

Prove $ALLM = \{ \langle M \rangle : L(M) = \Sigma^* \}$ is neither c.e. nor co-c.e.

Proof: By the All or Nothing Switch, $ATM \leq_m ALLM$



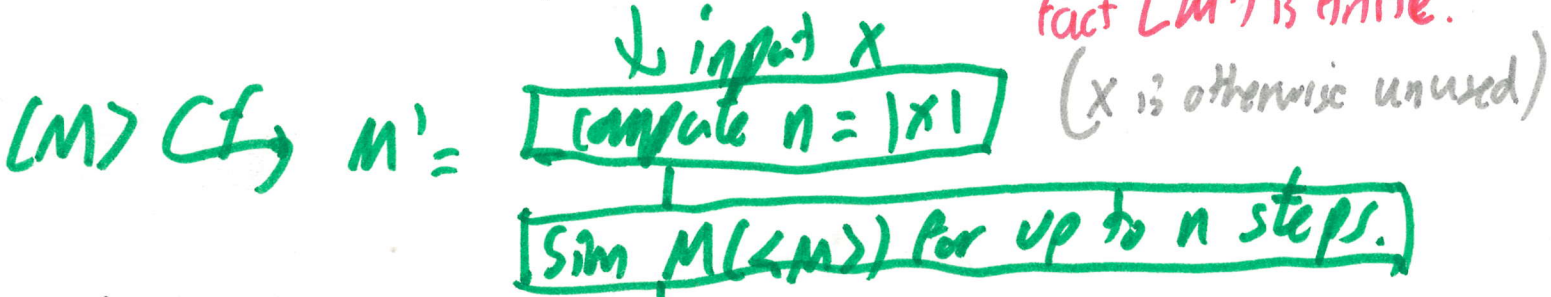
Since ATM is not co-c.e. $ALLM$ is not co-c.e. either.

Defn: A language B is hard for a class \mathcal{C} under \leq_m if for all $A \in \mathcal{C}$, $A \leq_m B$.
If also $B \in \mathcal{C}$, then B is \mathcal{C} -complete.

To show $ALLM$ is not c.e. either, show $DTM \leq_m ALLM$. I.e.

map $\langle M \rangle \mapsto M'$ such that

IF M does not accept $\langle M \rangle$, then $L(M') = \Sigma^*$
IF M does accept its own code, then $L(M') \neq \Sigma^*$; in fact $L(M')$ is finite.



Now define the problem

INST: A TM M

Ques: Is $L(M)$ infinite?

Language: $INF = \{ \langle M \rangle : L(M) \text{ is infinite} \}$.

Its complement $\{ \langle M \rangle : L(M) \text{ is finite} \}$

is called FIN or $IFIN$ because it is an index set.

Prove that INF and $IFIN$ are neither c.e. nor co-c.e.

$ATM \leq_m INF$ via $A \cup NS$
 $DTM \leq_m INF$ as above!