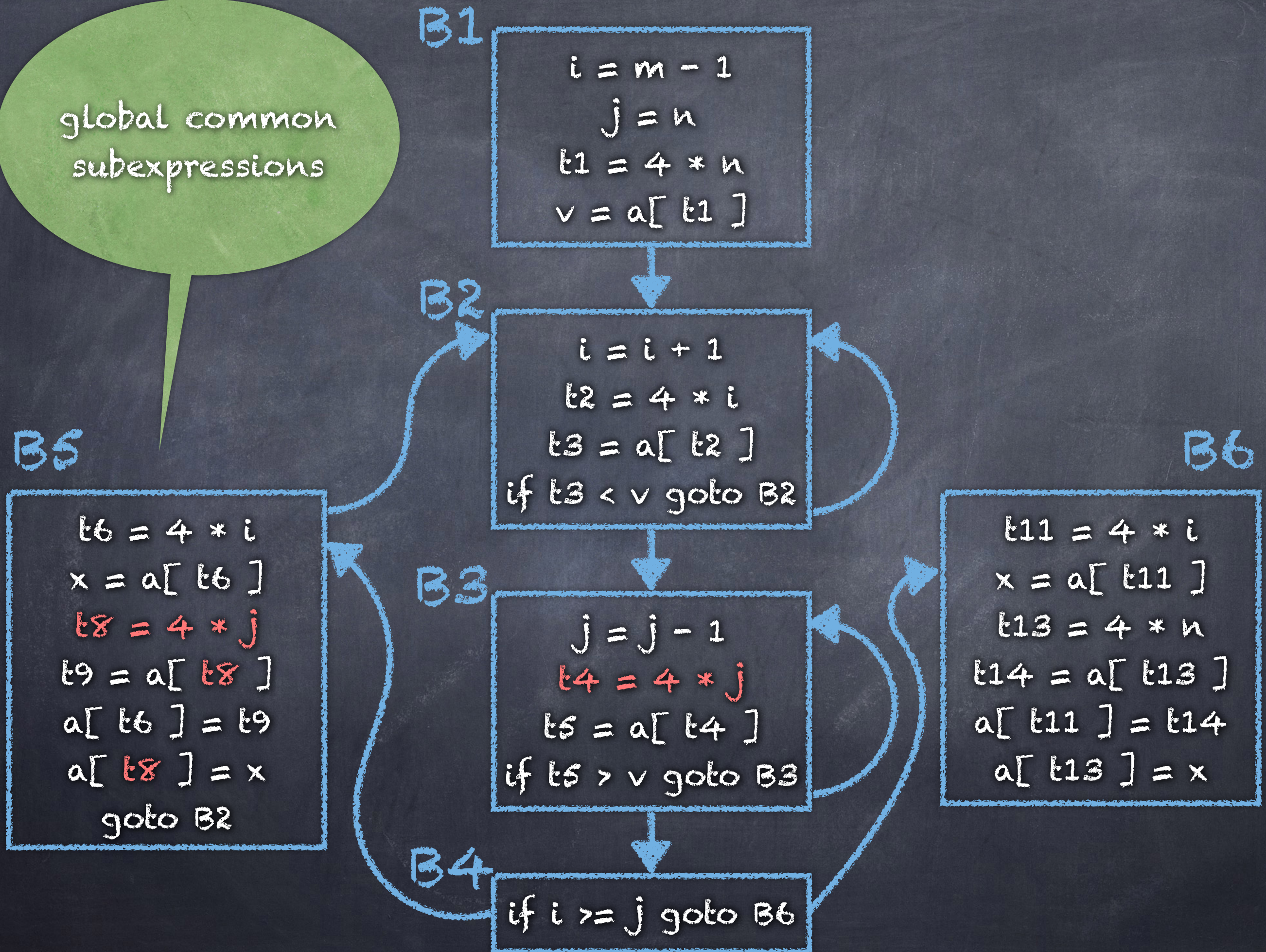
```
j = j - 1
t4 = 4 * j
t5 = a[ t4 ]
if t5 > v goto B3
```

B4

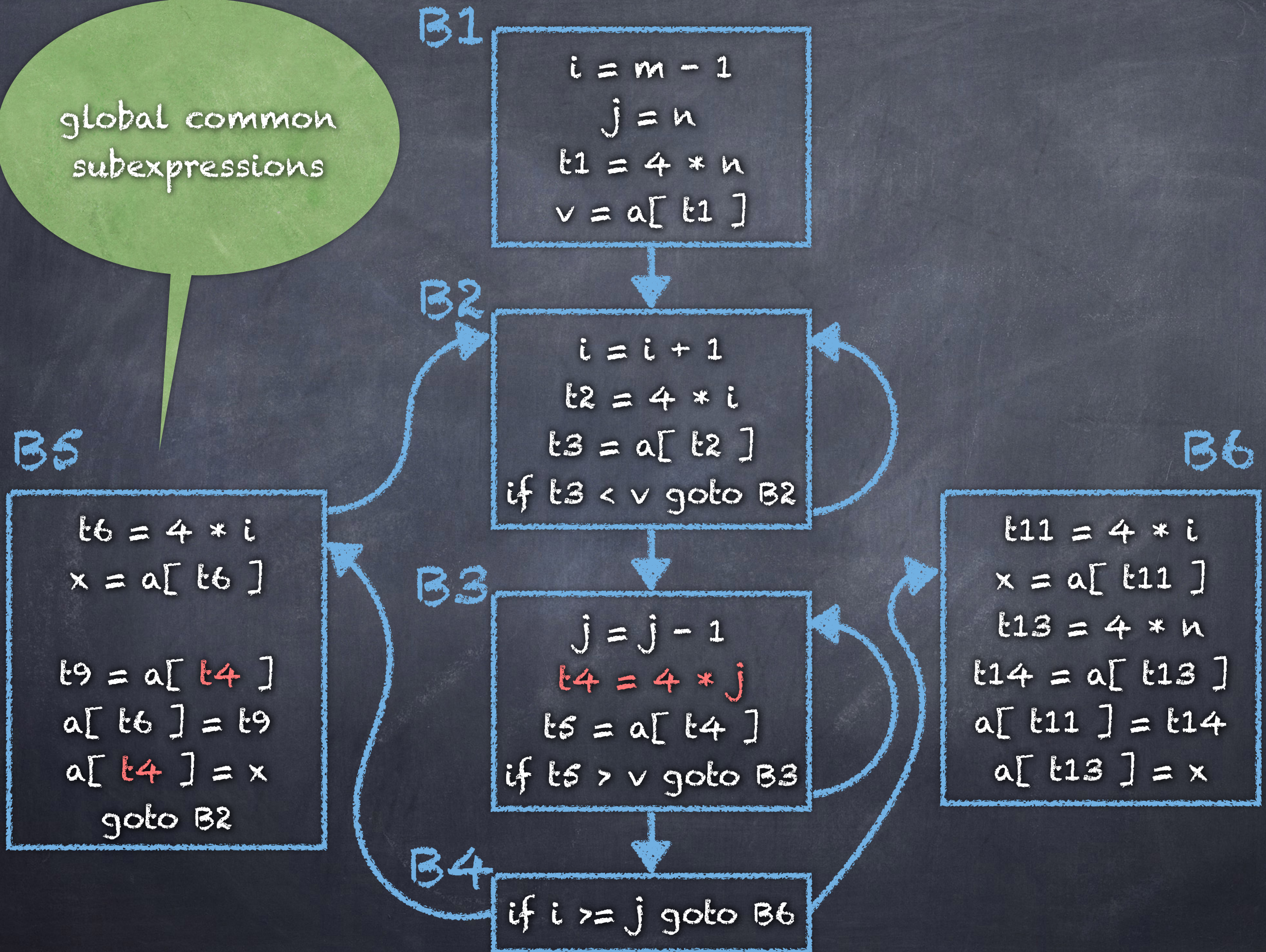
```
if i >= j goto B6
```

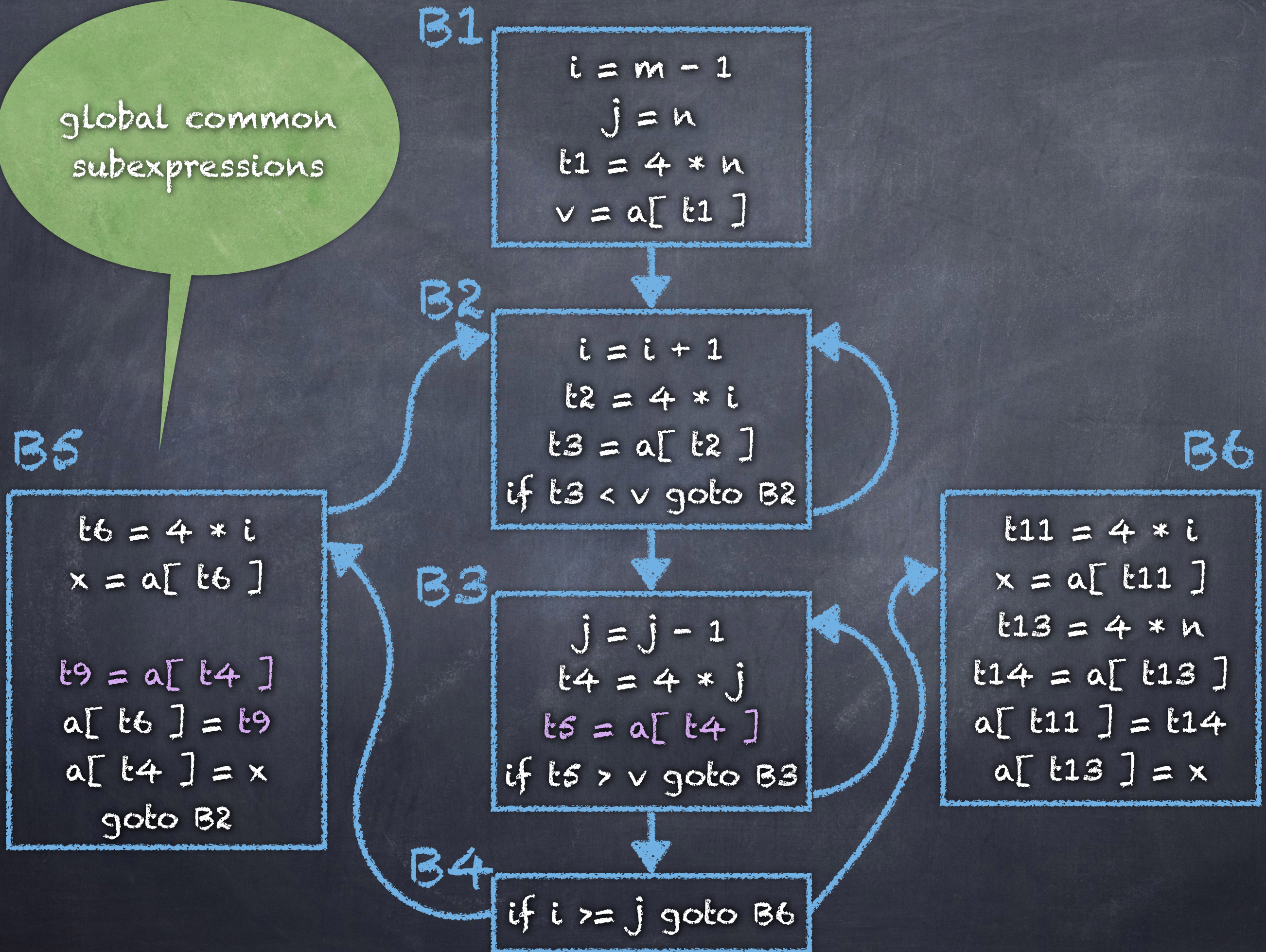
global common subexpressions



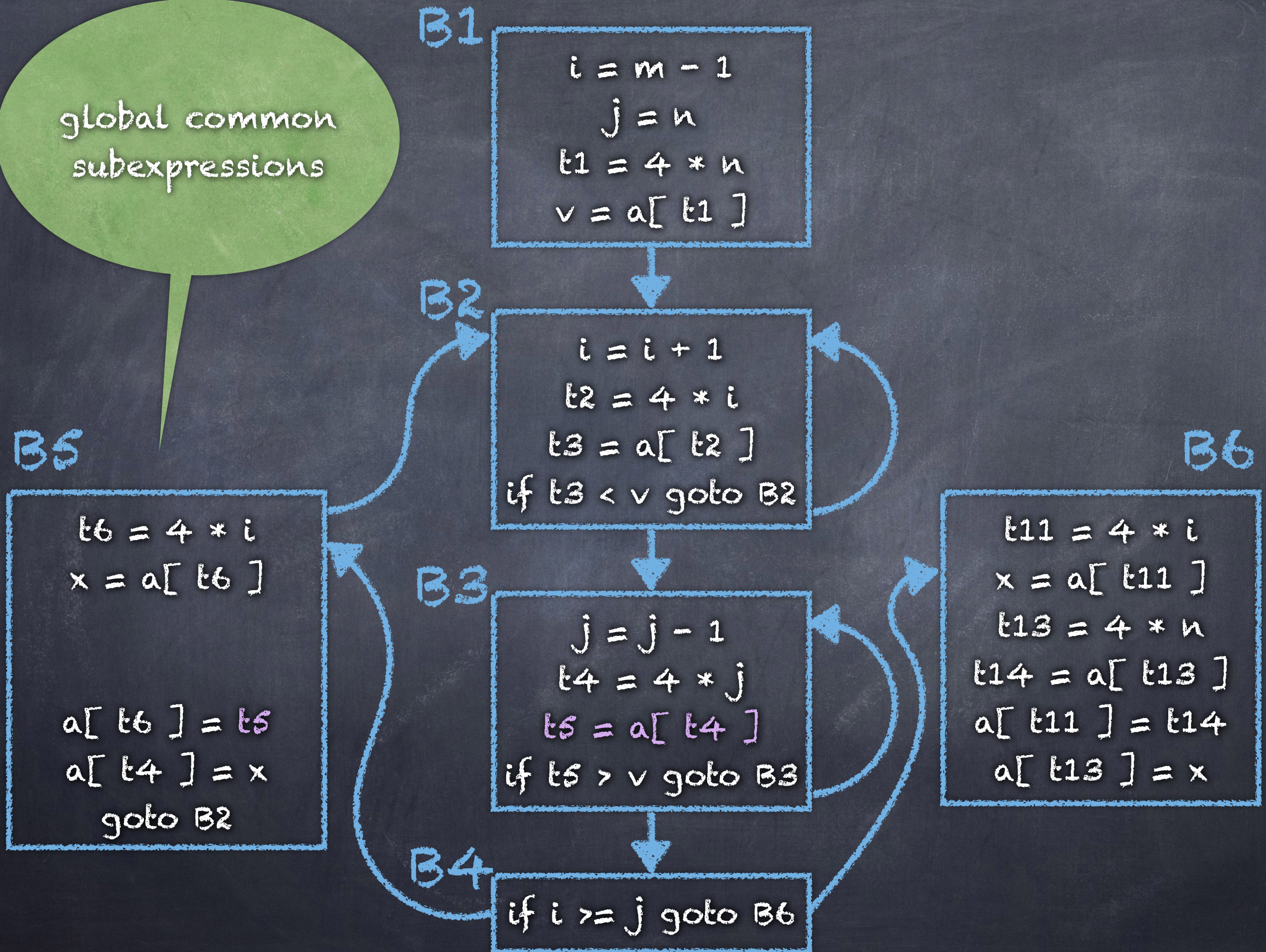
global common
subexpressions



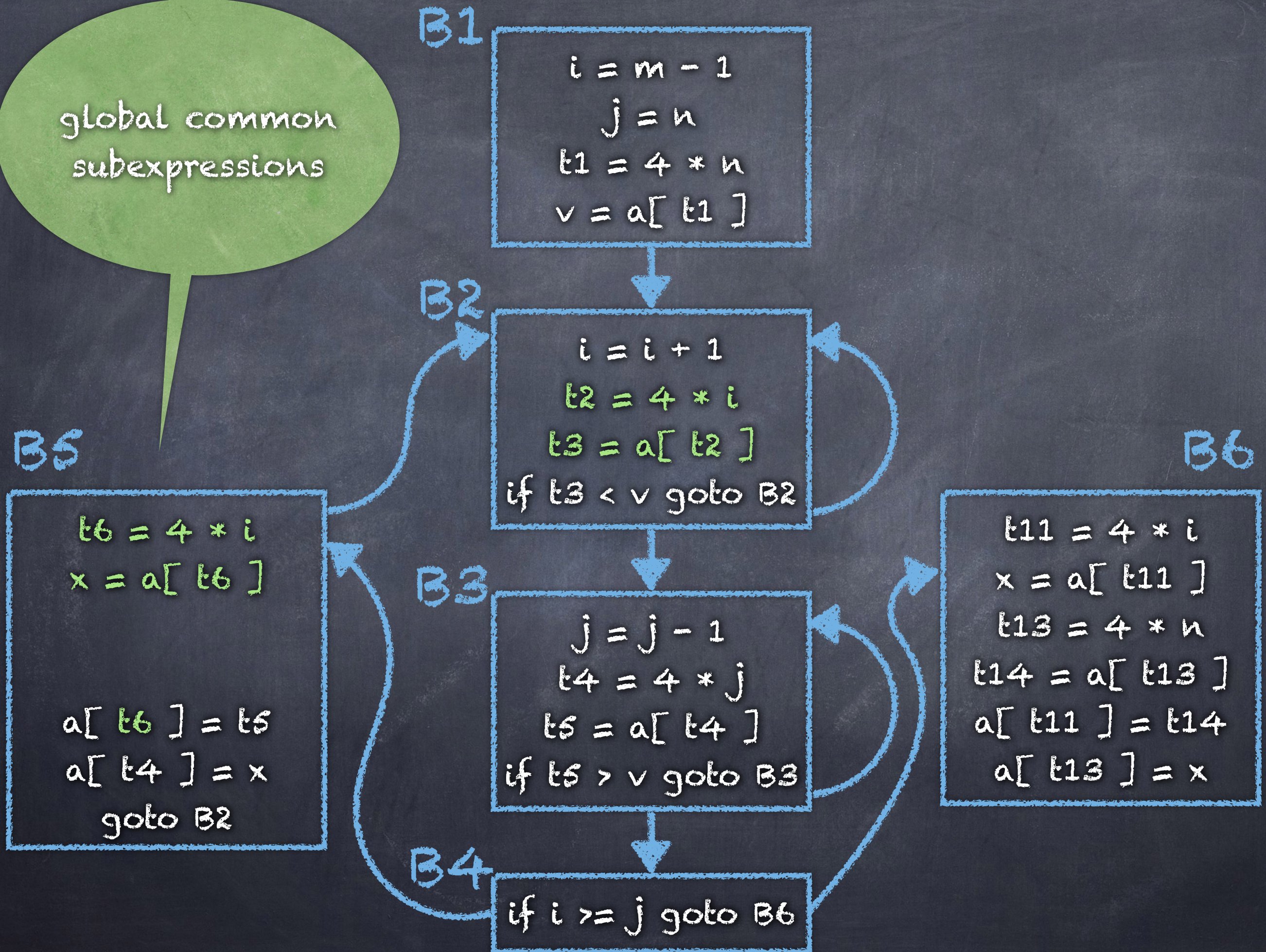
global common subexpressions



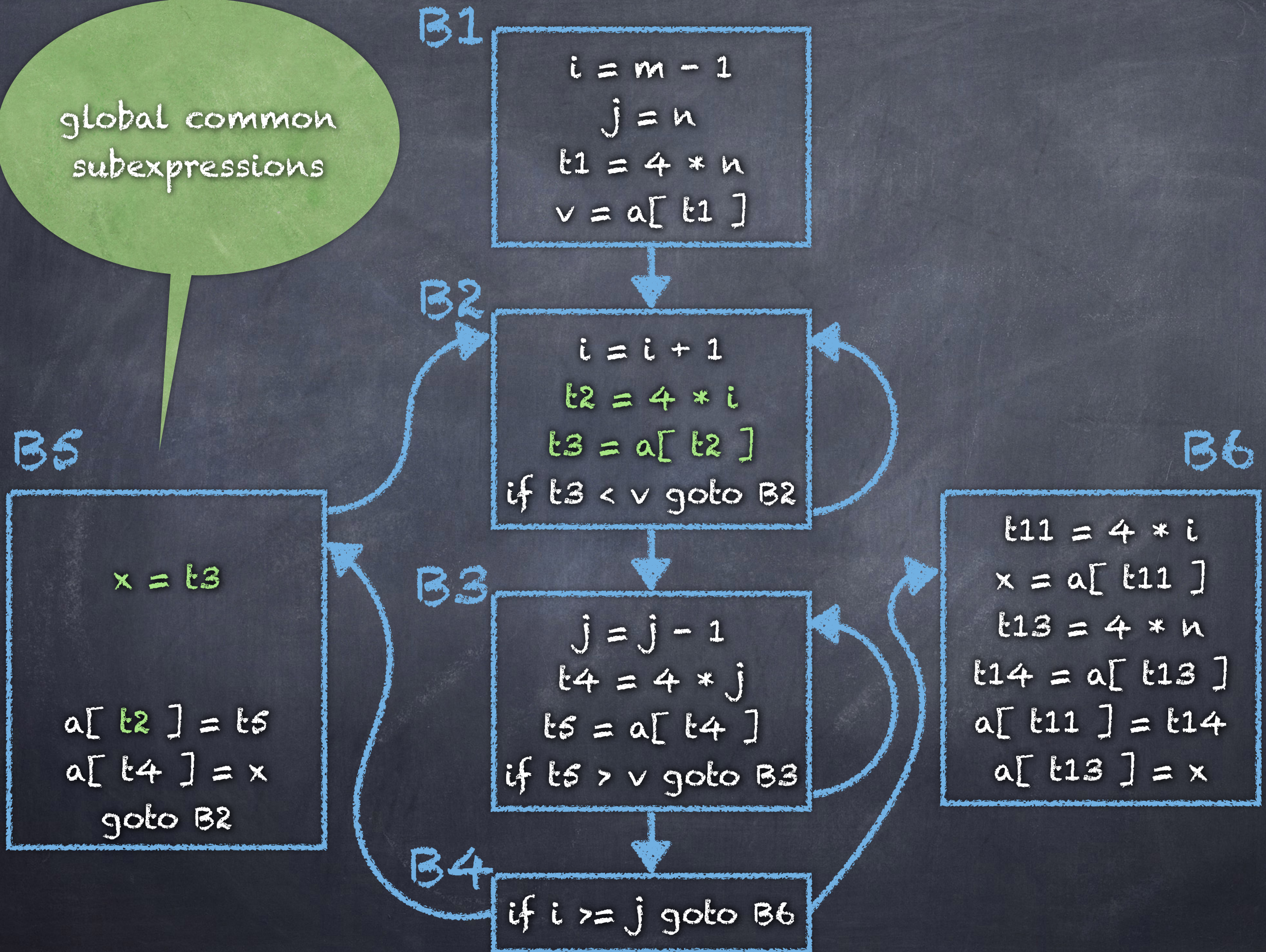
global common subexpressions



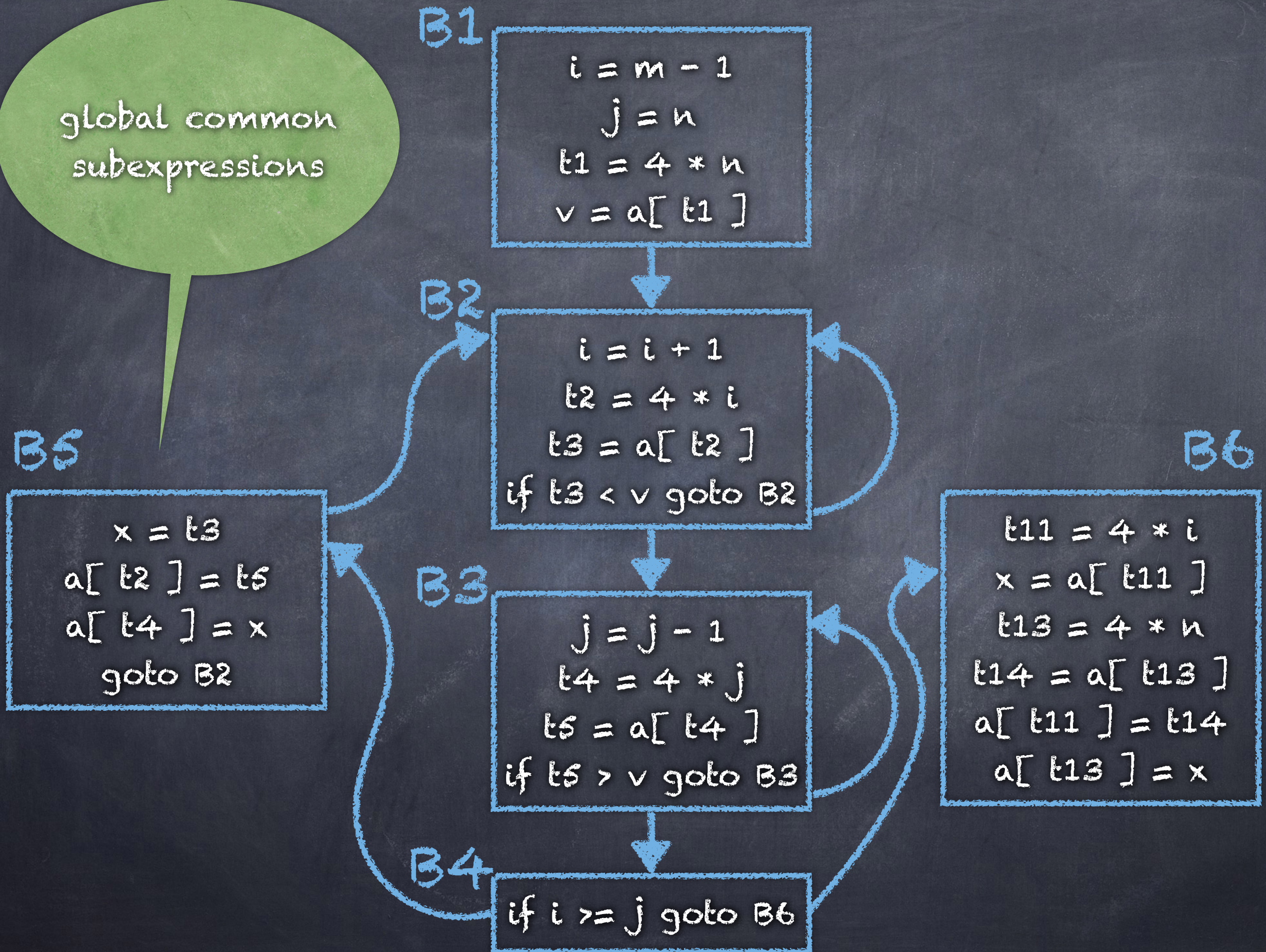
global common subexpressions

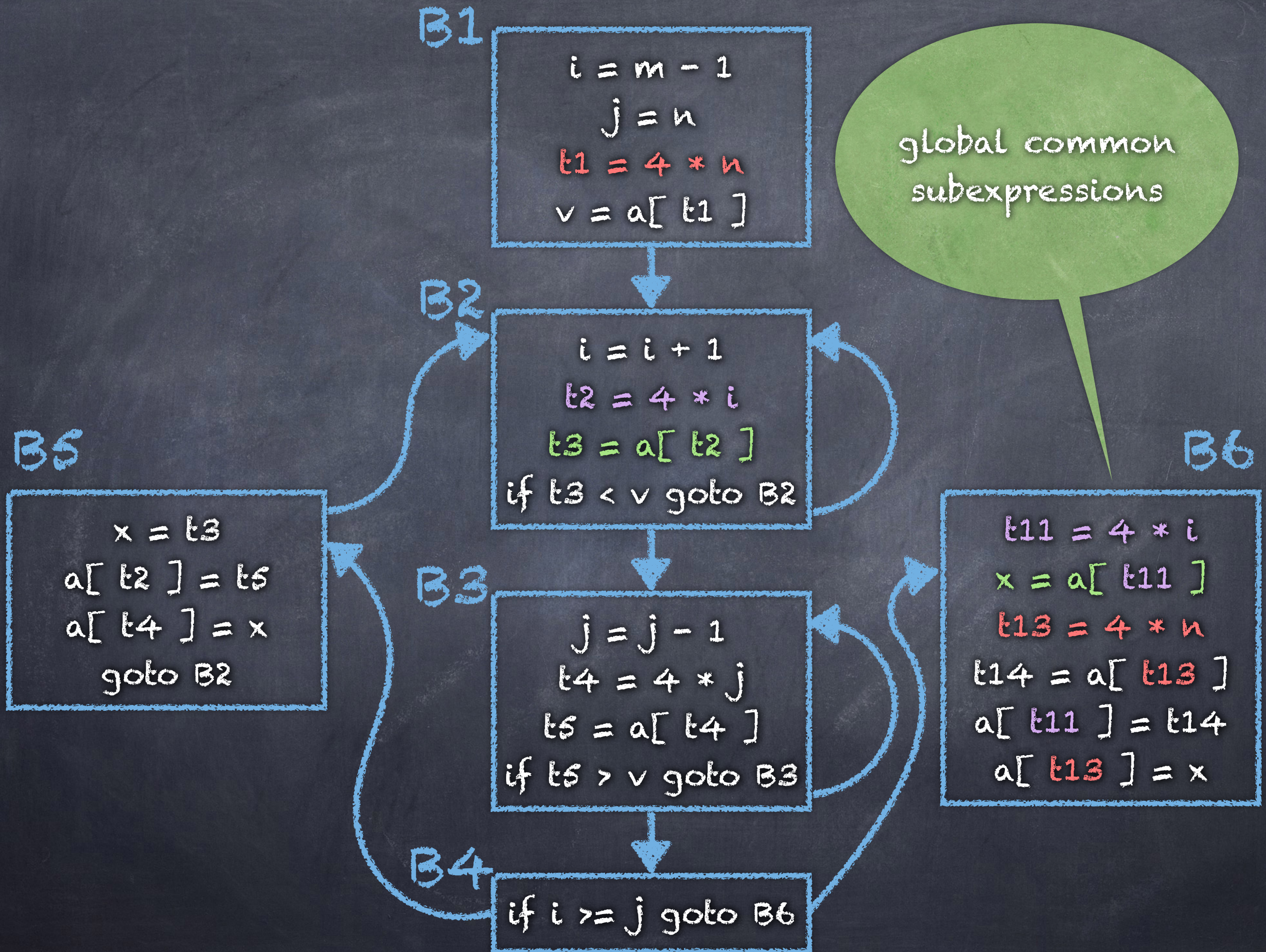


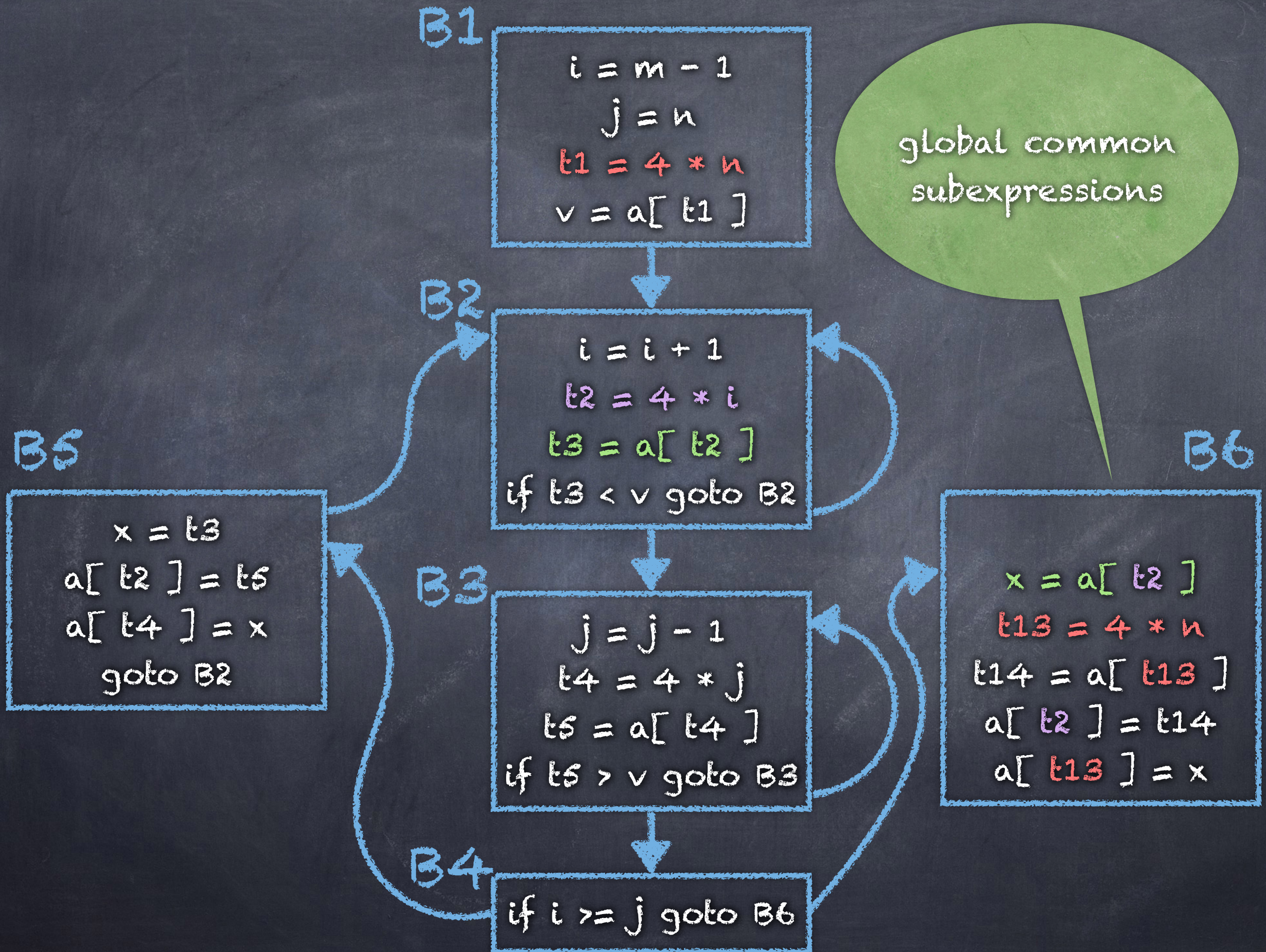
global common subexpressions

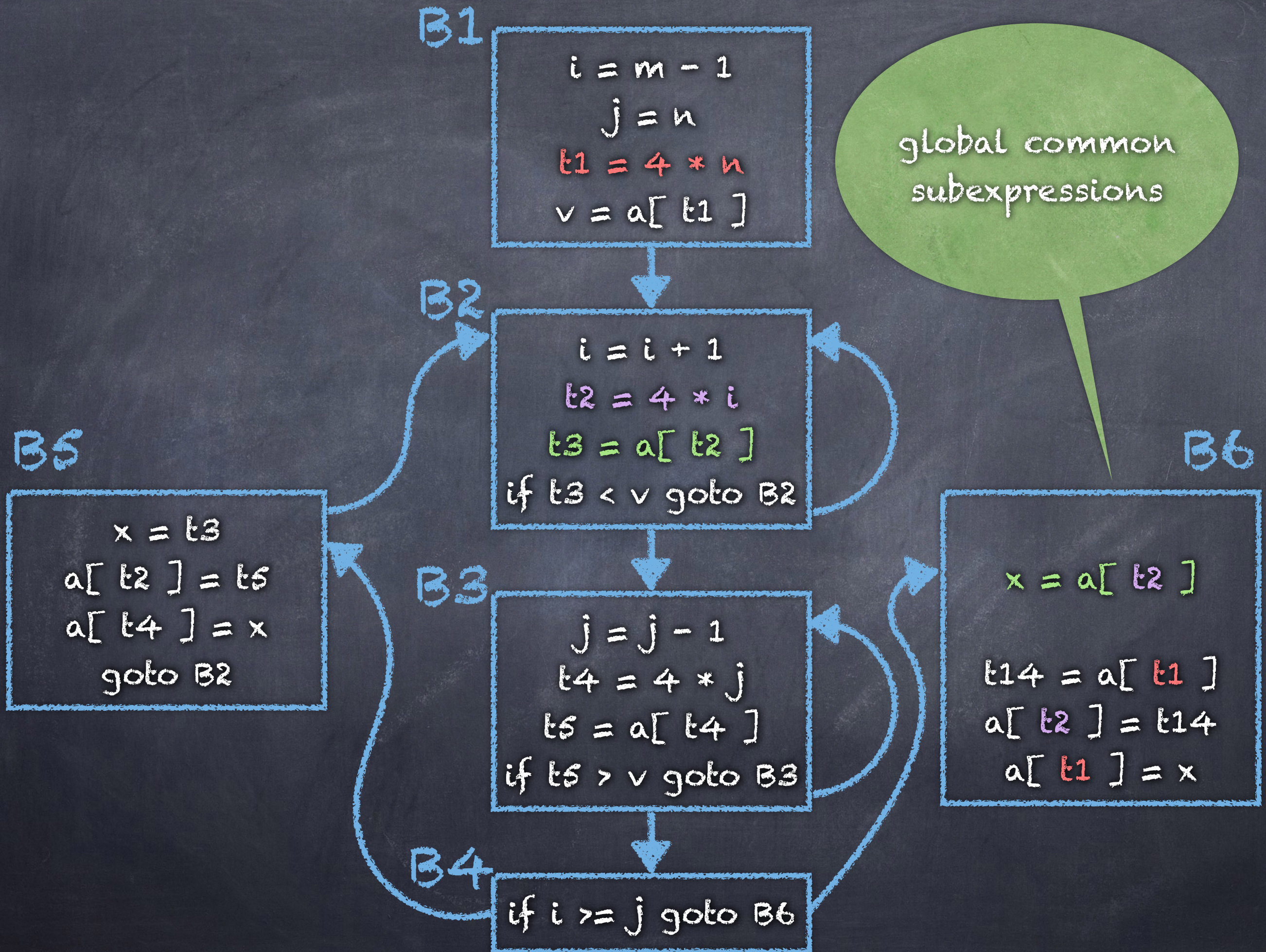


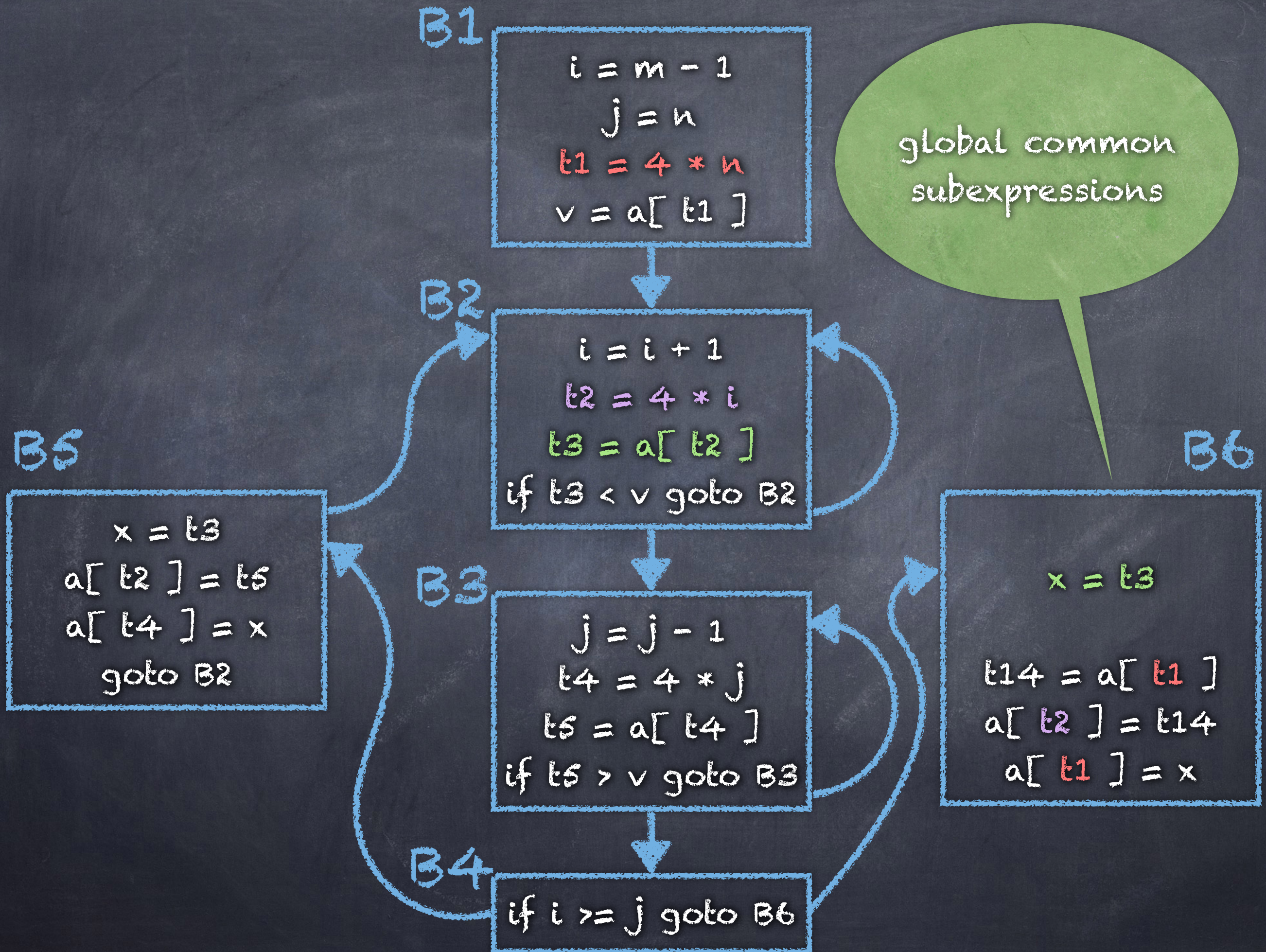
global common subexpressions



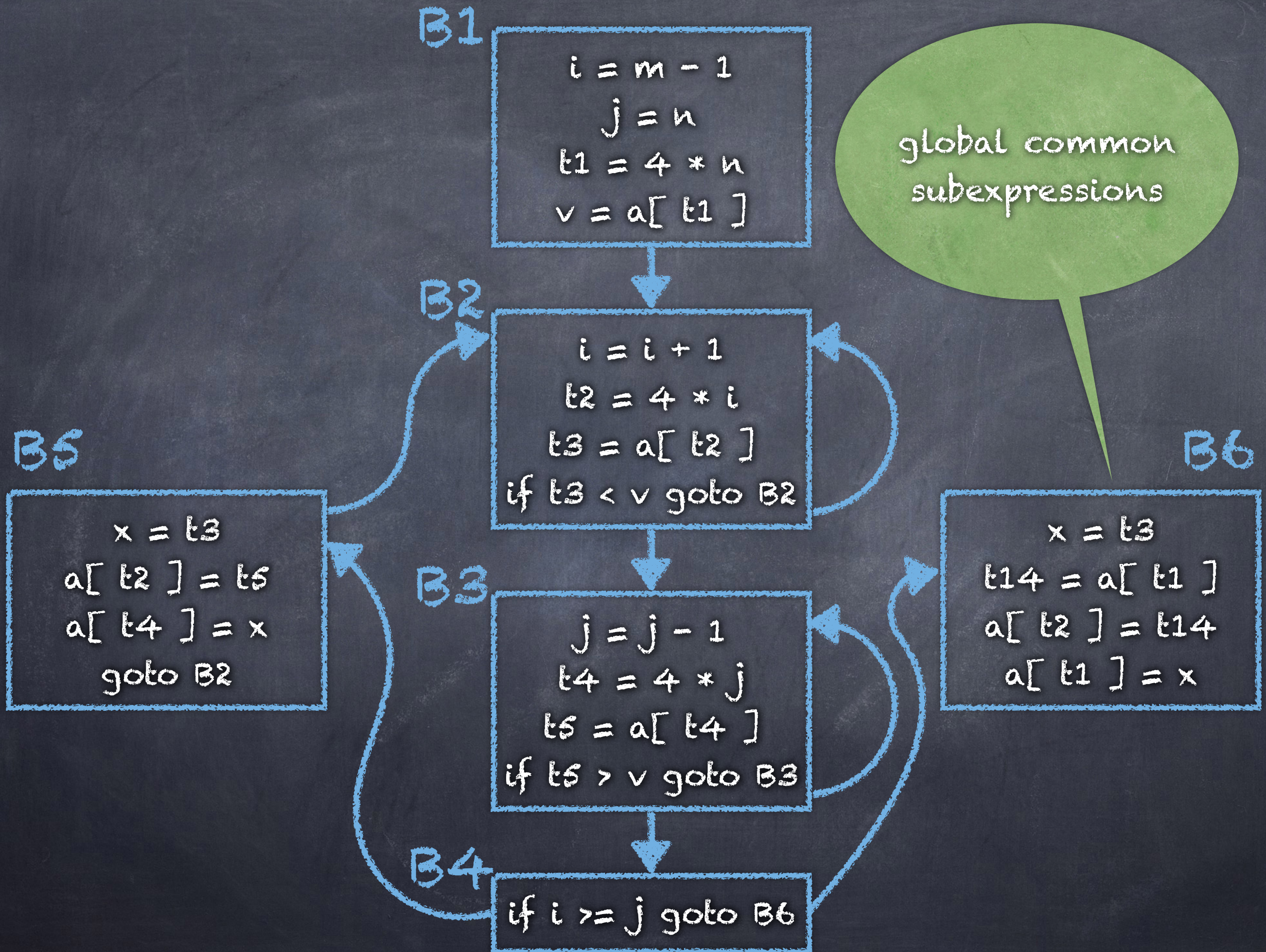




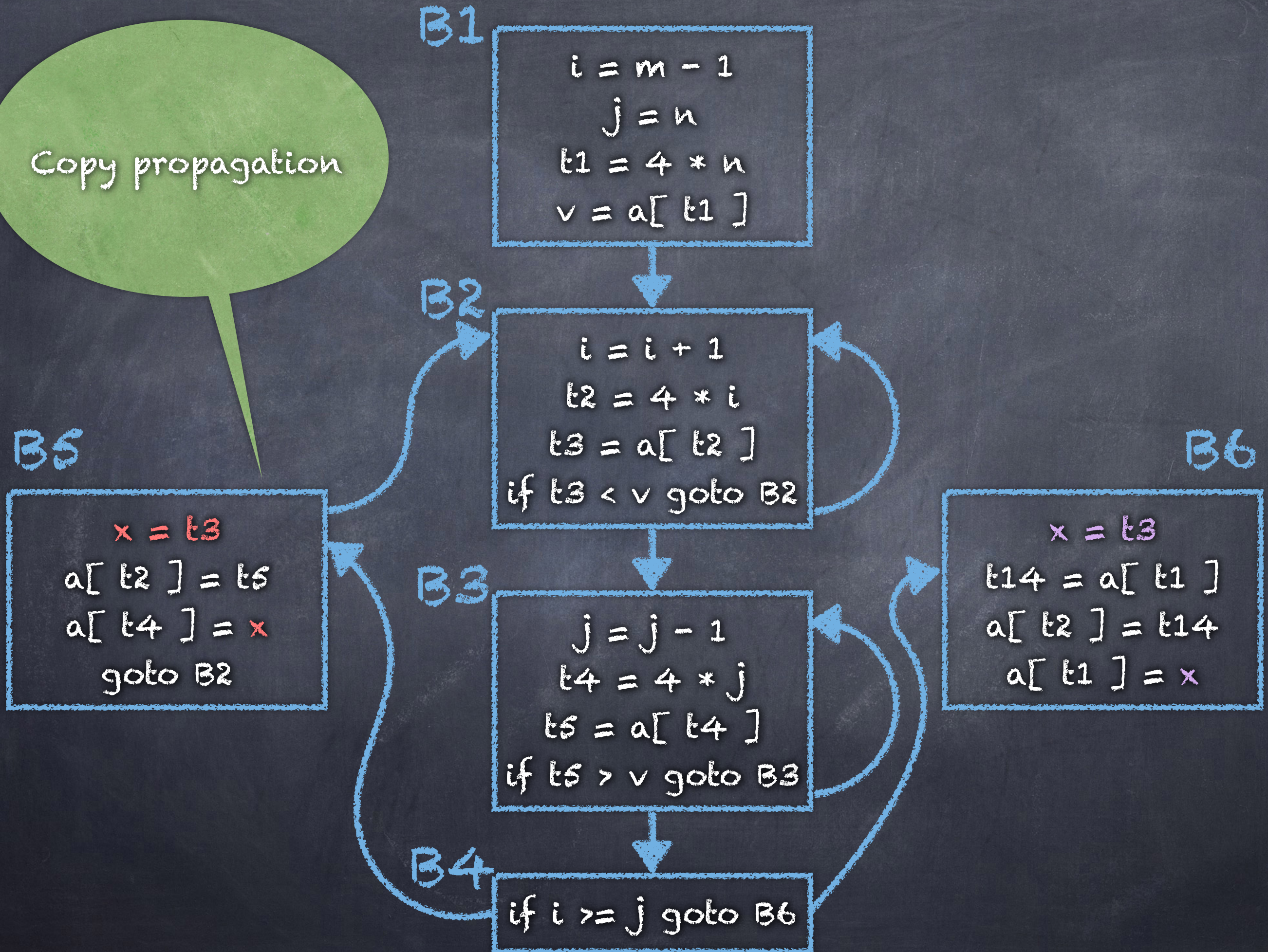




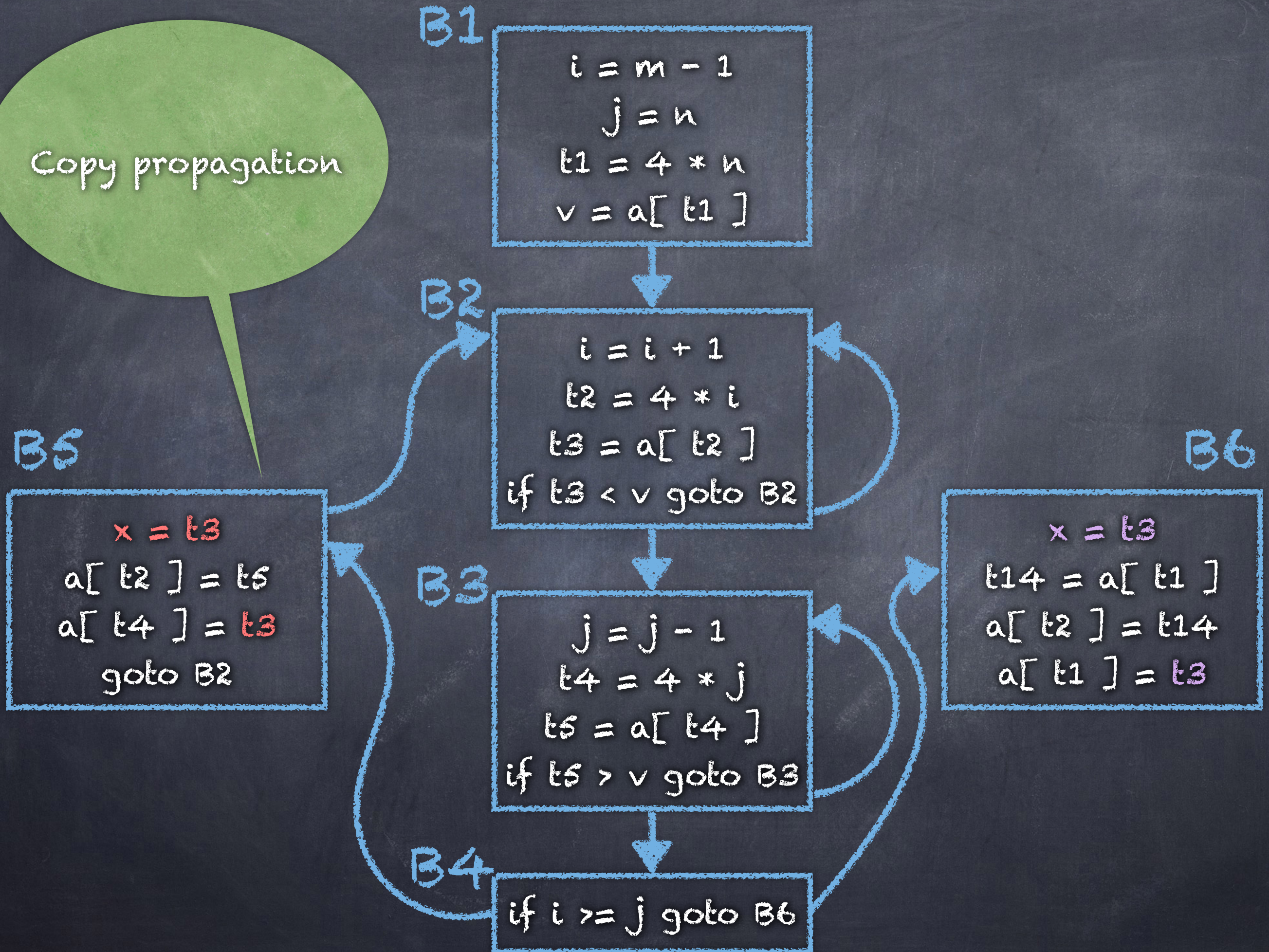
global common subexpressions



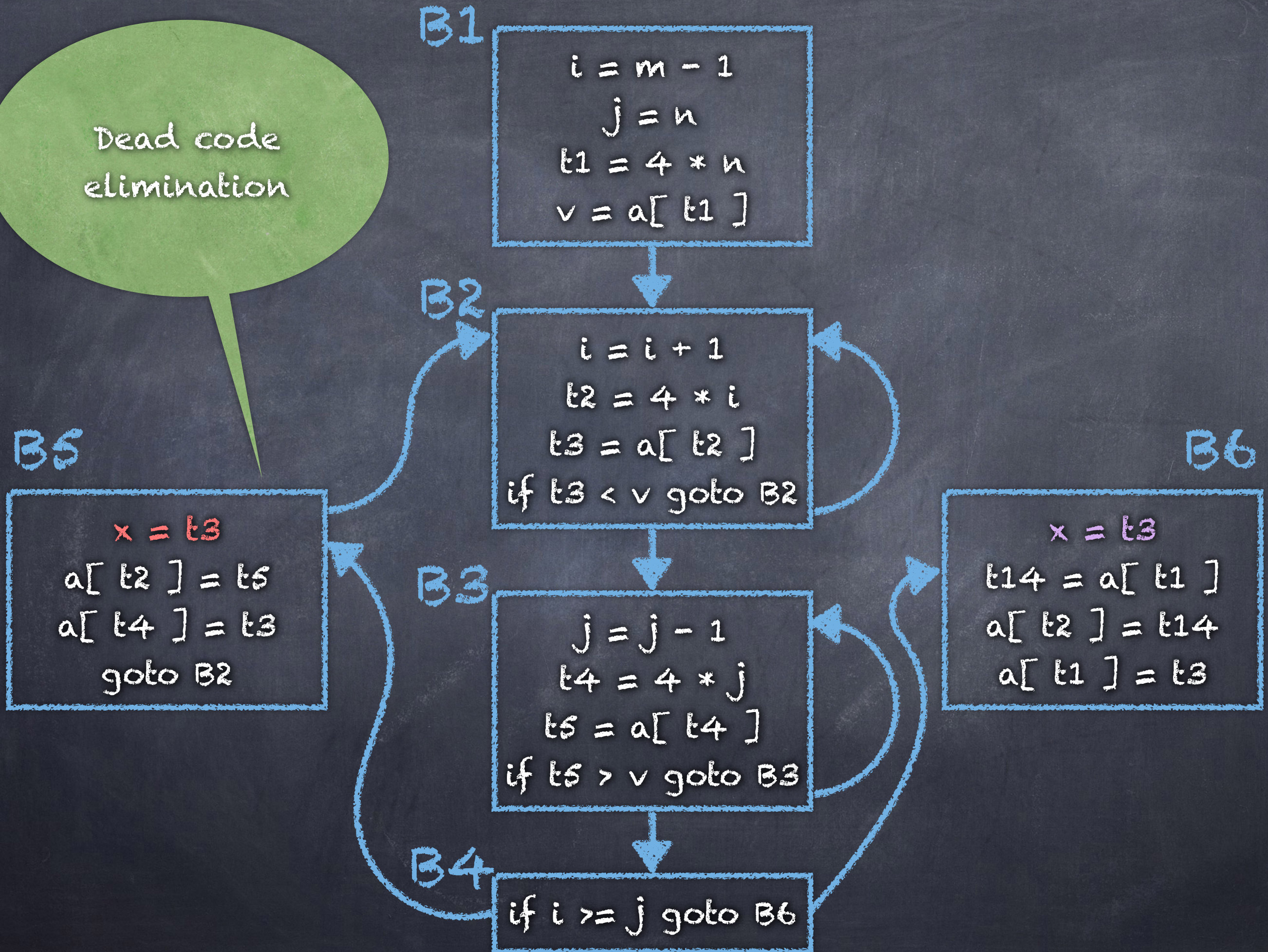
Copy propagation



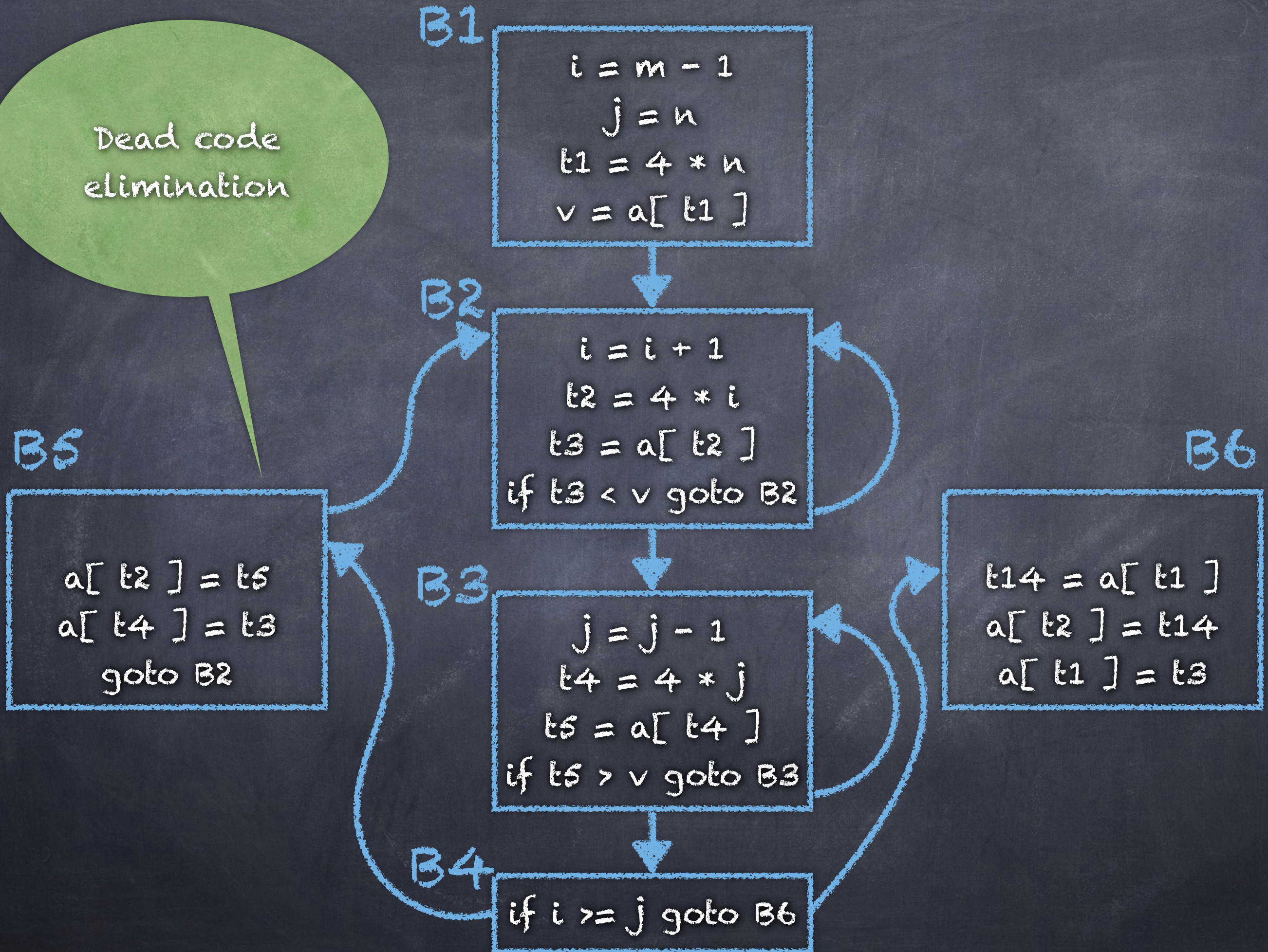
Copy propagation



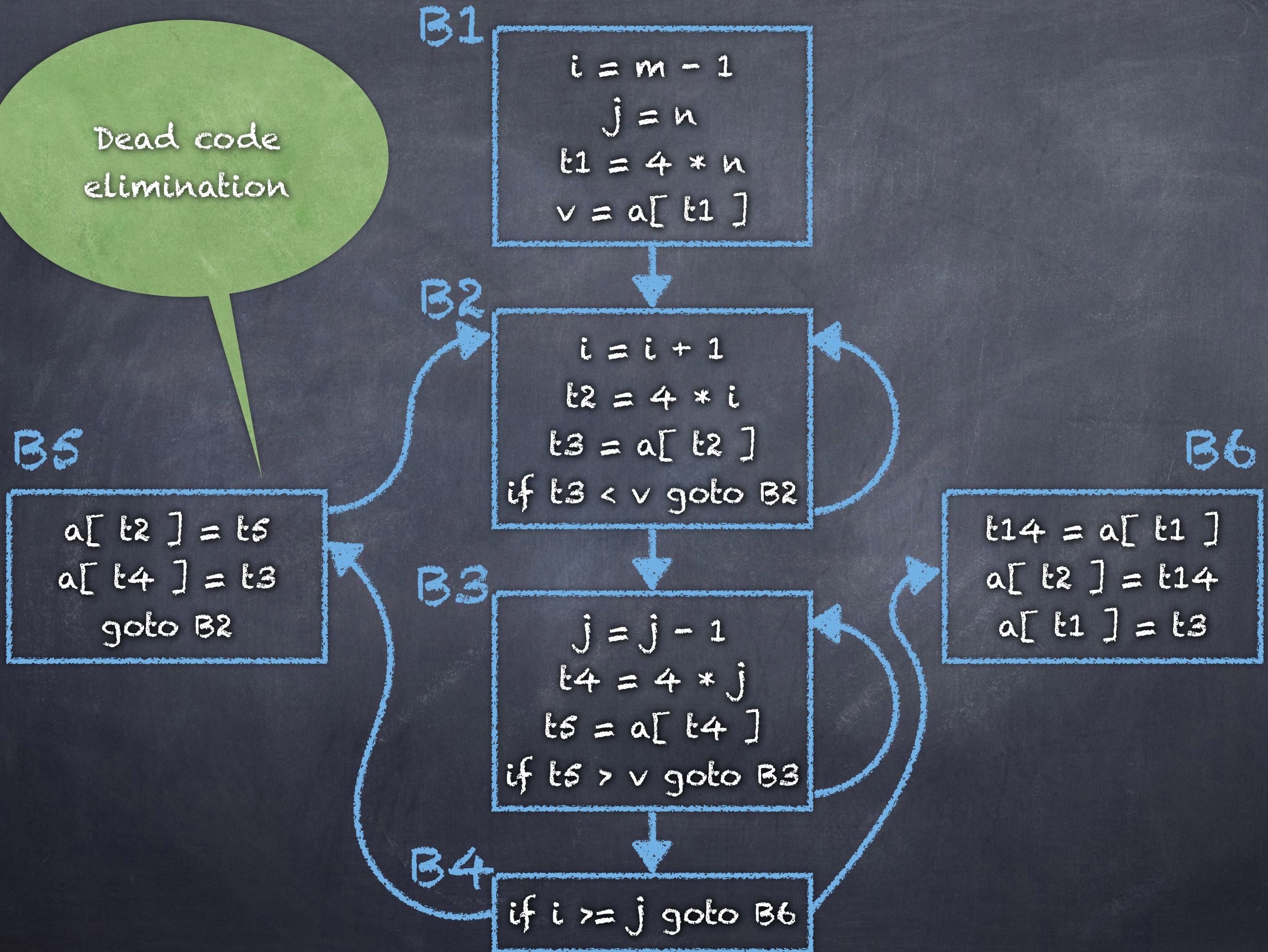
Dead code
elimination



Dead code
elimination



Dead code
elimination



Induction variables and reduction in strength

B1

```
i = m - 1  
j = n  
t1 = 4 * n  
v = a[ t1 ]
```

B2

```
i = i + 1  
t2 = 4 * i  
t3 = a[ t2 ]  
if t3 < v goto B2
```

B3

```
j = j - 1  
t4 = 4 * j  
t5 = a[ t4 ]  
if t5 > v goto B3
```

B4

```
if i >= j goto B6
```

B5

```
a[ t2 ] = t5  
a[ t4 ] = t3  
goto B2
```

B6

```
t14 = a[ t1 ]  
a[ t2 ] = t14  
a[ t1 ] = t3
```

Whenever j changes, so does t4.
* not necessary!

Induction variables and reduction in strength

Establish relationship between t4 and j

```

B1
  i = m - 1
  j = n
  t1 = 4 * n
  v = a[ t1 ]
  t4 = 4 * j

```

```

B2
  i = i + 1
  t2 = 4 * i
  t3 = a[ t2 ]
  if t3 < v goto B2

```

```

B3
  j = j - 1
  t4 = t4 - 4
  t5 = a[ t4 ]
  if t5 > v goto B3

```

```

B4
  if i >= j goto B6

```

```

B5
  a[ t2 ] = t5
  a[ t4 ] = t3
  goto B2

```

```

B6
  t14 = a[ t1 ]
  a[ t2 ] = t14
  a[ t1 ] = t3

```

Decrement t4 by 4 each time j decrements by 1

Induction variables and reduction in strength

B1
 $i = m - 1$
 $j = n$
 $t1 = 4 * n$
 $v = a[t1]$
 $t4 = 4 * j$

B2
 $i = i + 1$
 $t2 = 4 * i$
 $t3 = a[t2]$
if $t3 < v$ goto B2

B3
 $j = j - 1$
 $t4 = t4 - 4$
 $t5 = a[t4]$
if $t5 > v$ goto B3

B4
if $i \geq j$ goto B6

B5
 $a[t6] = t5$
 $a[t8] = t3$
goto B2

B6
 $t14 = a[t1]$
 $a[t2] = t14$
 $a[t1] = t3$

Induction variables and reduction in strength

Establish relationship between t_2 and I

Increment t_2 by 4 each time i increments by 1

B1

```
i = m - 1
j = n
t1 = 4 * n
v = a[ t1 ]
t4 = 4 * j
t2 = 4 * i
```

B2

```
i = i + 1
t2 = t2 + 4
t3 = a[ t2 ]
if t3 < v goto B2
```

B6

```
t14 = a[ t1 ]
a[ t2 ] = t14
a[ t1 ] = t3
```

B5

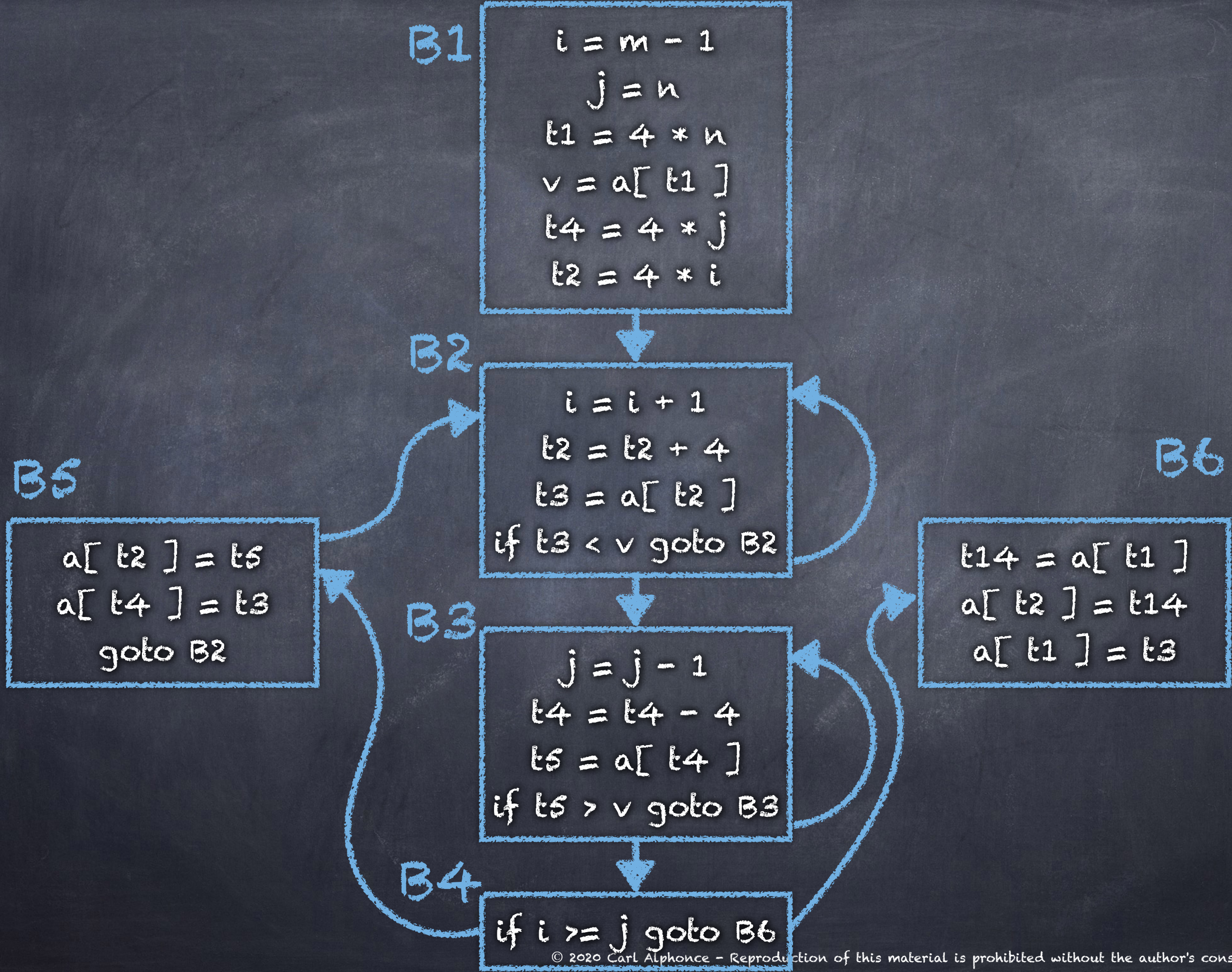
```
a[ t2 ] = t5
a[ t4 ] = t3
goto B2
```

B3

```
j = j - 1
t4 = t4 - 4
t5 = a[ t4 ]
if t5 > v goto B3
```

B4

```
if i >= j goto B6
```

Compare to
initial graph

B1

```
i = m - 1
j = n
t1 = 4 * n
v = a[ t1 ]
```

B2

```
i = i + 1
t2 = 4 * i
t3 = a[ t2 ]
if t3 < v goto B2
```

B3

```
j = j - 1
t4 = 4 * j
t5 = a[ t4 ]
if t5 > v goto B3
```

B4

```
if i >= j goto B6
```

B5

```
t6 = 4 * i
x = a[ t6 ]
t7 = 4 * i
t8 = 4 * j
t9 = a[ t8 ]
a[ t7 ] = t9
t10 = 4 * j
a[ t10 ] = x
goto B2
```

B6

```
t11 = 4 * i
x = a[ t11 ]
t12 = 4 * i
t13 = 4 * n
t14 = a[ t13 ]
a[ t12 ] = t14
t15 = 4 * n
a[ t15 ] = x
```