

Open book, open notes, closed neighbors, 170 minutes after a 10-minute read-through period. The exam totals 267 pts., subdivided as shown. Do *all seven problems* on these exam sheets. Always show your work—this may help for partial credit. Use of a quantum circuit simulation app—whether online or downloaded—or matrix calculator is forbidden. You may use—and cite—theorems and facts from lectures and homeworks without further justification.

*Notation:* You may freely mix “standard linear algebra notation” and “Dirac *bra-ket* notation” for quantum states. For example,  $e_{10}$  means the same thing as  $|10\rangle$  and is represented (in big-endian notation) by the unit column vector  $[0, 0, 1, 0]^T$ . Non-unit vectors should only use the standard notation. For matrix outer-products the *ket-bra* form typified by  $|+\rangle\langle-|$  is standard. The notation  $\langle u | v \rangle$  is the same as  $\langle u, v \rangle$  for inner product. Quantum state vectors use big-endian notation by default; if you switch to little-endian you must say so.

[Your actual final exam will use exam booklets. Its formatting will be close to this but will differ somewhat.]

**(1) ( $3 \times 15 = 45$  pts.)** *Short-answer questions.* A single pithy sentence may suffice; general expectation is a paragraph of 2–4 sentences. (This is technically the same as **true/false with justifications** but with longer questions.)

- (a) If  $A$  and  $B$  are unitary matrices, must  $\frac{1}{\sqrt{2}}(A + B)$  be unitary? If you say yes, prove it; if you say no, give a concrete counterexample.
- (b) Suppose you know in advance that a Grover search problem has exactly one solution  $y \in \{0, 1\}^n$ . The Grover oracle reflects about the “miss vector”  $\mathbf{m}_{\{y\}}$ , and the “hit vector”  $\mathbf{h}_{\{y\}}$  is a simple linear combination of  $\mathbf{m}_{\{y\}}$  and the all-1s vector  $\mathbf{j}$  (which is easily computed via  $\mathbf{H}^{\otimes n} |0^n\rangle$ ). Explain as best you can why Grover’s algorithm cannot simply jump to  $\mathbf{h}_{\{y\}}$  in one iteration step and then measure to give  $y$  with virtually 100% probability.
- (c) Suppose  $M = pq$  is a product of two odd primes and  $a$  is relatively prime to  $M$  with  $1 < a < M$ . Then  $z = a^2 \pmod{M}$  is a quadratic residue modulo  $M$ . Find a different number  $b < M$  such that  $b^2 \pmod{M} = z$ , show this, and say why  $b \neq a$ .

**(2) ( $9+6+15 = 30$  pts. total)**

Make a quantum circuit of two qubits—both initialized to  $|0\rangle$ —with the following gates:

1. A Hadamard on line 1 only.
2. A CNOT gate with control on line 1 and target on line 2.
3. An **S** gate on line 1. (This is also called the *phase gate* and has matrix  $\begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix}$ .)
4. Another CNOT gate with control on line 1 and target on line 2.

- (a) Calculate the resulting two-qubit state  $\phi$ .
- (b) Is  $\phi$  entangled? Prove your answer.
- (c) Now replace the second CNOT gate by a **CZ** gate on the same two qubits. Calculate the final state  $\phi'$  now and again say whether it is entangled.

**(3) (12 + 6 + 9 + 9 + 6 = 42 pts. total)**

Make a quantum circuit  $C$  on three qubits by placing the following gates:

1. Hadamard gates on all three lines.
2. A CZ gate between lines 1 and 2.
3. A CNOT gate with control on line 1 and target on line 3.
4. A CZ gate between lines 2 and 3.
5. Hadamard gates on all three lines again.

This differs from a graph-state circuit only in having the CNOT gate.

- (a) Compute  $\langle 0^n | C | 0^n \rangle$  *without* resorting to  $8 \times 8$  matrix multiplications. Using a “maze diagram” is fine, but there are methods even shorter than that.
- (b) Compute the state  $\phi$  that was present before the final three Hadamard gates—that is, after the second **CZ** gate. (This might already be part of your answer to part (a).)
- (c) Compute the density matrix  $\rho = |\phi\rangle \langle \phi|$ . (OK, this is tedious—but you get 9 pts. for it.)
- (d) Trace out qubit 3 from  $\rho$ .
- (e) Does the resulting 2-qubit density matrix  $\rho'$  represent a pure state? a completely mixed state? a mixed state but not completely mixed? Justify your answer.

**(4) (36 pts. total)**

Let  $A$  be the Hermitian PSD matrix  $\begin{bmatrix} 2 & 1 \\ 1 & 2 \end{bmatrix}$ .

- (a) Find its spectral decomposition  $A = \lambda_1 |u_1\rangle \langle u_1| + \lambda_2 |u_2\rangle \langle u_2|$ .
- (b) Calculate  $U = e^{i\pi A}$  by replacing  $\lambda_1$  with  $e^{i\pi\lambda_1}$  and ditto for  $\lambda_2$ .
- (c) Say briefly why  $U$  is unitary. What quantum operation (if any) does it represent?

**(5) (12 + 6 + 24 = 42 pts.)**

Make a quantum circuit of two qubits—both initialized to  $|0\rangle$ —by placing the following gates:

1. A Hadamard gate on line 1 only.
  2. Then a **T**-gate on line 1 only.
  3. A CNOT gate with control on line 1 and target on line 2.
  4. Another Hadamard gate on line 1.
- (a) Calculate the final quantum state  $\phi$  of two qubits.
- (b) Show that  $\phi$  is entangled.
- (c) Use the SVD truncation technique to approximate  $\phi$  by a separable state  $\alpha \otimes \beta$ .

**(6) (36 pts. total)**

Design a quantum circuit  $C$  of 3 qubits that on input  $|000\rangle$  prepares the pure state

$$\frac{1}{2}(|001\rangle - |011\rangle - |100\rangle + |111\rangle).$$

**(7) (6 × 6 = 36 pts. total)** *Multiple choice*—no justifications are required but may help for partial credit:

1. If a one-qubit quantum state  $|\phi\rangle$  is measured in an orthonormal basis  $B_1 = \{|u_1\rangle, |v_1\rangle\}$  versus another orthonormal basis  $B_2 = \{|u_2\rangle, |v_2\rangle\}$ , then:
  - (a) The probability of outcome  $|u_1\rangle$  using  $B_1$  will be the same as that of outcome  $|u_2\rangle$  using  $B_2$ .
  - (b) The probability of outcome  $|u_1\rangle$  using  $B_1$  will either be the same as that of the outcome  $|u_2\rangle$  using  $B_2$  or the outcome  $|v_2\rangle$  using  $B_2$ , because the basis could be listed either way.
  - (c) One could have the probability of outcomes  $|u_1\rangle$  versus  $|v_1\rangle$  be 50-50 while the probability of outcomes  $|u_2\rangle$  versus  $|v_2\rangle$  could be 75% versus 25%.
  - (d) It is possible for both the probability of getting  $|u_2\rangle$  and the probability of getting  $|v_2\rangle$  to be zero when the basis  $B_2$  is used.
2. A three-qubit quantum circuit  $C$  with the property that for any single-qubit pure state  $|\phi\rangle$ ,  $C$  on input  $|\phi\rangle \otimes |0\rangle \otimes |0\rangle$  outputs  $|0\rangle \otimes |\phi\rangle \otimes |\phi\rangle$  is:

- (a) Built simply by placing Hadamard on line 1, then two CNOT gates with controls on line 1 and targets on lines 2 and 3, respectively, and finally another Hadamard gate on line 1.
  - (b) Buildable by modifying the circuit in the text for “teleporting” an arbitrary qubit  $|\phi\rangle$ .
  - (c) Impossible, as shown by the No-Cloning Theorem.
  - (d) Impossible, by Holevo’s theorem that  $n$  qubits can yield only  $n$  bits of classical information.
3. A unitary  $4 \times 4$  matrix:
- (a) Always has positive real eigenvalues.
  - (b) May have 0 as an eigenvalue.
  - (c) Has only eigenvalues of the form  $e^{i\theta}$  for some angle  $\theta$ ,  $0 \leq \theta < 2\pi$ .
  - (d) Has eigenvalues whose squares sum to 1.
4. A Hermitian  $4 \times 4$  matrix:
- (a) Always has positive real eigenvalues.
  - (b) May have 0 as an eigenvalue.
  - (c) Has only eigenvalues of the form  $e^{i\theta}$  for some angle  $\theta$ ,  $0 \leq \theta < 2\pi$ .
  - (d) Has eigenvalues whose squares sum to 1.
5. The CHSH Game is significant because:
- (a) It proves that faster-than-light communication is possible.
  - (b) No physical system that operates only locally—without using non-local quantum entanglement—can achieve success probability over 75% in the long run of repeated trials.
  - (c) Experiments have achieved success probability over 80%—nearing the theoretical 85.3...%—by quantum means (winning a Nobel Prize, incidentally).
  - (d) Two of the above statements are true but the other is not true.
6. Shor’s Algorithm:
- (a) Has been used to factor integers of over 1,000 digits of the kind used in RSA encryption with consistent success.
  - (b) Solves a problem in quantum polynomial time that classical computers have been proved not to be able to solve in deterministic polynomial time.
  - (c) Challenges the assertion that natural processes can always be modeled precisely and efficiently in a computer language like C++.
  - (d) Cannot be executed using quantum circuits that have only Hadamard, CNOT, and **T**-gates.

END OF EXAM