

CSE 421/521 - Operating Systems
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LECTURE - XXIII

DISTRIBUTED SYSTEMS - I

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Motivation

- **Distributed system** is collection of loosely coupled processors that
 - do not share memory
 - interconnected by a communications network
- Reasons for distributed systems
 - Resource sharing
 - sharing and printing files at remote sites
 - processing information in a distributed database
 - using remote specialized hardware devices
 - Computation speedup - **load sharing**
 - Reliability - detect and recover from site failure, function transfer, reintegrate failed site
 - Communication - message passing

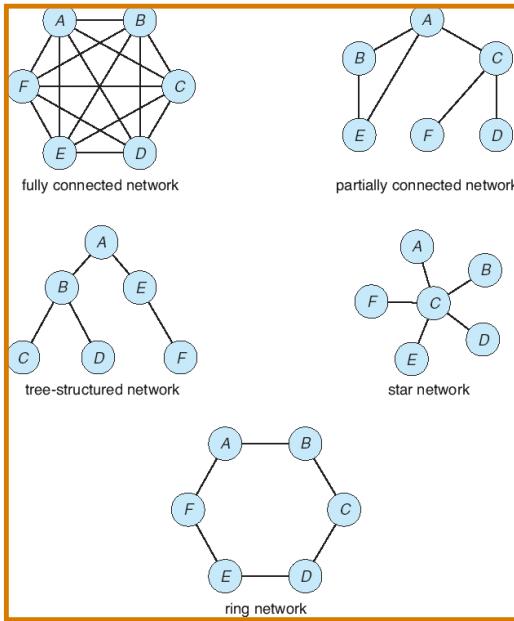
Distributed-Operating Systems

- Users not aware of multiplicity of machines
 - Access to remote resources similar to access to local resources
- Data Migration - transfer data by transferring entire file, or transferring only those portions of the file necessary for the immediate task
- Computation Migration - transfer the computation, rather than the data, across the system

Distributed-Operating Systems (Cont.)

- Process Migration - execute an entire process, or parts of it, at different sites
 - Load balancing - distribute processes across network to even the workload
 - Computation speedup - subprocesses can run concurrently on different sites
 - Hardware preference - process execution may require specialized processor
 - Software preference - required software may be available at only a particular site
 - Data access - run process remotely, rather than transfer all data locally

Network Topology



Robustness in Distributed Systems

- Failure detection
- Reconfiguration

Failure Detection

- Detecting hardware failure is difficult
- To detect a link failure, a handshaking protocol can be used
- Assume Site A and Site B have established a link
 - At fixed intervals, each site will exchange an *I-am-up* message indicating that they are up and running
- If Site A does not receive a message within the fixed interval, it assumes either (a) the other site is not up or (b) the message was lost
- Site A can now send an *Are-you-up?* message to Site B
- If Site A does not receive a reply, it can repeat the message or try an alternate route to Site B

Failure Detection (cont)

- If Site A does not ultimately receive a reply from Site B, it concludes some type of failure has occurred
- Types of failures:
 - Site B is down
 - The direct link between A and B is down
 - The alternate link from A to B is down
 - The message has been lost
- However, Site A cannot determine exactly **why** the failure has occurred

Reconfiguration

- When Site A determines a failure has occurred, it must reconfigure the system:
 1. If the link from A to B has failed, this must be broadcast to every site in the system
 2. If a site has failed, every other site must also be notified indicating that the services offered by the failed site are no longer available
- When the link or the site becomes available again, this information must again be broadcast to all other sites

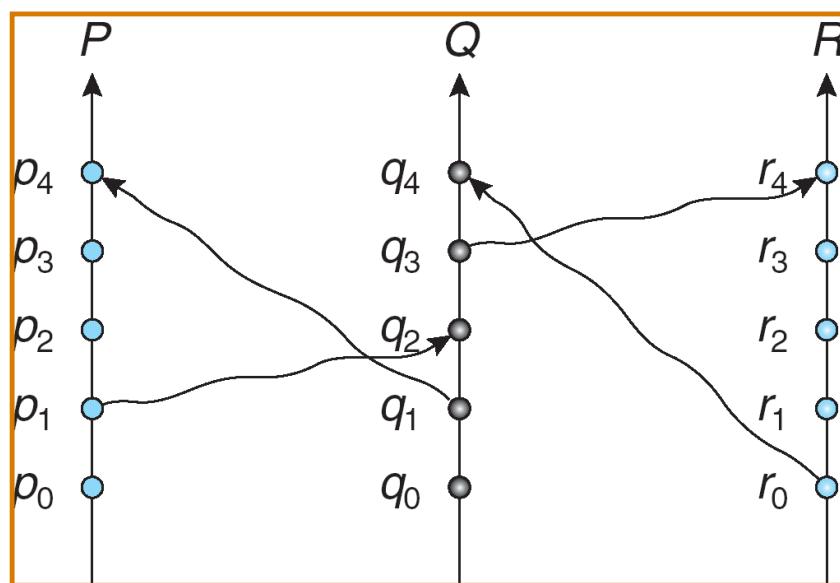
Distributed Coordination

- Ordering events and achieving synchronization in centralized systems is easier.
 - We can use common clock and memory
- What about distributed systems?
 - No common clock or memory
 - *happened-before* relationship provides partial ordering
 - How to provide total ordering?

Event Ordering

- **Happened-before** relation (denoted by \rightarrow)
 - If A and B are events in the same process (assuming sequential processes), and A was executed before B , then $A \rightarrow B$
 - If A is the event of sending a message by one process and B is the event of receiving that message by another process, then $A \rightarrow B$
 - If $A \rightarrow B$ and $B \rightarrow C$ then $A \rightarrow C$
 - If two events A and B are not related by the \rightarrow relation, then these events are executed **concurrently**.

Relative Time for Three Concurrent Processes



Which events are concurrent and which ones are ordered?

Exercise

Which of the following event orderings are true?

- (a) $p_0 \rightarrow p_3$:
- (b) $p_1 \rightarrow q_3$:
- (c) $q_0 \rightarrow p_3$:
- (d) $r_0 \rightarrow p_4$:
- (e) $p_0 \rightarrow r_4$:

Which of the following statements are true?

- (a) p_2 and q_2 are concurrent processes.
- (b) q_1 and r_1 are concurrent processes.
- (c) p_0 and q_3 are concurrent processes.
- (d) r_0 and p_0 are concurrent processes.
- (e) r_0 and p_4 are concurrent processes.

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Implementation of \rightarrow

- Associate a timestamp with each system event
 - Require that for every pair of events A and B, if $A \rightarrow B$, then the timestamp of A is less than the timestamp of B
- Within each process P_i , define a **logical clock**
 - The logical clock can be implemented as a simple counter that is incremented between any two successive events executed within a process
 - Logical clock is **monotonically increasing**
- A process advances its logical clock when it receives a message whose timestamp is greater than the current value of its logical clock
 - Assume A sends a message to B, $LC_1(A)=200$, $LC_2(B)=195 \rightarrow LC_2(B)=201$
- If the timestamps of two events A and B are the same, then the events are concurrent
 - We may use the process identity numbers to break ties and to create a total ordering

Distributed Mutual Exclusion (DME)

- Assumptions
 - The system consists of n processes; each process P_i resides at a different processor
 - Each process has a critical section that requires mutual exclusion
- Requirement
 - If P_i is executing in its critical section, then no other process P_j is executing in its critical section
- We present two algorithms to ensure the mutual exclusion execution of processes in their critical sections

DME: Centralized Approach

- One of the processes in the system is chosen to coordinate the entry to the critical section
- A process that wants to enter its critical section sends a request message to the coordinator
- The coordinator decides which process can enter the critical section next, and sends that process a reply message
- When the process receives a reply message from the coordinator, it enters its critical section
- After exiting its critical section, the process sends a release message to the coordinator and proceeds with its execution
- This scheme requires three messages per critical-section entry:
 - request
 - reply
 - release

DME: Fully Distributed Approach

- When process P_i wants to enter its critical section, it generates a new timestamp, TS , and sends the message *request* (P_i , TS) to all processes in the system
- When process P_j receives a *request* message, it may reply immediately or it may defer sending a reply back
- When process P_i receives a *reply* message from all other processes in the system, it can enter its critical section
- After exiting its critical section, the process sends *reply* messages to all its deferred requests

DME: Fully Distributed Approach (Cont.)

- The decision whether process P_j replies immediately to a *request* (P_i , TS) message or defers its reply is based on three factors:
 - If P_j is in its critical section, then it defers its reply to P_i
 - If P_j does *not* want to enter its critical section, then it sends a *reply* immediately to P_i
 - If P_j wants to enter its critical section but has not yet entered it, then it compares its own request timestamp with the timestamp TS
 - If its own request timestamp is greater than TS , then it sends a *reply* immediately to P_i (P_i asked first)
 - Otherwise, the reply is deferred
- **Example:** P1 sends a request to P2 and P3 (timestamp=10)
P3 sends a request to P1 and P2 (timestamp=4)

Undesirable Consequences

- The processes need to know the identity of all other processes in the system, which makes the dynamic addition and removal of processes more complex
- If one of the processes fails, then the entire scheme collapses
 - This can be dealt with by continuously monitoring the state of all the processes in the system, and notifying all processes if a process fails