

Apr 7

1 min checkin

~~RECAP~~ REMINDERS

- (o) HW 5 due by 11:59 pm Wed
- (o) Video due in ~ 5.5 weeks (last Sunday of semester)
 ↳ NO extension to the deadline (please plan in advance)

RECAP

- (o) $C \subseteq \Sigma^n$ is (p, L) -l.d. if $\forall \bar{y} \in \Sigma^n$,
 $|\{ \bar{c} \in C \mid \Delta(\bar{c}, \bar{y}) \leq pn \}| \leq L$
- (o) code with rel. dist δ is $(\frac{\delta}{2}, 1)$ -l.d.

Plan for today (o) Johnson bound

Q1: Can we correct $> \frac{\delta}{2}$ errors with $\text{poly}(n)$ list size?

Q2: What is the optimal tradeoff b/w R & p $(p, \text{poly}(n))$ -l.d. code?

A1: Yes! \exists codes! (Johnson bound)

THM (Johnson bound) $q \geq 2, C \subseteq [q]^n$ with dist. d

If $p \leq J_q(\frac{d}{n}) \Rightarrow C$ is (p, qdn) -l.d.
is $\text{poly}(n)$ if q is $\text{poly}(n)$

$$J_q(x) = \left(1 - \frac{1}{q}\right) \left(1 - \sqrt{1 - \frac{qx}{q-1}}\right)$$

$$0 \leq x \leq 1 - \frac{1}{q}$$

Lemma 1

$$\forall q \geq 2, 0 \leq x \leq 1 - \frac{1}{q}$$

- (pf. in book)

$$J_q(x) \geq 1 - \sqrt{1-x} > \frac{x}{2}$$

$$\Rightarrow J_q(\frac{\delta}{2}) > \frac{\delta}{2}!$$

\Rightarrow Any $(n, k, d)_q$ code is $(\frac{e}{n}, qdn)$ -l.d. from e errors

$$C \text{ is MDS } \Rightarrow \delta = 1 - R \quad e \leq n - \sqrt{n(n-d)}$$

$$\Rightarrow \text{Johnson bound } p \leq 1 - \sqrt{1 - (1-R)} \equiv \frac{e}{n} < 1 - \sqrt{\frac{n(n-d)}{n^2}}$$

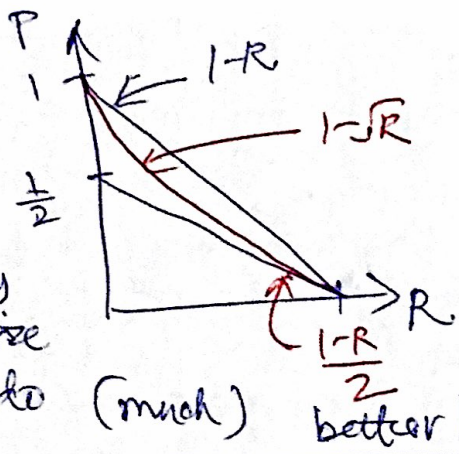
$$\text{Hamming setting (unique) rec } 1 - \sqrt{R} \quad p \leq \frac{\delta}{2} = 1 - \sqrt{1 - \frac{d}{n}}$$

Is JB tight?

Yes! \exists some code C
w/ rel. dist δ , s.t.

$P > 1 - \sqrt{1 - \delta} \Rightarrow$ super poly
code size

No: Random codes do (much) better!



Pf of Johnson bound ($q=2$)

Proof technique Double counting (to prove an inequality)

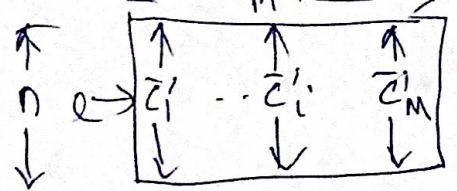
- ① Come up with some quantity S
- ② Show by (approx) counting that } *double counting*
- (i) $S \leq U$ (ii) $S \geq L$
- ③ $\Rightarrow L \leq U$
~~Star~~ (Somehow) \downarrow to prove the desired inequality
(generally algebraic manipulation)

Goal: $\forall C \subseteq \{0,1\}^n$ w/ $d = \delta n$, $\forall y \in \{0,1\}^n$
 $|B(y, \frac{\epsilon}{2}, pn) \cap C| \leq 2dn$ $P < (1 - \frac{\epsilon}{2}) (1 - \sqrt{1 - \frac{2\delta}{2-\epsilon}})$
 $\epsilon = pn$

fix $C \subseteq \{0,1\}^n$ w/ rel. dist δ
 $y \in \{0,1\}^n$ let $B(y, \frac{\epsilon}{2}, pn) \cap C = \{\bar{c}_1, \dots, \bar{c}_M\}$ for some M

Goal: $M \leq 2dn$

$\forall c \in [M]$ $\bar{c}'_i = \bar{c}_i - y$
 \Rightarrow (I) $wt(\bar{c}'_i) = \Delta(\bar{c}_i, y) \leq \epsilon$
(II) $\Delta(\bar{c}'_i, \bar{c}'_j) = \Delta(\bar{c}_i, \bar{c}_j) \geq d$



Def: $S = \sum_{(i,j)} \Delta(\bar{c}'_i, \bar{c}'_j)$
 $\bar{c}'_i[l] = c_i[l] - y[l]$
 $c'_j[l] = c_i[l] - y[l]$

(II) $\Rightarrow S = \sum_{i < j} \Delta(\bar{c}_i, \bar{c}_j) \geq \sum_{i < j} d = \binom{M}{2} d = \frac{M(M-1)}{2} \cdot d \quad \text{--- } \textcircled{\#}$

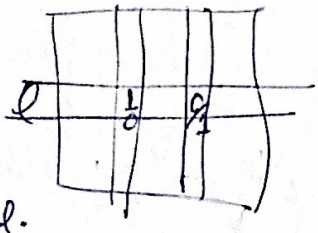
Upper bound: $k \in [n]$ ~~def~~ m_k def # of 1s in k^{th} row

$$S = \sum_{k=1}^n (\# \text{ pairs in } k^{\text{th}} \text{ row that differ})$$

$$= \sum_{k=1}^n m_k (M - m_k)$$

Goal: upper bound.

$$= M \sum_{k=1}^n m_k - \sum_{k=1}^n m_k^2$$



Aside (i) \bar{e} def $\frac{\sum_{k=1}^n m_k}{M} = \frac{\sum_{i \in [M]} \text{wt}(\bar{c}_i)}{M} \leq \frac{\sum_{i \in [M]} e}{M} = \frac{eM}{M} = e$

Aside (ii) Cauchy-Schwartz $\bar{x}, \bar{y} \in \mathbb{R}^N$

$$\langle \bar{x}, \bar{y} \rangle = \sum_{i=1}^N x_i y_i$$

$$\|\bar{x}\|^2 = \sum_{i=1}^N x_i^2$$

$$\langle \bar{x}, \bar{x} \rangle = \sum_{i=1}^N x_i^2$$

$$\langle \bar{x}, \bar{y} \rangle^2 \leq \|\bar{x}\|^2 \|\bar{y}\|^2$$

$$\equiv \|\bar{x}\|^2 \geq \langle \bar{x}, \bar{y} \rangle^2 / \|\bar{y}\|^2$$

Go back to S

$$S = M \sum_{k=1}^n m_k - \sum_{k=1}^n m_k^2$$

Cauchy-Schwartz $\Rightarrow M \cdot M \bar{e} - \frac{(\sum m_k)^2}{n}$

$$S \leq M^2 \left(\bar{e} - \frac{(\bar{e})^2}{n} \right)$$

$$M \leq 2dn$$