Introduction

- An important and fundamental feature in modern operating systems is concurrent execution of processes/threads. This feature is essential for the realization of multiprogramming, multiprocessing, distributed systems, and client-server model of computation.
- Concurrency encompasses many design issues including communication and synchronization among processes, sharing of and contention for resources.
- In this discussion we will look at the various design issues/problems and the wide variety of solutions available.

Topics for discussion

- The principles of concurrency
- Interactions among processes
- Mutual exclusion problem
- Mutual exclusion- solutions
  - Software approaches (Dekker's and Peterson's)
  - Hardware support (test and set atomic operation)
  - OS solution (semaphores)
  - PL solution (monitors)
  - Distributed OS solution (message passing)
- Reader/writer problem
- Dining Philosophers Problem

Principles of Concurrency

- Interleaving and overlapping the execution of processes.
- Consider two processes P1 and P2 executing the function echo:
  ```c
  { 
    input (in, keyboard);
    out = in;
    output (out, display);
  }
  ```

Interactions among processes

- In a multi-process application these are the various degrees of interaction:
  1. **Competing processes**: Processes themselves do not share anything. But OS has to share the system resources among these processes "competing" for system resources such as disk, file or printer.
  2. **Co-operating processes**: Results of one or more processes may be needed for another process.
  3. **Co-operation by sharing**: Example: Sharing of an IO buffer. Concept of critical section, (indirect)
  4. **Co-operation by communication**: Example: typically no data sharing, but co-ordination thru’ synchronization becomes essential in certain applications. (direct)
Among the three kinds of interactions indicated by 1, 2 and 3 above:
1 is at the system level: potential problems: deadlock and starvation.
2 is at the process level: significant problem is in realizing mutual exclusion.
3 is more a synchronization problem.
We will study mutual exclusion and synchronization here, and defer deadlock, and starvation for a later time.

Race condition: The situation where several processes access and manipulate shared data concurrently. The final value of the shared data depends upon which process finishes last.
To prevent race conditions, concurrent processes must be synchronized.

Successful use of concurrency among processes requires the ability to define critical sections and enforce mutual exclusion.
Critical section: is that part of the process code that affects the shared resource.
Mutual exclusion: in the use of a shared resource is provided by making its access mutually exclusive among the processes that share the resource.
This is also known as the Critical Section (CS) problem.

Any facility that provides mutual exclusion should meet these requirements:
1. No assumption regarding the relative speeds of the processes.
2. A process is in its CS for a finite time only.
3. Only one process allowed in the CS.
4. Process requesting access to CS should not wait indefinitely.
5. A process waiting to enter CS cannot be blocking a process in CS or any other processes.

Software Solutions: Algorithm 1

<table>
<thead>
<tr>
<th>Process 0</th>
</tr>
</thead>
<tbody>
<tr>
<td>...</td>
</tr>
<tr>
<td>while turn != 0 do</td>
</tr>
<tr>
<td>nothing;</td>
</tr>
<tr>
<td>// busy waiting</td>
</tr>
<tr>
<td>&lt; Critical Section&gt;</td>
</tr>
<tr>
<td>turn = 1;</td>
</tr>
<tr>
<td>...</td>
</tr>
</tbody>
</table>

Problems: Strict alternation, Busy Waiting

Algorithm 2

<table>
<thead>
<tr>
<th>PROCESS 0</th>
</tr>
</thead>
<tbody>
<tr>
<td>...</td>
</tr>
<tr>
<td>flag[0] = TRUE;</td>
</tr>
<tr>
<td>while flag[1] do</td>
</tr>
<tr>
<td>nothing;</td>
</tr>
<tr>
<td>&lt;CRITICAL SECTION&gt;</td>
</tr>
<tr>
<td>flag[0] = FALSE;</td>
</tr>
<tr>
<td>turn = 0;</td>
</tr>
<tr>
<td>...</td>
</tr>
</tbody>
</table>

Problems: Potential for deadlock, if one of the processes fails within CS.
Algorithm 3
Combined shared variables of algorithms 1 and 2.

Process \( P_i \)
do {
    flag [i] := true;
    turn = j;
    while (flag[j] and turn = j) ;
    critical section
    flag [i] = false;
    remainder section
} while (1);
Solves the critical-section problem for two processes.

Synchronization Hardware
Test and modify the content of a word atomically.

```java
boolean TestAndSet(boolean &target) {
    boolean rv = target;
    target = true;
    return rv;
}
```

Mutual Exclusion with Test-and-Set
Shared data:
boolean lock = false;

Process \( P_i \)
do {
    while (TestAndSet(lock)) ;
    critical section
    lock = false;
    remainder section
}

Synchronization Hardware
Atomically swap two variables.

```java
void Swap(boolean &a, boolean &b) {
    boolean temp = a;
    a = b;
    b = temp;
}
```

Mutual Exclusion with Swap
Shared data (initialized to false):
boolean lock;

Process \( P_i \)
do {
    key = true;
    while (key = true) 
    Swap(lock,key);
    critical section
    lock = false;
    remainder section
}

Semaphores
Think about a semaphore as a class
Attributes: semaphore value, Functions: init, wait, signal
Support provided by OS
Considered an OS resource, a limited number available: a limited number of instances (objects) of semaphore class is allowed.
Can easily implement mutual exclusion among any number of processes.
Critical Section of $n$ Processes

- Shared data:
  
  **Semaphore mutex; //initially mutex = 1**

- Process $P_i$:
  
  ```
  do {
      mutex.wait();
      critical section
      mutex.signal();
      remainder section
  } while (1);
  ```

Semaphore Implementation

- Define a semaphore as a class:
  ```
  class Semaphore {
      int value; // semaphore value
      ProcessQueue L; // process queue
      //operations
      wait()
      signal()
  }
  ```

  - In addition, two simple utility operations:
    - `block()` suspends the process that invokes it.
    - `Wakeup()` resumes the execution of a blocked process $P$.

Semantics of wait and signal

- Semaphore operations now defined as:
  ```
  S.wait():
  S.value--;
  if (S.value < 0) {
      add this process to S.L;
      block(); // block a process
  }

  S.signal():
  S.value++;
  if (S.value <= 0) {
      remove a process $P$ from S.L;
      wakeup(); // wake a process
  }
  ```

Semaphores for CS

- Semaphore is initialized to 1. The first process that executes a `wait()` will be able to immediately enter the critical section (CS). ($S.wait()$ makes $S$ value zero.)

- Now other processes wanting to enter the CS will each execute the `wait()` thus decrementing the value of $S$, and will get blocked on $S$. (If at any time value of $S$ is negative, its absolute value gives the number of processes waiting blocked.)

- When a process in CS departs, it executes $S.signal()$ which increments the value of $S$, and will wake up any one of the processes blocked. The queue could be FIFO or priority queue.

Two Types of Semaphores

- **Counting** semaphore – integer value can range over an unrestricted domain.

- **Binary** semaphore – integer value can range only between 0 and 1; can be simpler to implement. ex: nachos

- Can implement a counting semaphore using a binary semaphore.

Semaphore for Synchronization

- Execute $B$ in $P_b$ only after $A$ executed in $P_a$

  ```
  Pi
  A
  flag.wait()
  flag.signal()
  Pb
  ```

- Use semaphore `flag` initialized to 0

- Code:
  ```
  Pb
  A
  flag.wait()
  flag.signal()
  ```
Classical Problems of Synchronization

- Bounded-Buffer Problem
- Readers and Writers Problem
- Dining-Philosophers Problem

Producer/Consumer problem

- Producer
  repeat
  produce item v;
  b[in] = v;
  in = in + 1;
  MUTEX.signal();
  forever;

- Consumer
  repeat
  while (in <= out) nop;
  MUTEX.wait();
  w = b[out];
  out = out + 1;
  MUTEX.signal();
  consume w;
  forever;

Solution for P/C using Semaphores

- Producer
  repeat
  produce item v;
  MUTEX.wait();
  b[in] = v;
  in = in + 1;
  MUTEX.signal();
  forever;

- Consumer
  repeat
  while (in <= out) nop;
  MUTEX.wait();
  w = b[out];
  out = out + 1;
  MUTEX.signal();
  consume w;
  forever;

What if Producer is slow or late?

Ans: Consumer will busy-wait at the while statement.

P/C: improved solution

- Producer
  repeat
  produce item v;
  MUTEX.wait();
  b[in] = v;
  in = in + 1;
  MUTEX.signal();
  AVAIL.signal();
  forever;

- Consumer
  repeat
  AVAIL.wait();
  MUTEX.wait();
  w = b[out];
  out = out + 1;
  MUTEX.signal();
  consume w;
  forever;

What will be the initial values of MUTEX and AVAIL?

Ans: Initially MUTEX = 1, AVAIL = 0.

P/C problem: Bounded buffer

- Producer
  repeat
  produce item v;
  while((in+1)%n == out)
  NOP;
  b[in] = v;
  in = (in + 1)%n;
  forever;

- Consumer
  repeat
  while (in == out) NOP;
  w = b[out];
  out = (out + 1)%n;
  consume w;
  forever;

ANS: Using another counting semaphore.

P/C: Bounded Buffer solution

- Producer
  repeat
  produce item v;
  BUFSIZE.wait();
  MUTEX.wait();
  b[in] = v;
  in = (in + 1)%n;
  MUTEX.signal();
  AVAIL.signal();
  forever;

- Consumer
  repeat
  AVAIL.wait();
  MUTEX.wait();
  w = b[out];
  out = (out + 1)%n;
  MUTEX.signal();
  consume w;
  forever;

What is the initial value of BUFSIZE?

ANS: size of the bounded buffer.
Semaphores - comments

- Intuitively easy to use.
- `wait()` and `signal()` are to be implemented as atomic operations.
- **Difficulties:**
  - `signal()` and `wait()` may be exchanged inadvertently by the programmer. This may result in deadlock or violation of mutual exclusion.
  - `signal()` and `wait()` may be left out.
  - Related `wait()` and `signal()` may be scattered all over the code among the processes.

Monitors

- Monitor is a predecessor of the "class" concept.
- Initially it was implemented as a programming language construct and more recently as library. The latter made the monitor facility available for general use with any PL.
- Monitor consists of procedures, initialization sequences, and local data. Local data is accessible only thru’ monitor’s procedures. Only one process can be executing in a monitor at a time. Other process that need the monitor wait suspended.

```
monitor monitor-name
{
  shared variable declarations
  procedure body P1 (...) {
    . . .}
  procedure body P2 (...) {
    . . .}
  procedure body Pn (...) {
    . . .}
  {
    initialization code
  }
}
```

To allow a process to wait within the monitor, a condition variable must be declared, as

```
condition x, y;
```

Condition variable can only be used with the operations `wait` and `signal`.

- The operation `x.wait();` means that the process invoking this operation is suspended until another process invokes `x.signal();`
- The `x.signal` operation resumes exactly one suspended process. If no process is suspended, then the `signal` operation has no effect.

Schematic View of a Monitor

Monitor With Condition Variables
Message passing
- Both synchronization and communication requirements are taken care of by this mechanism.
- More over, this mechanism yields to synchronization methods among distributed processes.
- Basic primitives are:
  send (destination, message);
  receive (source, message);

Issues in message passing
- Send and receive: could be blocking or non-blocking:
  - Blocking send: when a process sends a message it blocks until the message is received at the destination.
  - Non-blocking send: After sending a message the sender proceeds with its processing without waiting for it to reach the destination.
- Blocking receive: When a process executes a receive it waits blocked until the receive is completed and the required message is received.
- Non-blocking receive: The process executing the receive proceeds without waiting for the message(!).
  - Blocking Receive/non-blocking send is a common combination.

Reader/Writer problem
- Data is shared among a number of processes.
- Any number of reader processes could be accessing the shared data concurrently.
- But when a writer process wants to access, only that process must be accessing the shared data. No reader should be present.
- Solution 1: Readers have priority; If a reader is in CS any number of readers could enter irrespective of any writer waiting to enter CS.
- Solution 2: If a writer wants CS as soon as the CS is available writer enters it.

Reader/writer: Priority
Readers
- Writer:
  ForCS.wait();
  CS;
  ForCS.signal();
- Reader:
  ES.wait();
  NumRdr = NumRdr + 1;
  if NumRdr = 1 ForCS.wait();
  ES.signal();
  CS;
  ES.wait();
  NumRdr = NumRdr -1;
  If NumRdr = 0 ForCS.signal();
  ES.signal();

Dining Philosophers Example
```c
monitor dp
{
    enum {thinking, hungry, eating}
    state[5];
    condition self[5];
    void pickup(int i) // following slides
    void putdown(int i) // following slides
    void test(int i) // following slides
    void init() {
        for (int i = 0; i < 5; i++)
            state[i] = thinking;
    }
}
```
void test(int i) {
    if ((state[(i + 4) % 5] != eating) &&
            (state[i] == hungry) &&
            (state[(i + 1) % 5] != eating))
    {
        state[i] = eating;
        self[i].signal();
    }
}