

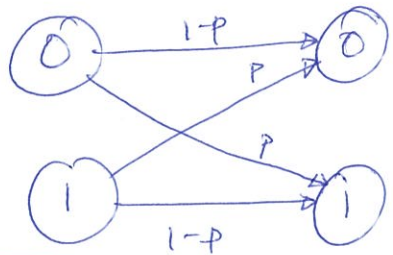
March 14, 2019

REMINDERS

- (.) HAPPY π DAY!
- (.) HW 3 out by tonight & due in 2 weeks (Mar 28 ~~April 2nd~~)
- (.) 2-page report due in ~ 2.5 weeks (April 2)
- (.) 2nd round of proof reading has started.

RECAP:

(.) BSC_p:



Each of the n transmitted bits independently get flipped w/ prob $0 \leq p \leq \frac{1}{2}$.
 $[\bar{e} \sim \text{BSC}_p$: noise vector sampled from $\text{BSC}_p]$

Shannon's Capacity Thm for BSC_p:

- $\exists \delta > 0$, $E: \{0,1\}^k \rightarrow \{0,1\}^n$, $D: \{0,1\}^n \rightarrow \{0,1\}^k$ for $k \leq \lfloor (1 - H(p) + \epsilon)n \rfloor$ s.t. $\Pr [D(E(\bar{m}) + \bar{e}) \neq \bar{m}] \leq 2^{-\delta n}$
 where $\bar{m} \in \{0,1\}^k$, $\bar{e} \sim \text{BSC}_p$.
 (Note: $1 - H(p) + \epsilon$ is the quantitative diff)
- If $k \geq \lfloor (1 - H(p) + \epsilon)n \rfloor$, then $\forall E: \{0,1\}^k \rightarrow \{0,1\}^n$, $D: \{0,1\}^n \rightarrow \{0,1\}^k$
 $\exists \bar{m} \in \{0,1\}^k$ s.t. $\Pr_{\bar{e} \sim \text{BSC}_p} [D(E(\bar{m}) + \bar{e}) \neq \bar{m}] \geq \frac{1}{2}$.

$\Rightarrow$ Capacity of BSC_p is $1 - H(p)$

Pf sketch of ①: Pick E at random. $D = \text{MLDE}$

(Step 1) For random E , average (over $\bar{m} \in \{0,1\}^k$) decoding error is $2^{-\Omega(n)}$.

(Step 2) Show similar result $\forall \bar{m} \in \{0,1\}^k$ by dropping half of codewords.

TODAY: Pf reader $\rightarrow$ Anna.

- (.) Prove ① of Shannon's result.
- (.) Overview of Pf of part ② $\rightarrow$ Reading assignment: Sec 6.3.1 in book.

2 sources of randomness:

- ① choice of $E \leftarrow$ for prob method
- ② noise from BSC $\leftarrow$ contribute to decoding error
 $\bar{e} \sim \text{BSC}_p$

"Pf" by picture:

fix $m \in \{0,1\}^k$

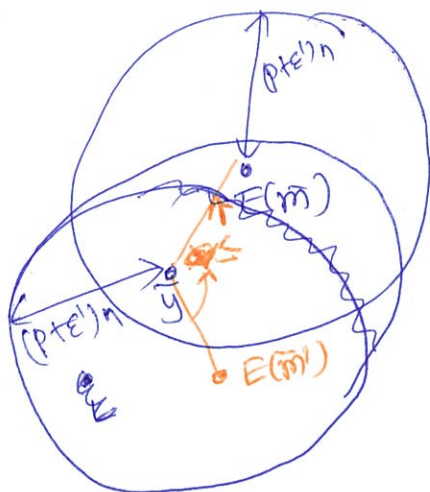
Upper bound

$$\bar{y} = E(m) + \bar{e}$$

$$\Pr_{\bar{e} \sim \text{BSC}_p} [D(E(m) + \bar{e}) \neq m] \quad (\leq 2^{-\Omega(n)})$$

By Chernoff bound up $\leq 2^{-\Omega(n)}$

$$\Pr [\text{wt}(\bar{e}) \geq (p + \epsilon')n] \leq 2^{-\Omega(n)}$$



$$m' = D(E(m) + \bar{e}) \neq m$$

$$\Rightarrow \Delta(\bar{y}, E(m)) \leq \Delta(\bar{y}, E(m')) \leq (p + \epsilon')n$$

Dec error prob (via $m' \neq m$)

$$\leq \frac{\text{Vol}_2((p + \epsilon')n, n)}{2^n} \leq \frac{2^{H(p + \epsilon')n}}{2^n}$$

Since $\leq 2^k - 1$ choices of m'

Dec error prob

$$\begin{aligned} &\leq (2^k - 1) \cdot 2^{-n(1 - H(p + \epsilon'))} = 2^{-n(1 - H(p + \epsilon'))} \\ &\leq 2^k \cdot 2^{-n(1 - H(p + \epsilon'))} \\ &\leq 2^{n(1 - H(p + \epsilon')) - n(1 - H(p + \epsilon'))} = 2^{-n(H(p + \epsilon') - H(p + \epsilon))} \\ &= 2^{-n(H(p + \epsilon') - H(p + \epsilon))} \leq 2^{-\Omega(n)} \text{ if } \epsilon' = \frac{\epsilon}{2}. \end{aligned}$$

Probability tools

(I) Chernoff bound: $X_1, \dots, X_m$ i.i.d. r.v. $X = \sum_{i=1}^n X_i$

$$\Pr [X - \mathbb{E}[X] > \delta m] \leq e^{-\delta^2 m / 2}$$

(II) Linearity expectation: r.v.s $Y_1, \dots, Y_m$

$$\mathbb{E} \left(\sum_{i=1}^m Y_i \right) = \sum_{i=1}^m \mathbb{E}[Y_i]$$

(III) Union bound:

$$\Pr \left[\bigcup_{i=1}^m E_i \right] \leq \sum_{i=1}^m \Pr[E_i]$$

(IV) For independent r.v. X, Y ; $\mathbb{E}[X \cdot Y] = \mathbb{E}[X] \cdot \mathbb{E}[Y]$

Application of (I)+(II): $\bar{e} = (e_1, \dots, e_n) \sim \text{BSC}_p$

$$\mathbb{E}[e_i] = p \cdot 1 + (1-p) \cdot 0 = p$$

$$\mathbb{E}[\text{wt}(\bar{e})] = \mathbb{E} \left(\sum_{i=1}^n e_i \right)$$

$$\stackrel{(II)}{\Rightarrow} = \sum_{i=1}^n \mathbb{E}[e_i] = pn.$$

$$(I) \Rightarrow \Pr [\text{wt}(\bar{e}) - pn > \delta n] \leq e^{-\delta^2 n / 2}$$

Pf of part (I): $\bar{y} = E(\bar{m}) + \bar{e}$

fix $\bar{m} \in \{0, 1\}^k$

$\bar{e} \sim \text{BSC}_p$

$$\Pr_{\bar{e} \sim \text{BSC}_p} [D(\bar{y}) \neq \bar{m}] = \sum_{\bar{y} \in \{0, 1\}^n} \Pr[\bar{y} | E(\bar{m})] \cdot \mathbb{1}_{D(\bar{y}) \neq \bar{m}}$$

was rec'd
was transmitted

$$= \sum_{\bar{y} \in B(E(\bar{m}), (p+\epsilon')n)} \Pr[\bar{y} | E(\bar{m})] \cdot \mathbb{1}_{D(\bar{y}) \neq \bar{m}} + \sum_{\bar{y} \notin B(E(\bar{m}), (p+\epsilon')n)} \Pr[\bar{y} | E(\bar{m})] \cdot \mathbb{1}_{D(\bar{y}) \neq \bar{m}}$$

$$\Rightarrow \left(\text{ERR}^E(\bar{m}) \right) \leq e^{-\epsilon'^2 n / 2}$$

$$+ \sum_{\bar{y} \in B(E(\bar{m}), (p+\epsilon')n)} \Pr[\bar{y} | E(\bar{m})] \cdot \mathbb{1}_{D(\bar{y}) \neq \bar{m}}$$

$$\leq \sum_{\bar{y} \notin B(E(\bar{m}), (p+\epsilon')n)} \Pr[\bar{y} | E(\bar{m})] \leq 1$$

$$\stackrel{II}{=} \Pr[\text{wt}(\bar{e}) > (p+\epsilon')n] \leq e^{-\epsilon'^2 n / 2}$$

By linearity of exp:

$$\begin{aligned} \mathbb{E}_E [\text{Err}^E(\bar{m})] &\leq e^{-\epsilon'^2 n / 2} + \sum_{y \in B(E(\bar{m}), p + \epsilon' n)} \mathbb{E}_E \left(\Pr[y | E(\bar{m})] \right) \mathbb{1}_{D(y) \neq \bar{m}} \\ &= e^{-\epsilon'^2 n / 2} + \sum_{y \in B(E(\bar{m}), p + \epsilon' n)} \Pr[y | E(\bar{m})] \cdot \mathbb{E}_E [\mathbb{1}_{D(y) \neq \bar{m}}] \end{aligned}$$

depends on BSCp but not ϵ
 $\mathbb{1}_{D(y) \neq \bar{m}}$
 $\uparrow$ depends on E but not on BSCp

Claim 1: $\mathbb{E}_E [\mathbb{1}_{D(y) \neq \bar{m}}] \leq 2^{-n (H(p + \epsilon) - H(p + \epsilon'))}$

$$\begin{aligned} \Rightarrow \mathbb{E}_E [\text{Err}^E(\bar{m})] &\leq e^{-\epsilon'^2 n / 2} + 2^{-n (H(p + \epsilon) - H(p + \epsilon'))} \\ &\leq e^{-\epsilon'^2 n / 2} + 2^{-n (H(p + \epsilon) - H(p + \epsilon'))} \sum_{y \in B(E(\bar{m}), p + \epsilon' n)} \Pr[y | E(\bar{m})] \\ &\leq e^{-\epsilon'^2 n / 2} + 2^{-n (H(p + \epsilon) - H(p + \epsilon'))} \underbrace{\sum_{y \in \{0,1\}^n} \Pr[y | E(\bar{m})]}_{=1} \\ &\leq 2^{-\delta' n} \quad \text{with } \epsilon' = \frac{\epsilon}{2}, \delta' = \Theta(\epsilon^2) \end{aligned}$$

$$\Rightarrow \mathbb{E}_E (\text{Err}^E(\bar{m})) \leq 2^{-\delta' n}$$

Since $\bar{m}$ was arbitrary $\forall \bar{m} \in \{0,1\}^k$, $\mathbb{E}_E (\text{Err}^E(\bar{m})) \leq 2^{-\delta' n}$

Prob. method: $\Rightarrow \forall \bar{m} \in \{0,1\}^k, \exists E \text{ s.t. } \text{Err}^E(\bar{m}) \leq 2^{-\delta' n}$

$$\begin{aligned} \Rightarrow \mathbb{E}_E \mathbb{E}_m (\text{Err}^E(\bar{m})) &\leq 2^{-\delta' n} \\ \text{uniform over } \{0,1\}^k \rightarrow \bar{m} & \\ \Rightarrow \mathbb{E}_E (\mathbb{E}_m (\text{Err}^E(\bar{m}))) &\leq 2^{-\delta' n} \\ m \text{ \& } E \text{ are indep.} & \\ \text{prob. method} \Rightarrow \exists E \text{ s.t. } \mathbb{E}_m (\text{Err}^E(\bar{m})) &\leq 2^{-\delta' n} \end{aligned}$$

As error $\leq 2^{-s'n} \Rightarrow$ Max error $\leq 2^{-s'n}$
(dropping half the codewords)

Claim 2: $\{0,1\}^k = \bar{m}_1, \bar{m}_2, \dots, \bar{m}_{2^k}$ $P_i = \text{Err}^E(\bar{m}_i)$

$$P_1 \leq P_2 \leq \dots \leq P_{2^k} \quad \frac{1}{2^k} \sum_{i=1}^{2^k} P_i \leq 2^{-s'n}$$

$$\Rightarrow P_{2^{k/2}} \leq 2 \cdot 2^{-s'n} \quad s^k = s' - \frac{1}{2^k} \Rightarrow 2 \cdot 2^{-s'n} \leq 2^{-s'n}$$

$\Rightarrow$ Only we $\bar{m}_1, \dots, \bar{m}_{2^{k/2}}$ in final code (dim: $k \rightarrow k-1$)

DONE! (except Claims 1+2)

Pf of Claim 2: Pf by contradiction.

Ass $P_{2^{k-1}} > 2 \cdot 2^{-s'n}$

$$2^k \cdot 2^{-s'n} \geq \sum_{i=1}^{2^k} P_i \geq \sum_{i=2^{k-1}+1}^{2^k} P_i \geq \sum_{i=2^{k-1}+1}^{2^k} P_{2^{k-1}}$$

$$> \sum_{i=2^{k-1}+1}^{2^k} 2 \cdot 2^{-s'n} = \frac{2^k}{2} \cdot 2 \cdot 2^{-s'n} = 2^k \cdot 2^{-s'n} !!$$

Claim 1: $\Pr_E [\mathbb{1}_{D(y) \neq m}] \leq 2^{-n} (H(p+\epsilon) - H(p-\epsilon'))$

$$\Pr_E (\mathbb{1}_{D(y) \neq m}) = \Pr_E [\mathbb{1}_{D(y) \neq m} | E(m)]$$

$$= \Pr_E [\exists \bar{m}' \neq m \text{ s.t. } \Delta(y, E(\bar{m}')) \leq \Delta(y, E(m)) | E(m)]$$

Def of MLD $\Rightarrow$ Union bound $\Rightarrow$

$$\leq \sum_{\bar{m}' \neq m} \Pr [\Delta(y, E(\bar{m}')) \leq \Delta(y, E(m)) | E(m)]$$

$$\Rightarrow \Delta(y, E(\bar{m}')) \leq (p+\epsilon')n$$

$$\leq \Pr [\Delta(y, E(\bar{m}')) \leq (p+\epsilon')n | E(m)]$$

$E(m)$ & $E(\bar{m})$ are indep $\Rightarrow$

$$= \sum_{\bar{m}' \neq m} \Pr [\Delta(y, E(\bar{m}')) \leq (p+\epsilon')n] = \frac{\text{Vol}_2((p+\epsilon')n, n)}{2^n}$$

$$\leq \sum_{\bar{m}' \neq m} \frac{\text{Vol}_2((p+\epsilon')n, n)}{2^n - n[1-H(p+\epsilon') - (1-H(p-\epsilon))]} \leq 2^k \frac{2^{H(p+\epsilon')n}}{2^n}$$