

CSE 421/521 - Operating Systems
Fall 2011

LECTURE - XV

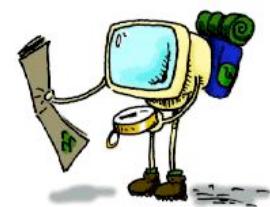
MEMORY MANAGEMENT & VIRTUAL MEMORY

Tevfik Koşar

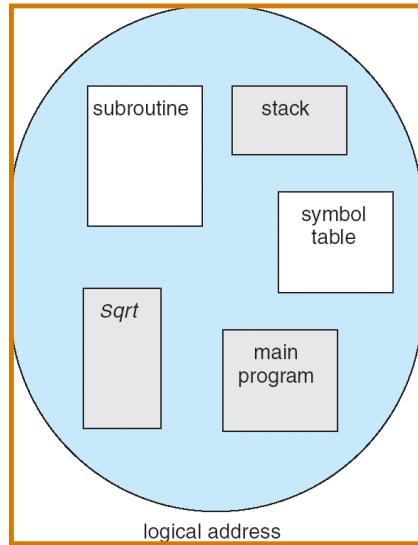
University at Buffalo
October 25th, 2011

Roadmap

- Main Memory Management
 - Segmentation
- Virtual Memory
 - Demand Paging
 - Page Faults
 - Page Replacement
 - Page Replacement Algorithms
 - Performance of Demand Paging



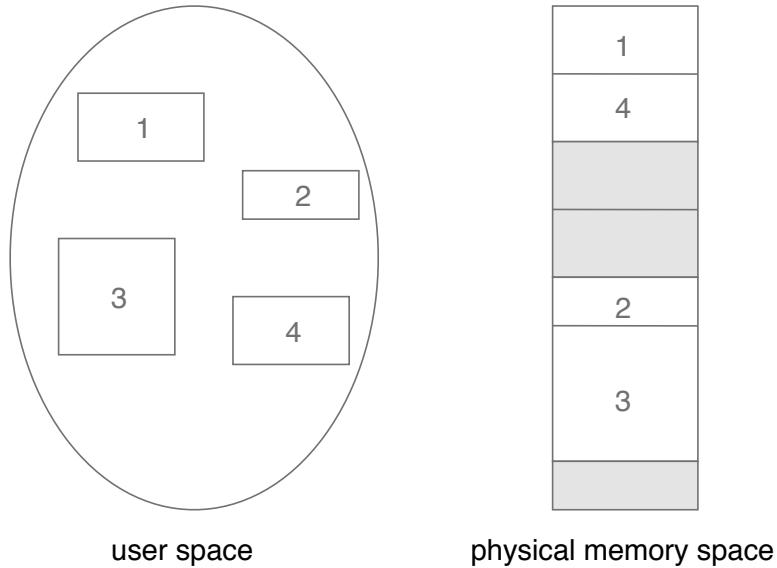
User's View of a Program



Segmentation

- Memory-management scheme that supports user view of memory
- A program is a collection of segments. A segment is a logical unit such as:
 - main program,
 - procedure,
 - function,
 - method,
 - object,
 - local variables, global variables,
 - common block,
 - stack,
 - symbol table, arrays

Logical View of Segmentation



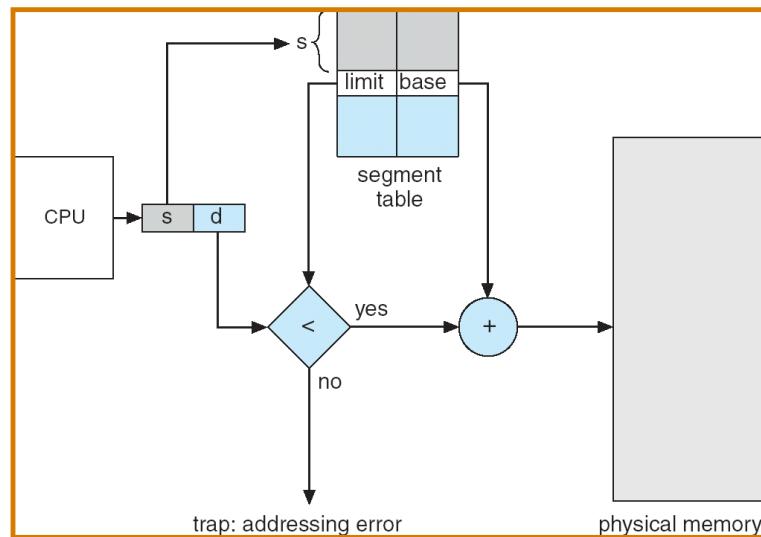
Segmentation Architecture

- Logical address consists of a two tuple:
 \langle segment-number, offset \rangle ,
- **Segment table** - maps two-dimensional physical addresses; each table entry has:
 - base - contains the starting physical address where the segments reside in memory
 - *limit* - specifies the length of the segment
- **Segment-table base register (STBR)** points to the segment table's location in memory
- **Segment-table length register (STLR)** indicates the length (limit) of the segment
- segment addressing is d (offset) $<$ STLR

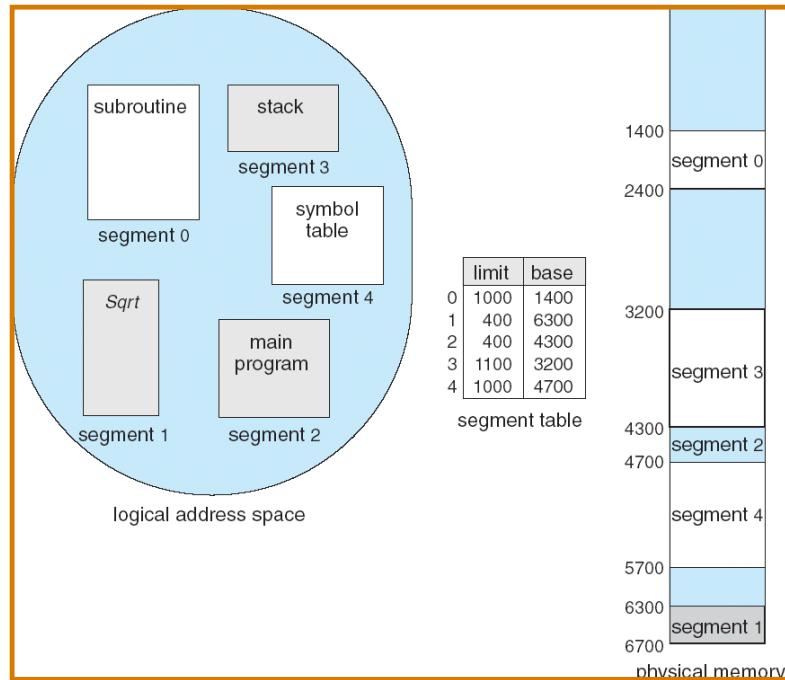
Segmentation Architecture (Cont.)

- **Protection.** With each entry in segment table associate:
 - validation bit = 0 \Rightarrow illegal segment
 - read/write/execute privileges
- Protection bits associated with segments; code sharing occurs at segment level
- Since segments vary in length, memory allocation is a dynamic storage-allocation problem
- A segmentation example is shown in the following diagram

Address Translation Architecture



Example of Segmentation



Exercise

- Consider the following segment table:

<u>Segment</u>	<u>Base</u>	<u>Length</u>
0	219	600
1	2300	14
2	90	100
3	1327	580
4	1952	96

What are the physical addresses for the following logical addresses?

a. 1, 100

b. 2, 0

c. 3, 580

Solution

- Consider the following segment table:

Segment	Base	Length
0	219	600
1	2300	14
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What are the physical addresses for the following logical addresses?

a. 1, 100

illegal reference (2300+100 is not within segment limits)

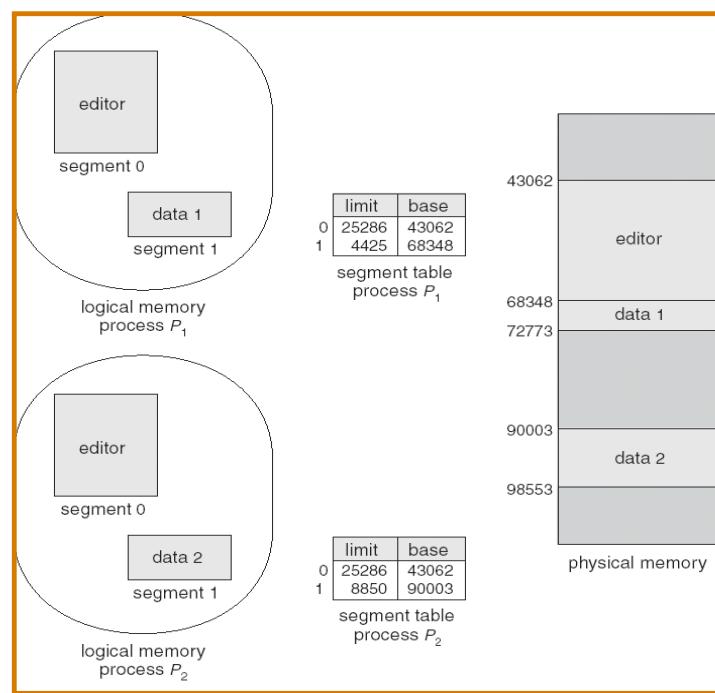
b. 2, 0

physical address = 90 + 0 = 90

c. 3, 580

illegal reference (1327 + 580 is not within segment limits)

Sharing of Segments



Virtual Memory

Background

- **Virtual memory** - separation of user logical memory from physical memory.
 - Only part of the program needs to be in memory for execution.
 - Logical address space can therefore be much larger than physical address space.
 - Allows address spaces to be shared by several processes.
 - Allows for more efficient process creation.
- Virtual memory can be implemented via:
 - Demand paging
 - Demand segmentation

Demand Paging

- Bring a page into memory only when it is needed
 - Less I/O needed
 - Less memory needed
 - Faster response
 - More users
- Page is needed \Rightarrow reference to it
 - invalid reference \Rightarrow abort
 - not-in-memory \Rightarrow bring to memory

Valid-Invalid Bit

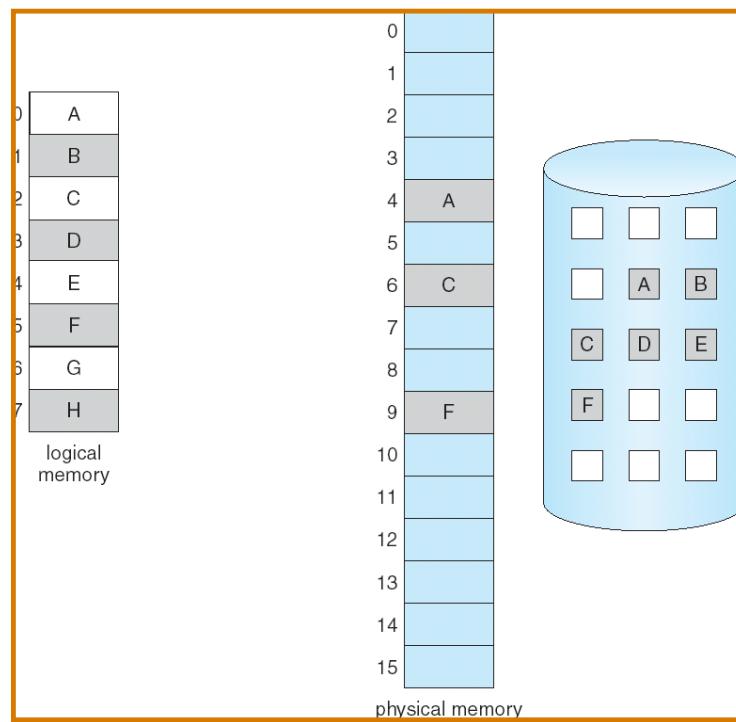
- With each page table entry a valid-invalid bit is associated (1 \Rightarrow in-memory and legal, 0 \Rightarrow not-in-memory or invalid)
- Initially valid-invalid bit is set to 0 on all entries
- Example of a page table snapshot:

Frame #	valid-invalid bit
	1
	1
	1
	1
	0
:	
	0
	0

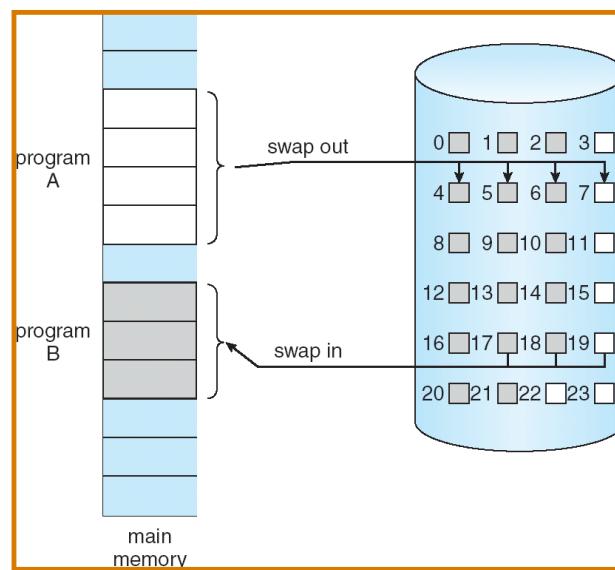
page table

- During address translation, if valid-invalid bit in page table entry is 0 \Rightarrow page fault

Page Table When Some Pages Are Not in Main Memory



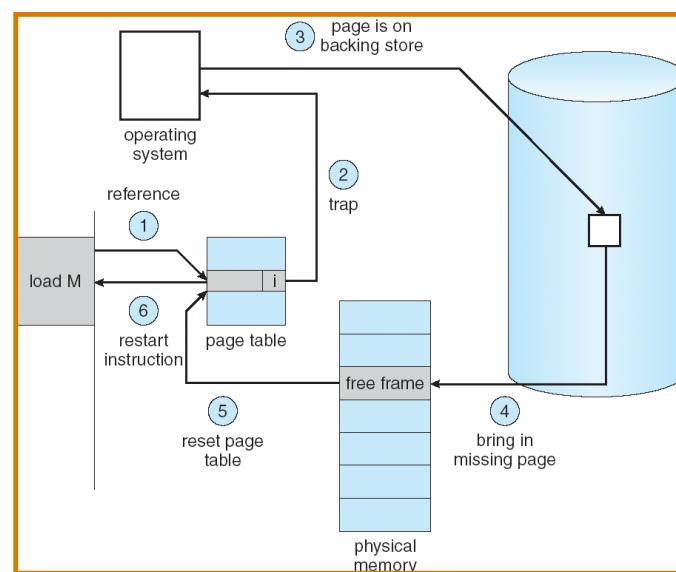
Transfer of a Paged Memory to Contiguous Disk Space



Page Fault

- If there is ever a reference to a page not in memory, first reference will trap to OS \Rightarrow page fault
- OS looks at another table (in PCB) to decide:
 - Invalid reference \Rightarrow abort.
 - Just not in memory. \Rightarrow page-in
- Get an empty frame.
- Swap (read) page into the new frame.
- Set validation bit = 1.
- Restart instruction

Steps in Handling a Page Fault



What happens if there is no free frame?

- Page replacement - find some page in memory, but not really in use, swap it out
 - Algorithms (FIFO, LRU ..)
 - performance - want an algorithm which will result in minimum number of page faults
- Same page may be brought into memory several times

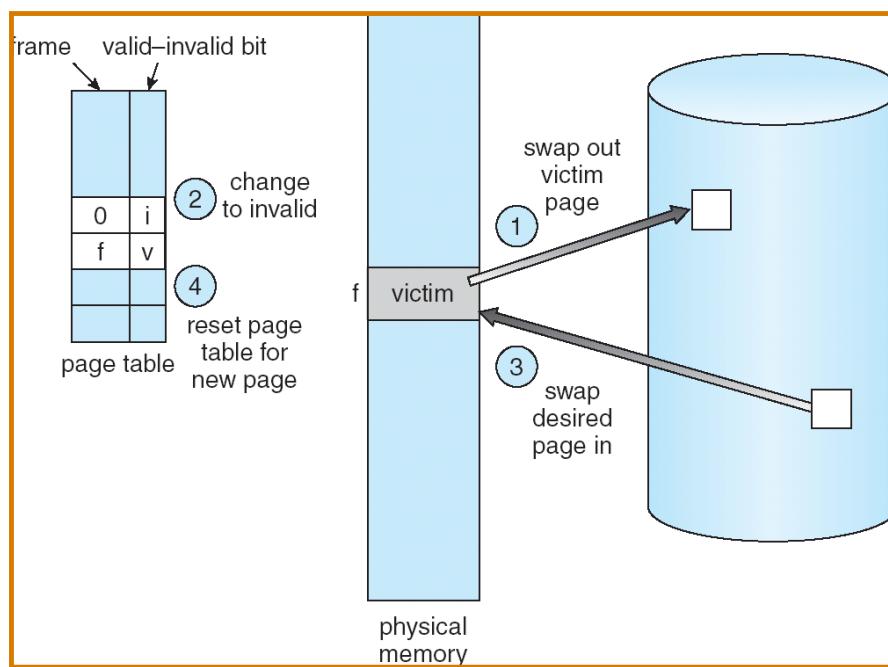
Page Replacement

- Prevent over-allocation of memory by modifying page-fault service routine to include page replacement
- Use **modify (dirty) bit** to reduce overhead of page transfers - only modified pages are written to disk
- Page replacement completes separation between logical memory and physical memory - large virtual memory can be provided on a smaller physical memory

Basic Page Replacement

1. Find the location of the desired page on disk
2. Find a free frame:
 - If there is a free frame, use it
 - If there is no free frame, use a page replacement algorithm to select a **victim** frame
3. Read the desired page into the (newly) free frame. Update the page and frame tables.
4. Restart the process

Page Replacement



Page Replacement Algorithms

- Want lowest page-fault rate
- Evaluate algorithm by running it on a particular string of memory references (reference string) and computing the number of page faults on that string
- In all our examples, the reference string is
1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5

Graph of Page Faults Versus The Number of Frames



First-In-First-Out (FIFO) Algorithm

- Reference string: 1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5
- 3 frames (3 pages can be in memory at a time per process)



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1	4	5
2	1	3
3	2	4

9 page faults

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- 4 frames

-

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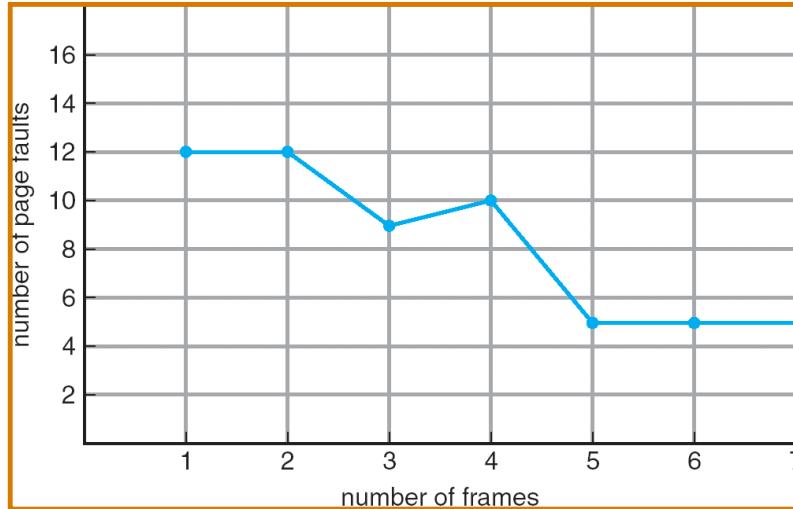
9 page faults

- 4 frames

1	1	5	4
2	2	1	5
3	3	2	
4	4	3	

- FIFO Replacement - **Belady's Anomaly**
 - more frames \Rightarrow more page faults

FIFO Illustrating Belady's Anomaly



Performance of Demand Paging

- Page Fault Rate $0 \leq p \leq 1.0$
 - if $p = 0$ no page faults
 - if $p = 1$, every reference is a fault
- Effective Access Time (EAT)
$$\text{EAT} = (1 - p) \times \text{memory access} + p \times (\text{page fault overhead} + [\text{swap page out}] + \text{swap page in} + \text{restart overhead})$$

Demand Paging Example

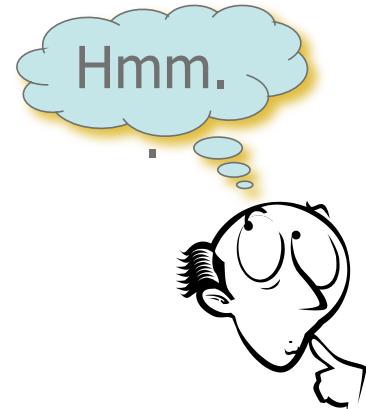
- Memory access time = 1 microsecond
- 50% of the time the page that is being replaced has been modified and therefore needs to be swapped out
- Swap Page Time = 10 msec = 10,000 microsec
- EAT = ?

Demand Paging Example

- Memory access time = 1 microsecond
- 50% of the time the page that is being replaced has been modified and therefore needs to be swapped out
- Swap Page Time = 10 msec = 10,000 microsec
- $$\begin{aligned} EAT &= (1 - p) \times 1 + p \times (10,000 + 1/2 \times 10,000) \\ &= 1 + 14,999 \times p \quad (\text{in microsec}) \end{aligned}$$
- What if 1 out of 1000 memory accesses cause a page fault?
- What if we only want 30% performance degradation?

Summary

- Main Memory Management
 - Segmentation
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- Next Lecture: Virtual Memory - II
- Reading Assignment: Chapter 9 from Silberschatz.

Acknowledgements

- “Operating Systems Concepts” book and supplementary material by A. Silberschatz, P. Galvin and G. Gagne
- “Operating Systems: Internals and Design Principles” book and supplementary material by W. Stallings
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