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CSE115 / CSE503 Introduction to Computer Science I Dr. Carl Alphonce 343 Davis Hall alphonce@buffalo.edu Office hours: Tuesday 10:00 AM – 12:00 PM\* Wednesday 4:00 PM – 5:00 PM Friday 11:00 AM - 12:00 PM OR request appointment via e-mail

<sup>\*</sup>Tuesday adjustments: 11:00 AM – 1:00 PM on 10/11, 11/1 and 12/6

## ANNOUNCEMENTS



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## Undergraduate TA office hours will be ramped up to meet demand.

This week:

Kira - Tuesday at 5:00 Corwyn - Wednesday at 2:00 Steven - Thursday at 1:00

### See "UTA Office Hours" table here:

www.cse.buffalo.edu/faculty/alphonce/cse115/people.php

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DATE: Tuesday October 4 TIME: 8:45 PM – 9:45 PM LOCATION: various rooms in NSC specific room/seat assignments to come COVERAGE:

lecture material up to and including 9/23 (this week) lab material up to and including lab 3 (next week) readings: all assigned up to and including 3.2 HAVE A CONFLICT?

I will ask for documentation 9/26 – 9/30 BRING: your UB card NO ELECTRONICS: cell phone, calculator, etc.

# **ELECTRONICS:**

off & away

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### Last time

variables expression evaluation object diagrams

Today class definitions in detail (terminology review) variable scope & lifetime method definitions

Coming up class relationships

## REVIEW

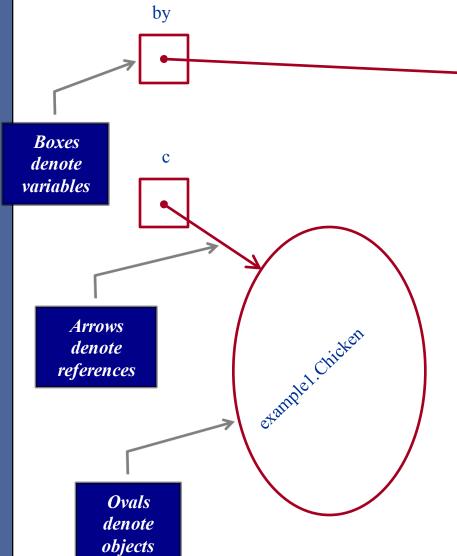


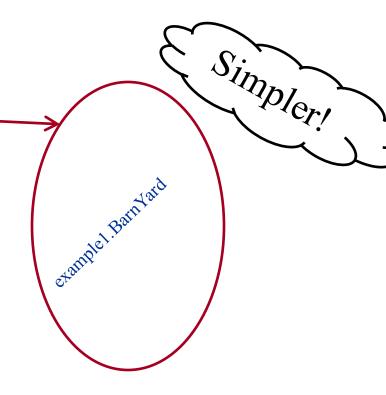
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12203	used	497362	example1.Chicken object
12204	available	497363	example1.Chicken example1.Chicken object
12205	available	497364	available
12206by	497366	497365	available
12207 c	497362	497366	example1.BarnYard object
12208	available	497367	example1.BarnYard object
12209	available	497368	example1.BarnYard object
12210	available	497369	example1.BarnYard object
12211	available	497370	used



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This diagram is an abstraction of the one on the previous slide: it ignores irrelevant details, such as the addresses and sizes of the two objects being shown. An abstraction is thus a simplification.

## MOVING ON



Here's a minimal class definition. We will label and discuss each part of it in detail next class. For now we identify the major parts:

package lab2;



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### Package declaration is shown in green:

Syntax

package lab2;



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#### package is a reserved word:

Syntax

}

package lab2;



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#### lab2 is the name of the package – you choose this

Syntax

}

package lab2;



}

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### A semicolon ; marks the end of the declaration:

You package lab2;



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### The class definition is shown in green:

Syntax

}

package lab2;



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### The class definition consists of a header . . .

Syntax

}

package lab2;



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... and a body:

Syntax

}

package lab2;



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### The class header consists of an access control modifier . . .

Syntax

package lab2;



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. . . the reserved word class . . .

Syntax

}

package lab2;



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... and a class name:

Syntax

}

package lab2;



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### The class body begins with an opening brace '{'...

Syntax

}

package lab2;



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... and ends with the matching closing brace '}':

Syntax

}

package lab2;



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### Syntax ba

}

In this example, the body consists of a single constructor definition:

package lab2;



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#### The constructor definitions consists of a header . . .

Syntax

}

package lab2;



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### ... and a body:

Syntax

}

package lab2;



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### The constructor header consists of an access control modifier . . .

Syntax

package lab2;



Syntax

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... the constructor name (which is the same as the class name) ...

package lab2;



Syntax

}

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... and a parameter list:

package lab2;



Syntax

}

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### The constructor body begins with an opening brace '{'...

package lab2;



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... and ends with the matching closing brace '}':

Syntax

}

package lab2;

# VARIABLES (more detail)



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A variable exists at a storage location in memory. For example, location 12207:

12203	
12204	
12205	
12206	
12207	
12208	
12209	
12210	
12211	

space for a variable



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a name

a location

- $\rightarrow$  in the HLL (Java)
- $\rightarrow$  in memory

- a type
- a value
- a scope

a lifetime

- → representation scheme/size
  → contents
- We'll discuss these next



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### a name → determined by declaration a location a type → determined by declaration

- a value
- a scope

a lifetime



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A variable has:

a name a location

- a type
- a value

### $\rightarrow$ determined by assignment

a scope

a lifetime

## SCOPE

(no, not the mouthwash...)



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The <u>scope</u> of a variable is the part of a program where a variable declaration is in effect.

Variables declared in different ways have different scope:

local variables

instance variables



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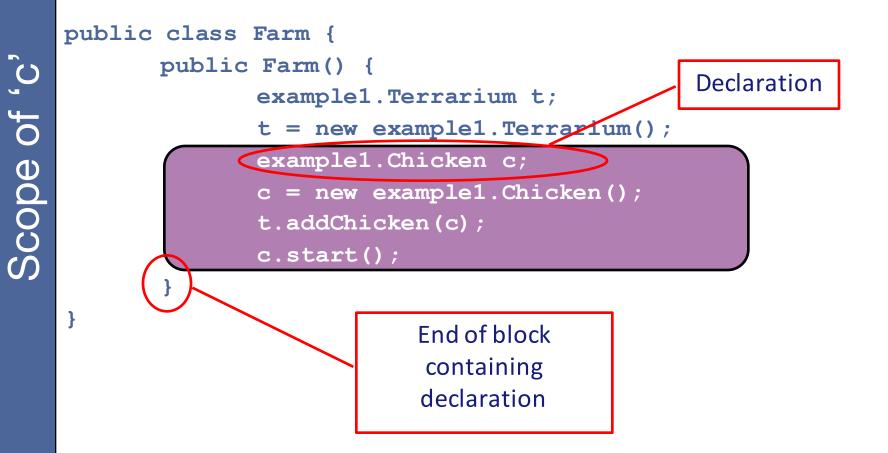
A variable declared within a constructor (or a method) is called a *local variable*.

The <u>scope</u> of a local variable is from the point of the declaration to the end of the brace-delimited block containing the declaration.



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## package lab2;





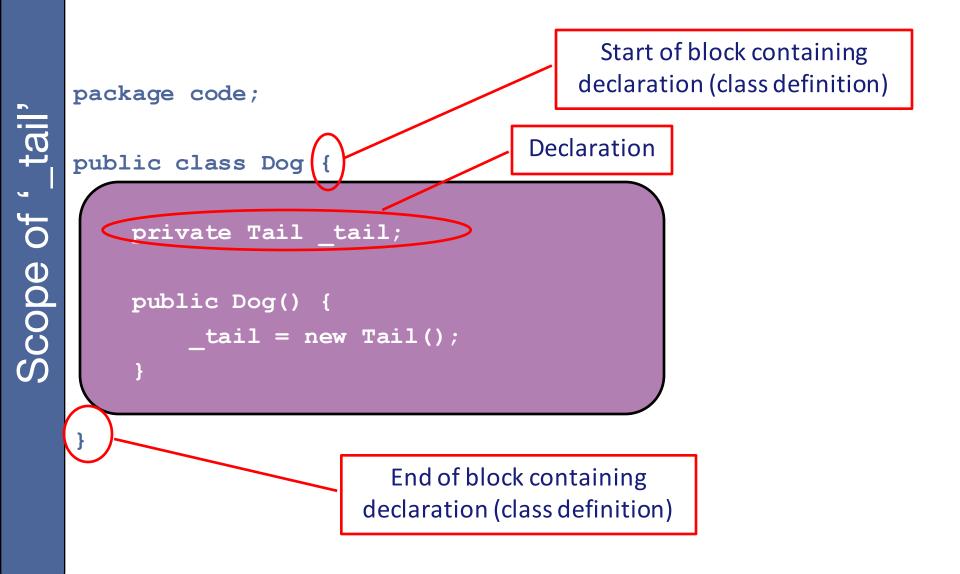
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A variable declared within a class but outside of any method is called an *instance variable*.

The <u>scope</u> of an instance variable is the entire class body.







## ETIVE

(sorry, no pun here)



Lifetime of a variable

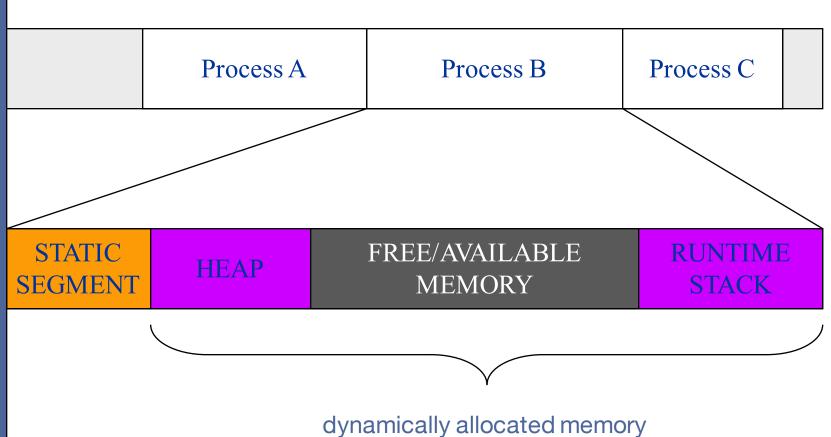
The *lifetime* of a variable is the period of time during execution of a program that the variable exists in memory. This is a dynamic property (one relating to runtime).

Variables declared in different ways have different lifetimes:

local variables

instance variables





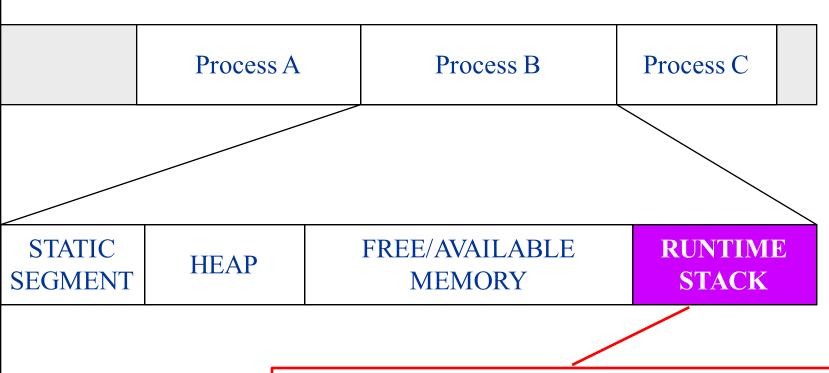


A local variable comes into existence when a method is called, and disappears when the method is completed.

Space for a local variable is allocated in a special region of memory, called the *runtime stack*.

All the local variables of a method are allocated space in the same area, called a *stack frame* (or *invocation record*).





Local variables are stored on the runtime stack. Each method invocation (call) results in an invocation record (stack frame) being added to the top of the stack. When a method exits, its invocation record is removed from the top of the stack.

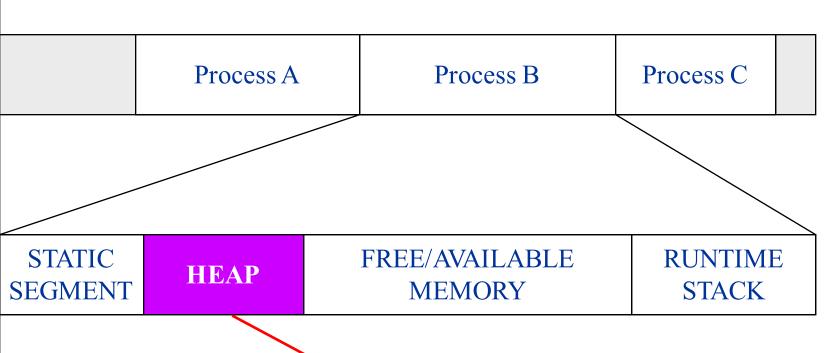


Instance variables are created when a class is instantiated. 'new' allocates memory from the heap

Each object has its own set of instance variables. the variables are the constituents of an object instance variables therefore exist on the heap

Instance variables persist as long as their objects persist as far as we know right now, objects persist until the end of the runtime of the program.





All memory allocated by 'new' comes from the heap.

Objects are allocated space by 'new', and their representations (which contain their instance variables) therefore exist on the heap.