

Mar 18

# REMINDERS

- (\*) Mini project 2-pg report due by 11:59pm on Wed, Mar 30
- (\*) Please vote on prepared HW grade policy change
- (\*) Deadline to sign up for future proof reading: 5pm, Mon
- (\*) Piazza off for next week (Enjoy your spring break!) Mar 21

## RECAP

BSC<sub>p</sub> → Each of the  $n$  transmitted bits independently get flipped w/ prob  $p$  ( $0 \leq p \leq \frac{1}{2}$ )

(\*)  $\bar{e} \sim \text{BSC}_p \rightarrow \bar{e} \in \{0,1\}^n$  sampled according to BSC<sub>p</sub>  
 $\rightarrow \Pr_{\bar{e} \sim \text{BSC}_p} [\text{wt}(\bar{e}) > (p+r)n] \leq e^{-r^2 n / 2}$  for  $r > 0$

(\*) Shannon's Capacity thm for BSC<sub>p</sub> (positive part)

$\forall 0 \leq p < \frac{1}{2}, 0 \leq \epsilon \leq \frac{1}{2} - p, \exists \delta > 0$  s.t following holds for large enough  $n$ :

$\rightarrow \forall k \leq \lfloor (1 - H(p+\epsilon))n \rfloor, \exists E: \{0,1\}^k \rightarrow \{0,1\}^n$   
 $D: \{0,1\}^n \rightarrow \{0,1\}^k$

s.t  $\forall \bar{m} \in \{0,1\}^k, \text{ERR}^E(\bar{m}) \stackrel{\text{def}}{=} \Pr_{\bar{e} \sim \text{BSC}_p} [D(E(\bar{m}) + \bar{e}) \neq \bar{m}] \leq 2^{-\delta n}$

(\*) Pf. sketch:  $\rightarrow$  Pick  $E$  at random  
 $\rightarrow D = \text{MLD}_E$

$\rightarrow$  Step 1: Show for fixed  $\bar{m} \in \{0,1\}^k, \mathbb{E}_E (\text{ERR}^E(\bar{m})) \leq 2^{-\delta(n)}$

Step 2: Drop half of messages to show  $\exists E$  s.t  $\forall \bar{m} \in \{0,1\}^k, \text{ERR}^E(\bar{m}) \leq 2^{-\delta n}$

PROBABILITY FACTS: (i) Linearity of expectation (ii) Union Bound

(iii) For indep r.v.  $\mathbb{E}(X \cdot Y) = \mathbb{E}(X) \cdot \mathbb{E}(Y)$

(iv)  $\Pr(X=x) = \sum_y \Pr(X=x | Y=y) \cdot \Pr(Y=y)$

PROOF READERS for today: Sai Kiam, Sai Neeraj, Vignesh

Plan for today: Prove +ve part of Shannon's capacity thm for BSC<sub>p</sub>

Pf of part (i) of Shannon's Thm  $\bar{e} = (e_1, \dots, e_n) \sim \text{BSC}_p$

fix  $\bar{m} \in \{0, 1\}^k$

$$\bar{y} = E(\bar{m}) + \bar{e}$$

$$D = \text{MLD}_E$$

$$\text{ERR}^E(\bar{m}) \stackrel{\text{def}}{=} \Pr_{\bar{e} \sim \text{BSC}_p} [D(\bar{y}) \neq \bar{m}]$$

$$\mathbb{1}_p = \begin{cases} 1 & \text{if } p \leq \frac{1}{2} \\ 0 & \text{else} \end{cases}$$

$$= \sum_{\bar{y} \in \{0, 1\}^n} \Pr[\bar{y} | E(\bar{m})] \cdot \mathbb{1}_{D(\bar{y}) \neq \bar{m}}$$

↑ was received
↑ was transmitted

$$= \sum_{\bar{y} \in B(E(\bar{m}), (\phi + \sigma)n)} \Pr[\bar{y} | E(\bar{m})] \cdot \mathbb{1}_{D(\bar{y}) \neq \bar{m}}$$

$$\bar{y} \in B(E(\bar{m}), (\phi + \sigma)n)$$

$$\equiv \bar{y}: \Delta(E(\bar{m}), \bar{y}) \leq (\phi + \sigma)n$$

$$+ \sum_{\substack{\bar{y} \notin B(E(\bar{m}), (\phi + \sigma)n) \\ \equiv \bar{e} = \bar{y} - E(\bar{m}) \text{ has} \\ \text{wt}(\bar{e}) > (\phi + \sigma)n}} \Pr[\bar{y} | E(\bar{m})] \cdot \mathbb{1}_{D(\bar{y}) \neq \bar{m}} \leq 1$$

$$\text{ERR}^E(\bar{m}) \leq e^{-r^2 n / 2} + \mathbb{B}$$

$$+ \sum_{\bar{y} \in B(E(\bar{m}), (\phi + \sigma)n)} \Pr[\bar{y} | E(\bar{m})] \cdot \mathbb{1}_{D(\bar{y}) \neq \bar{m}}$$

$$\Rightarrow \mathbb{E}(\text{ERR}^E(\bar{m})) \leq e^{-r^2 n / 2}$$

$$+ \mathbb{E} \left( \sum_{\bar{y} \in B} \Pr[\bar{y} | E(\bar{m})] \cdot \mathbb{1}_{D(\bar{y}) \neq \bar{m}} \right) \leq e^{-r^2 n / 2}$$

$$= e^{-r^2 n / 2}$$

lim.  $\rightarrow$

$$+ \sum_{\bar{y} \in B} \mathbb{E} \left( \Pr[\bar{y} | E(\bar{m})] \cdot \mathbb{1}_{D(\bar{y}) \neq \bar{m}} \right)$$

constant under  $\mathbb{E}$  { only depends on randomness in BSC

Depends only on  $E + \text{BSC}_p$

$$= e^{-r^2 n / 2}$$

$$+ \sum_{\bar{y} \in B} \Pr[\bar{y} | E(\bar{m})] \mathbb{E} \left[ \mathbb{1}_{D(\bar{y}) \neq \bar{m}} \right]$$

Goal: bound this



Claim 1:  $\mathbb{E}_E \left[ \mathbb{1}_{D(y) \neq \bar{m}} \right] \leq 2^{-n(H(p_{TE}) - H(p_{Tr}))}$

$\Rightarrow \mathbb{E}_E \left( \text{ERR}^E(\bar{m}) \right) \leq e^{-\delta^2 n / 2} + \sum_{y \in B} \text{Pr}[y | E(\bar{m})] \cdot 2^{-n(H(p_{TE}) - H(p_{Tr}))}$

$= e^{-\delta^2 n / 2} + 2^{-n(H(p_{TE}) - H(p_{Tr}))} \sum_{y \in B} \text{Pr}[y | E(\bar{m})]$

$\leq e^{-\delta^2 n / 2} + 2^{-n(H(p_{TE}) - H(p_{Tr}))} \leq \sum_{y \in \{0,1\}^n} \text{Pr}[y | E(\bar{m})]$   
 if  $r = \frac{\epsilon}{2} > 0$   
 $\leq 2^{-\delta' n}$ , where  $\delta' = \Theta(\epsilon^2)$   
 $= \sum_{\bar{e} \in \{0,1\}^n} \text{Pr}[\bar{e} \text{ is error pattern}] = 1$

$\Rightarrow$  For every fixed  $\bar{m} \in \{0,1\}^k$   
 $\mathbb{E}_E \left( \text{ERR}^E(\bar{m}) \right) \leq 2^{-\delta' n}$

By prob method  $\Rightarrow \forall \bar{m} \in \{0,1\}^k, \exists E$  s.t.  $\text{ERR}^E(\bar{m}) \leq 2^{-\delta' n}$   
 Goal show:  $\exists E, \forall \bar{m}, \dots$

$\Rightarrow \mathbb{E}_{\bar{m} \in \{0,1\}^k} \left( \mathbb{E}_E \left( \text{ERR}^E(\bar{m}) \right) \right) \leq 2^{-\delta' n}$

Uniform  $\mathbb{E}_E \left( \mathbb{E}_{\bar{m} \in \{0,1\}^k} \text{ERR}^E(\bar{m}) \right) \leq 2^{-\delta' n}$

since choices for  $E$  &  $\bar{m}$  are indep.

By prob method  $\exists E$  s.t.  $\mathbb{E}_{\bar{m} \in \{0,1\}^k} \text{ERR}^E(\bar{m}) \leq 2^{-\delta' n}$   
 fix such an  $E$

Average error  $\leq 2^{-\delta' n} \Rightarrow$  max error  $\leq 2 \cdot 2^{-\delta' n}$   
 by dropping strategically half of message

Claim 2: Let  $\{0,1\}^{2^k} = \bar{m}_1, \dots, \bar{m}_{2^k}$

$$P_i = \mathbb{E}_{ERR}^E(\bar{m}_i)$$

$$P_1 \leq P_2 \leq \dots$$

$$\leq P_{2^k}$$

$$\frac{1}{2^k} \sum_{i=1}^{2^k} P_i \leq 2^{-\delta n}$$

$$\Rightarrow P_{\frac{2^k}{2}} \leq 2 \cdot 2^{-\delta n}$$

$\Rightarrow$  final code just use messages  $\bar{m}_1, \dots, \bar{m}_{\frac{2^k}{2}}$

$$\Rightarrow \text{max dec err} \leq 2 \cdot 2^{-\delta n} = 2^{-\delta n}$$

$k \rightarrow k-1 \rightarrow$  change in rate is negligible

$\Rightarrow$  DONE!

$$\delta = \delta' - \frac{1}{n}$$