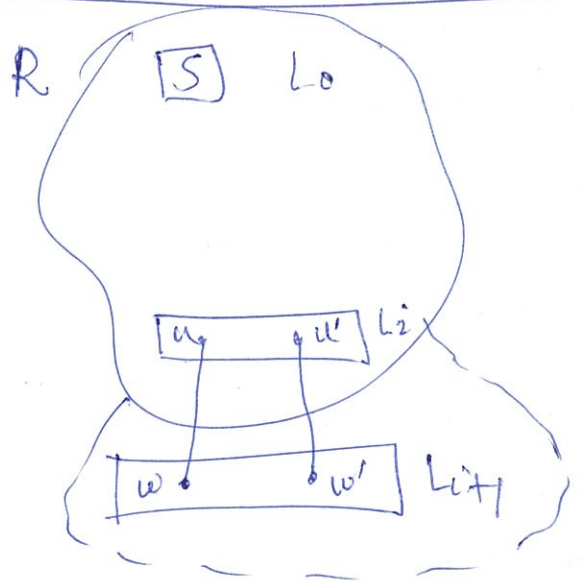


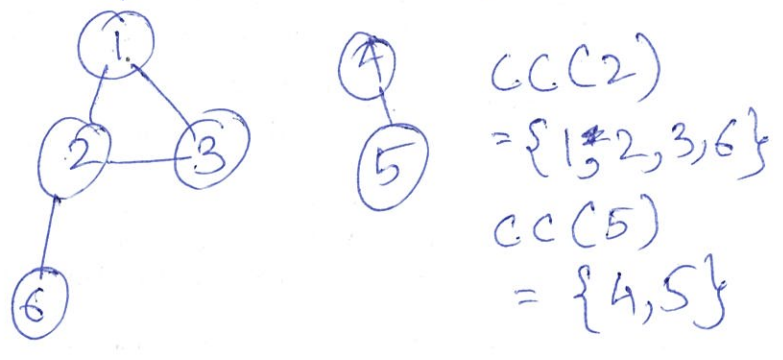
Sep 23

Explore (s)

0. $R = \{s\}$
 1. While $\exists (u,w) \in E$ s.t. $u \in R$ and $w \notin R$
Add w to R .
 2. Output $R^* = R$
- such that
- Lemo: Explore always terminates



Def: Set of all vertices connected to s is called its connected component $CC(s)$



THEOREM: for all G , start vertices s , $R^* = CC(s)$

General trick: to show $A = B \iff A \subseteq B$ and $B \subseteq A$

Lemma 1: $R^* \subseteq CC(s)$

Lemma 1 + Lemma 2 \implies THM.

Lemma 2: $CC(s) \subseteq R^*$

THM \implies BFS is correct.

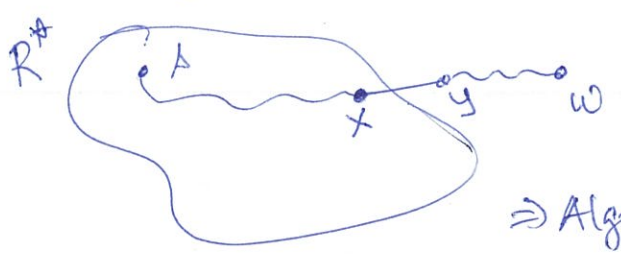
Ex: By induction

Pf idea of Lemma 2:

Pf. by contradiction

Assume $CC(s) \not\subseteq R^* \implies \exists w \in CC(s)$ BUT $w \notin R^*$

$\iff \exists s-w$ path P in G but $w \notin R^*$
 $[w \in CC(s)]$



Since P starts inside of R^* but ends up outside of R^* ,
 $\implies \exists (x,y) \in P$ s.t. $x \in R^*$ and $y \notin R^*$
 $\implies y$ should have been added to R by Explore
 \implies Algo has not terminated \implies contradicts the existence of R^*