

Sep 30

Interval Scheduling Problem of $[3, 7] = \{3, 4, 5, 6\}$

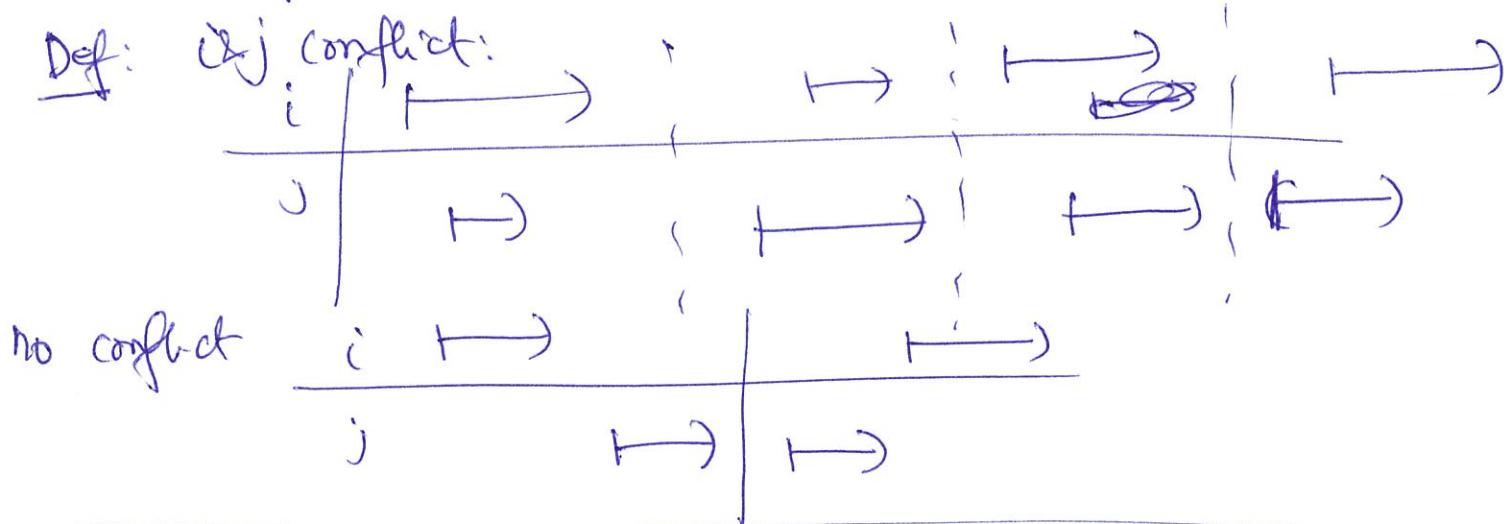
Input: n intervals : i^{th} interval $(s(i), f(i))$

Output: A valid schedule with max # intervals

Def: A schedule $S \subseteq [n]$ ($= \{1, \dots, n\}$)

Def: A valid schedule S has no conflicts

Def: i, j conflict:



Obs: A valid schedule sorted $\rightarrow \rightarrow \rightarrow \rightarrow \rightarrow \cdots$
— by start or finish time gives same order

Assume: Input intervals are sorted by finish time
 $(f(1) \leq f(2) \leq \dots \leq f(n))$

If not sort in $O(n \log n)$ time

Greedy Algo

0. $R = [n]$

1. $S = \emptyset$

2. While $R \neq \emptyset$

(2.1) Let i be the smallest index in R

(2.2) Add i to S

(2.3) Remove ~~all~~ i from R

(2.4) Delete $\{j\}$ from R that conflict with i

3. Return $S^* = S$

Thm 1: S^* is an optimal solution
↳ # inputs, among all possible valid schedules for given input, S^* has max # jobs.

Ex1: Algo terminates

Ex2: S^* is a valid schedule

Pf. of correctness of greedy algo $\xrightarrow{\text{Greedy stays ahead (next)}}$ Exchange argument (min, max, lateness: see 4.2)

Let Θ be an optimal solution

Ex3: Convince yourself that such an Θ ∃.

Idea: $S^* = \Theta$ $\xrightarrow{\quad}$ Θ

↳ problem: Can have $\xrightarrow{\quad} S^*$ > 1 optimal soln.

Idea': $|S^*| = |\Theta|$