

The  
Dan & Bobby  
Lecture Series

Day 2

The EXP3 Algorithm

## Multi-Armed Bandit: Review.

$n$  experts/slot machines,  $T$  time steps  
at time  $t$ , choose  $x_t$  & discover  $c_t(x_t)$   
but do not know  $c_t(x)$ ,  $x \neq x_t$ .

## PSim (Phased Simulation)

$p$  phases, of length  $L = \frac{T}{p}$ .  
randomly choose  $n$  of  $L$  steps for EXPLORÉ  
→ build  $\hat{c}_j(x)$

For remaining  $L-n$  steps, sample  $X_t$  from BEX using  $\hat{c}_1, \dots, \hat{c}_t$

$$\hat{R}(\text{PSim}, T) = O\left(\sqrt[3]{\frac{n \log n}{T}}\right)$$

The EXP3 algorithm (Exponential Exploration & Exploitation).

We want to maximize the exploit steps from PSim, so we want  $L=1$ .

Let's define  $P_t[x_t] = \hat{p}_t(x)$

$$X_t = \begin{cases} \text{sample uniformly w/ Pr } \gamma \\ \text{from BEX w/ Pr } 1-\gamma. \end{cases}$$

$$\text{So, } \hat{p}_t(x) = \gamma \cdot \frac{1}{n} + (1-\gamma) p_t(x).$$

(Note, every step is a phase)

Let EXPLORE be  $\{t \mid \text{choose uniform sampling}\}$

$$\text{EXPLOIT} = [t] \setminus \text{EXPLORE.}$$

$$\text{Let } \hat{C}_t(x) = \begin{cases} C_t(x) / \hat{p}_t(x), & x = x_t \\ 0 & \end{cases}$$

$$\begin{aligned} C_t(x) < 1, \hat{p}_t(x) > \frac{\delta}{n} \\ \Rightarrow \hat{C}_t(x) < \frac{n}{\delta} \end{aligned}$$

Lemma 6  $\mathbb{E}[\hat{C}_t | \mathcal{F}_{<t}] = C_t.$

Proof:  $= \sum_x \hat{C}_t(x) \hat{p}_t(x) \stackrel{\text{def}}{=} C_t(x).$

Lemma 7.  $\mathbb{E}[\hat{C}_t(x_t) | \mathcal{F}_{<t}]$   
 $= \mathbb{E}[C_t(x_t) | \mathcal{F}_{<t}].$

Proof.  $= \sum_x \mathbb{E}[\hat{C}_t(x) | \mathcal{F}_{<t}] p_t(x) =$   
 $= \sum_x C_t(x) p_t(x) = \mathbb{E}[C_t(x_t) | \mathcal{F}_{<t}]$

Lemma 8. For all  $x$  in  $[n]$ ,

$$\frac{1}{T} \mathbb{E} \left[ \sum_{t \in \text{EXPLOIT}} (c_t(x_t) - c_t(x)) \right] < \frac{\epsilon \Delta}{\delta} + \frac{n \log n}{\epsilon \gamma T}.$$

By Lemmas 6 & 7,

$$= \frac{1}{T} \mathbb{E} \left[ \sum_{t \in \text{EXPLOIT}} (\hat{c}_t(x_t) - \hat{c}_t(x)) \right]$$

↳ multiply previous proof by factor of  $\frac{n}{\delta}$ .

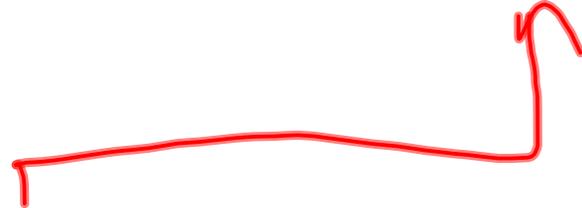
Lemma  $\uparrow$   $\frac{1}{T} \mathbb{E} \left[ \sum_{t \in \text{EXPLORE}} (c_t(v_t) - c_t(v)) \right] < \gamma.$

$$\leq \frac{1}{T} \cdot \mathbb{E} [ |\text{EXPLORE}| ]$$

$$= \frac{1}{T} \cdot \gamma T = \gamma.$$

Theorem 10. Let  $\gamma = 4 \sqrt{\frac{n \log n}{T}}$  and  $\epsilon = \gamma^2$

Then  $\hat{R}(\text{EXP3}, T) = O(\epsilon)$ .



Pretty good, but we can do better!

Lemma 11. If  $0 < \alpha < 1$  and  $\hat{C}_t : [n] \rightarrow \mathbb{R}_+$ .

$$\sum_{t=1}^T \sum_{x \in [n]} p_t(x) \hat{C}_t(x) \leq$$

$$\max_{x \in [n]} \sum_{t=1}^T \hat{C}_t(x) - \log(1 - \alpha) \sum_{t=1}^T \sum_{x \in [n]} p_t(x) \hat{C}_t(x)^2 + \frac{\log(n)}{\alpha}.$$

Proof.

First, we want to bound  $(1-\alpha)^x$

$$(1) \quad (1-\alpha)^x \leq \underbrace{1 + \log(1-\alpha)x + \log^2(1-\alpha)x^2}_{g'(x)}$$

$\parallel$   
 $f(x)$

$$f(0) = g(0) = 1$$

$$f'(0) = g'(0) = \log(1-\alpha)$$

$$f''(0) = \log^2(1-\alpha) < 2\log^2(1-\alpha) = g''(0)$$

$$f''(x) < f''(0) < g''(0) = g''(x).$$

Define  $h(x) = g(x) - f(x)$ .

$$h(0) = h'(0) = 0$$

$$h''(x) > 0 \Rightarrow \text{for all } x, f(x) \leq g(x),$$

Hedge says:

$$w_t(x) = (1-\alpha) \sum_{s=1}^t \hat{c}_s(x)$$

↳

$$W_t = \sum_x w_t(x)$$

$$P_t(x) = \frac{w_{t-1}(x)}{W_{t-1}}$$

$$\begin{aligned}
\frac{w_t}{w_{t-1}} &= \sum_x P_t(x) (1-\alpha)^{\hat{c}_t(x)} \\
&= \sum_x \frac{w_{t-1}(x)}{w_{t-1}} (1-\alpha)^{\hat{c}_t(x)} \\
&= \frac{1}{w_{t-1}} \sum_x (1-\alpha)^{\sum_{s=1}^t \hat{c}_s(x)} \cdot \left( \frac{w_t}{w_{t-1}} \right) \\
&= \frac{w_t}{w_{t-1}}
\end{aligned}$$

$$\frac{w_t}{w_{t-1}} = \sum_x P_{\mathcal{L}}(x) (1-\alpha)^{\hat{\mathcal{L}}_{\mathcal{L}}(x)}$$

$$\log(1+b) \leq b \quad \text{if } b > -1$$

$$\leq |+\log(1-\alpha)| \sum_x P_{\mathcal{L}}(x) \hat{\mathcal{L}}_{\mathcal{L}}(x) + \log^2(1-\alpha) \sum_x P_{\mathcal{L}}(x) \hat{\mathcal{L}}_{\mathcal{L}}(x)^2$$

$$|\log(1-\alpha) \hat{\mathcal{L}}_{\mathcal{L}}(x)| < 1 \quad |\log(1-\alpha) \mathcal{L}_{\mathcal{L}}(x)| > 1$$

$$\log(w_t) - \log(w_{t-1}) \leq \log(1-\alpha) \sum_x P_{\mathcal{L}}(x) \hat{\mathcal{L}}_{\mathcal{L}}(x) + \log^2(1-\alpha) \sum_x P_{\mathcal{L}}(x) \hat{\mathcal{L}}_{\mathcal{L}}(x)^2$$

$$\log(W_T) - \log(W_0) \leq \log(1-\alpha) \sum_{t,x} P_t(x) \Delta_t(x)$$

$$+ \log^2(1-\alpha) \sum_{t,x} P_t(x) \Delta_t^2(x)$$

Also

$$\log(W_T) \geq \log(1-\alpha) \sum_{t=1}^T \hat{C}_t(x)$$

$$W_T \geq (1-\alpha)^{\sum_{t=1}^T \hat{C}_t(x)}$$

$$0 < \alpha < 1$$

$$\frac{-1}{\log(1-\alpha)} < \frac{1}{\alpha}$$

$$-1 > \frac{\log(1-\alpha)}{\alpha}$$

Lemma 12 If  $0 < \gamma \leq \frac{1}{2}$

AND hedge  $(\delta_n)$  is used

$$\Rightarrow \alpha = \frac{\delta}{n}$$

$$\frac{1}{T} \mathbb{E} \left[ \sum_{t \in \mathbb{F}_{\text{exploit}}} \zeta_t(x_t) - c_t(x_t) \right] \leq (4 \ln 2) \gamma T + \frac{h \ln n}{\delta}$$

$$\frac{1}{T} \# \left[ \sum_{\lambda \in \Lambda} c_{\lambda}(\lambda) - c_{\lambda}(\lambda) \right] \leq$$

$$- \log_2 \left( 1 - \frac{\delta}{n} \right) \# \left[ \sum_{\lambda \in \Lambda} P_{\lambda}(\lambda) \hat{c}_{\lambda}^2(\lambda) \right] + \frac{n \log n}{\delta}$$

$$\alpha \frac{\delta}{n} \leq \frac{1}{2} \Rightarrow -\log_2 \left( 1 - \frac{\delta}{n} \right) \leq [2 \ln(2)] \frac{\delta}{n}$$

$$P_{\lambda}(\lambda) = \frac{1}{1-\delta} \left( \hat{P}_{\lambda}(\lambda) - \frac{\delta}{n} \right) \leq 2 \hat{P}_{\lambda}(\lambda)$$

$$\text{So } P_{\lambda}(\lambda) \hat{c}_{\lambda}(\lambda) \leq 2 \hat{P}_{\lambda}(\lambda) \hat{c}_{\lambda}(\lambda) \leq 2$$

$$-k_B \left(1 - \frac{\delta}{n}\right) \# \left[ \sum_{x,z} P_z(x) \Delta_z^2(x) \right]$$

$$< (4 \ln 2) \frac{\delta}{n} \sum_{x,z} \# [\Delta_z(x)]$$

$$= (4 \ln 2) \frac{\delta}{n} \sum_{x,z} \leftarrow z(x)$$

$$\leq 4 \ln 2 \frac{\delta}{n} (nT)$$

$$= (4 \ln 2) \delta T$$

THM 13

Regret of Exp 3 is  $O(\sqrt{T \ln \ln n})$

lets pick  $\gamma = \sqrt{\frac{n \ln n}{T}}$

then we see thm 13 is true

Example / Demos

or Something

(For those of you  
who cared enough  
to show up)