

Lecture 15

CSE 331

Oct 4, 2021

Please have a face mask on

Masking requirement



UR requires all students, employees and visitors – regardless of their vaccination status – to wear face coverings while inside campus buildings.

<https://www.buffalo.edu/coronavirus/health-and-safety/health-safety-guidelines.html>

If you need it, ask for help



Project groups

note @244

stop following 4 views

Delay in setting up random groups for 331 project

I had originally planned to finalize the random groups (for those who had signed up to be assigned a random group) by tonight but I'm running a bit behind schedule-- my current plan is to get this done by tomorrow (or Tue at the absolute latest). I'll email the groups once they are set (as well as confirm the groups for those who signed up as a group of size 3 but had not confirmed those groups before in [@217](#)).

Apologies for the delay!

project

edit

good note

Updated 7 minutes ago by AMR Rubin

Quiz 1 this FRIDAY

note @220     stop following **33** views

Quiz 1 on Friday, Oct 8

The first quiz will be from 10:20-10:30am in class on Friday, October 8. We will have a 5 mins break after the quiz and the lecture will start at 10:35am.

We will hand out the quiz paper at 10:15am but you will **NOT** be allowed to open the quiz to see the actual questions till 10:20am. However, you can use those 5 minutes to go over the instructions and get yourself in the zone.

There will be two T/F with justification questions (like those in the sample mid term 1: [@197](#).) Also quiz 1 will cover all topics we cover in class till Friday, Oct 4.

Also like the mid-term y'all can bring in one letter sized cheat-sheet (you can use both sides). But other than cheatsheet and writing implements nothing else is allowed.

[edit](#)

[good note](#) 

Updated 2 hours ago by [Abi-Ruthe](#)

Mid-term post 1

note @218

stop following 63 views

The mid-term post

First, midterm-I is on **Monday, Oct 11** and midterm-II is on **Wednesday, Oct 13** during the usual class timings (i.e. 10:20-11:10am in Knox 110). Below are some comments that might be helpful to prepare for the mid-term.

(Thoughts on what to do during the exam here: 219)

- Work through the sample mid-term exams ([@197](#)). Do **not** use the sample mid-term to deduce **anything** about the relative coverage of different topics. (See points below for more on the coverage.) The sample mid-terms are meant for you to see the format of the questions. The *actual* mid-term exams will be harder than the sample mid-term exams. The actual mid-terms will follow the exact same format for the sample mid-terms: i.e. first mid-term will be only T/F while the second ones will be longer ones.
- I encourage you to not look at the solutions to the sample mid-terms before you have spent some quality time by yourself on the mid-term questions first.
- Use the quiz on Oct 8 ([@220](#)) to get some practice in solving T/F questions under some time pressure. Also review the T/F polls for more examples of such T/F questions.
- Review the HW problems/solutions. HW solutions are here: [@176](#).
- You **will** be under (a bit of) time pressure in the mid-term exams-- it might be useful for you to use the sample mid-term to decide on how much time you are going to spend on each question. Also read the instructions on the first page and keep them in mind during the exam (the instructions will of course be repeated on the exam sheet).
- If you need help attend the usual recitation, office hours. We will have extra office hours (details TBA) next week and the week after.
- The exam will be closed book and closed notes. However, you can bring in **one** 8.5" X 11" review sheet. (If you prefer you can bring in different review sheets for the two mid-term exams.) You can write anything that you want on the sheet as long as it is one sheet (you can use both sides). It can hand-written or typed up doesn't matter-- however, you are not allowed to bring in a magnifying glass. The review sheet is to make sure you do not spend time memorizing definitions etc. but can concentrate on the main ideas in the material we have covered. The exam (as you can probably make out from the sample mid-term) will focus on how well you understand the material and not how well you can memorize. However, see next point.
- Do not spend too much time cramming stuff into the review sheet. In my experience (both as a student and instructor), it never helps to just put in arbitrary stuff. However, you should use the review sheet to write down references for various algos etc. we have seen in class/HWs/recitation notes etc., so that you can just

Mid-term post 2

note @219

stop following 67 views

Few thoughts on what to do during the exam

In a previous post @218, I listed some pointers on how I think you should prepare for the mid-term exams.

Below are (in no particular order) some thoughts on how you should work on the actual exam:

1. **Do NOT panic (or delay it as much as possible)!** And I don't mean this in either a joking way or a scary way. In these kinds of exams once you panic everything else that follows will not be good. (Believe me I have been there.) So the idea for you will be to avoid panicking as much as possible or mitigate its effects. Here are some specific pointers in this regard:
 - Read **all** the questions even before you start writing anything. This way if you are short on time and you are not done at least you will be working on a question that you have read before: trying to make sense of a question that you are reading for the first time and under time pressure never ends well.
 - You know the structure and number of questions. Make sure you setup a time table on how much time you want to spend on each questions and stick to that plan. Make sure you keep at least 10 mins at the end to go over all your answers to make sure you were not missing something.
 - Make sure you stick to your timetable and avoid the sunk cost fallacy. Thinking that I have already spent 5 mins on a question so let me spend a couple more mins to try and crack the question often leads to you spending 15 mins on the question and then you are terribly short on time.
 - I try to order the questions from easiest to hardest and I think I do fine on the average but the ordering might not match with yours. E.g. for some reason you might have studied a particular part of the book the night before the exam and that part might be relevant to say the last question. So what I think might a hard question for an average student in the class might be easy for you. Reading through all questions upfront will also help you identify these "out of order" questions.
2. **Try to reinvent as little of the wheel as possible.**
 - Your first attack on any problem should be to see if you can sufficiently modify the question/input to the algorithm so that you can use a solution from a previous HW problem/the book/stuff on piazza as a blackbox. Note this is the same philosophy as to why you should libraries instead of writing code from scratch.
 - Remember how easy it was to get most points on 1(a) and 2(a) by just referring to the recitation notes. Y'all should try to do that as much as possible!
 - If you try and build something from scratch (like an algorithm or a proof) that you could have just referenced away, then you will be short on time. The exam is timed in such a way that if something can be just referenced, then you are expected to do so.
 - If the above fails then you should see if you can slightly tweak a previous solution to work in the current situation. Most of the problems in the mid terms will either be in the previous category or this one.
 - If both of the above fail, then try to answer from scratch but **this should be your last option.**

Feedback on CSE 331

note @245    0 views Actions

Feedback on CSE 331

Every year, I ask y'all to give feedback on CSE 331, so here is the feedback form for this year:

https://docs.google.com/forms/d/1FAIpQL5cQDQCd11Igm9oYUzWm65KA2HUQISjKs6uFp6pZFnbtv1m5A/viewform?usp=ref_link

Filling in this form is **completely optional and anonymous**.

In particular, I would love feedback (even if it is critical). Many of the aspects of CSE 331 that you like were suggested by someone in a previous incarnation of CSE 331. While I'll try and incorporate as much as I can this fall, some of your suggestions might have to wait for the next offering.

I might also dis-agree with your feedback but after a week or so, I'll post my response to the feedback from y'all. So at the very least y'all would get to hear my reasoning for why certain things are the way they are in CSE 331. And then we can agree to disagree :-)

[Feedback](#)

 undo good note |  Updated Just now by Aar Rudra

Questions?



Interval Scheduling Problem

Input: n intervals $[s(i), f(i))$ for $1 \leq i \leq n$



$\{ s(i), \dots, f(i)-1 \}$

Output: A schedule S of the n intervals

No two intervals in S conflict

$|S|$ is maximized

Algorithm with examples

Interval Scheduling via examples

In which we derive an algorithm that solves the Interval Scheduling problem via a sequence of examples.

The problem

In these notes we will solve the following problem:

Interval Scheduling Problem

Input: An input of n intervals $[s(i), f(i))$, or in other words, $\{s(i), \dots, f(i) - 1\}$ for $1 \leq i \leq n$ where i represents the intervals, $s(i)$ represents the start time, and $f(i)$ represents the finish time.

Output: A schedule S of n intervals where no two intervals in S conflict, and the total number of intervals in S is maximized.

Sample Input and Output

Input:

Re-define problem on the board...

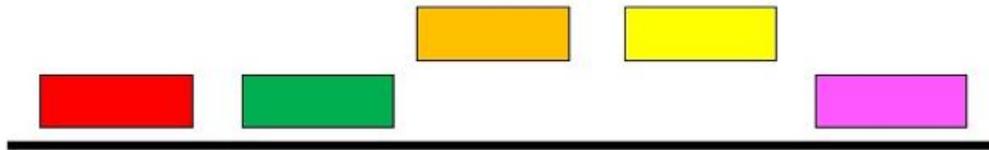


Example 1

No intervals overlap



Algorithm?



No intervals overlap

R : set of requests

Set S to be the empty set

While R is not empty

 Choose i in R

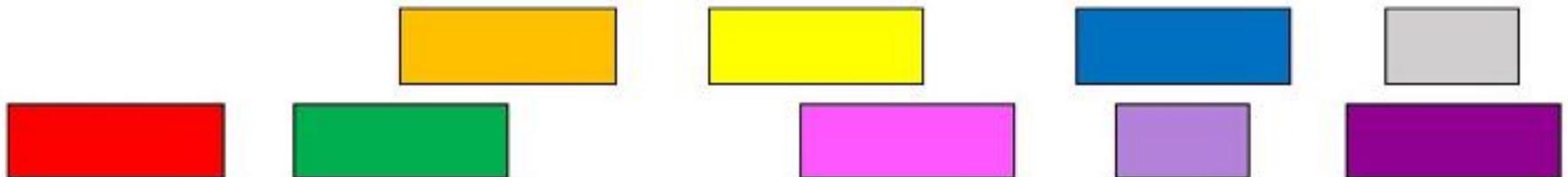
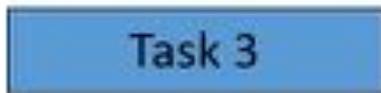
 Add i to S

 Remove i from R

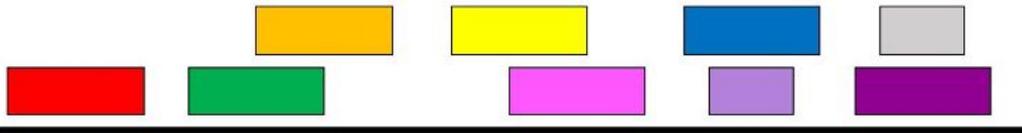
Return $S^* = S$

Example 2

At most one overlap/task



Algorithm?



At most one overlap

R : set of requests

Set S to be the empty set

While R is not empty

 Choose i in R

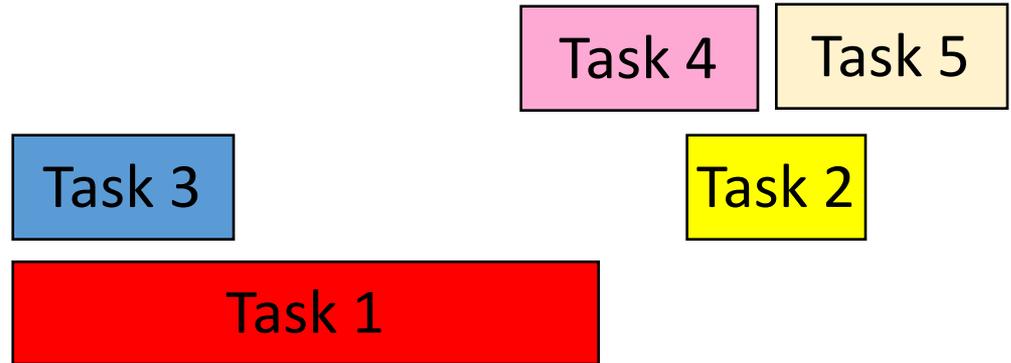
 Add i to S

 Remove all tasks in R that conflict with i from R

Return $S^* = S$

Example 3

More than one conflict



Set S to be the empty set

While R is not empty

 Choose i in R

 Add i to S

 Remove all tasks that conflict with i from R

Return $S^* = S$

Greedily solve your blues!

Arrange tasks in some order and iteratively pick non-overlapping tasks



Write up a term paper

Party!

Exam study

331 HW

Project

Saturday

Sunday

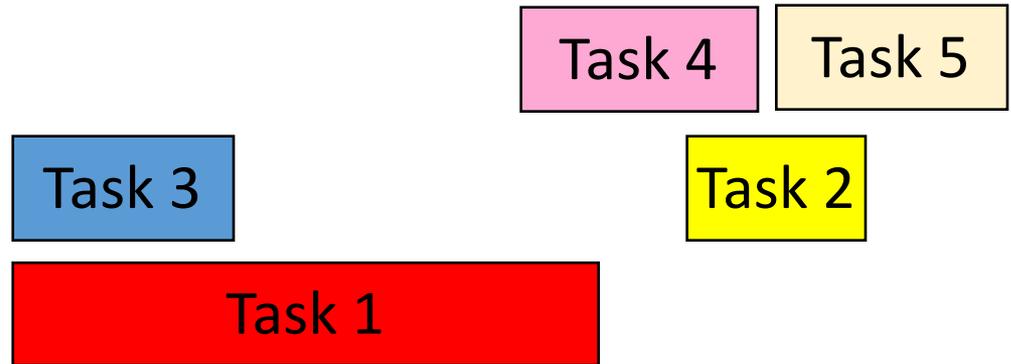
Monday

Tuesday

Wednesday

Making it more formal

More than one conflict



Set S to be the empty set

While R is not empty

Choose i **in** R that minimizes $v(i)$

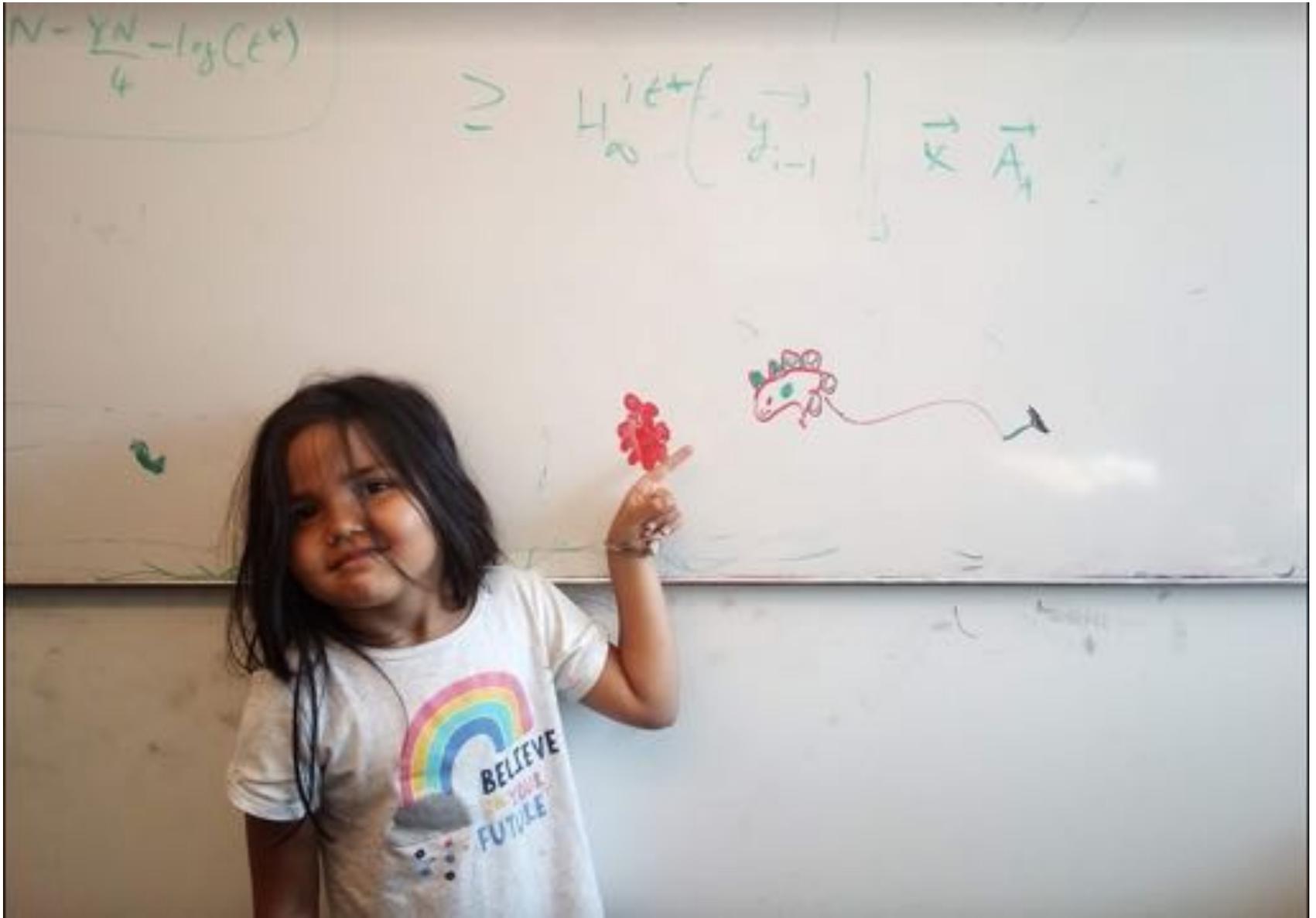
 Add i to S

 Remove all tasks that conflict with i from R

Return $S^* = S$

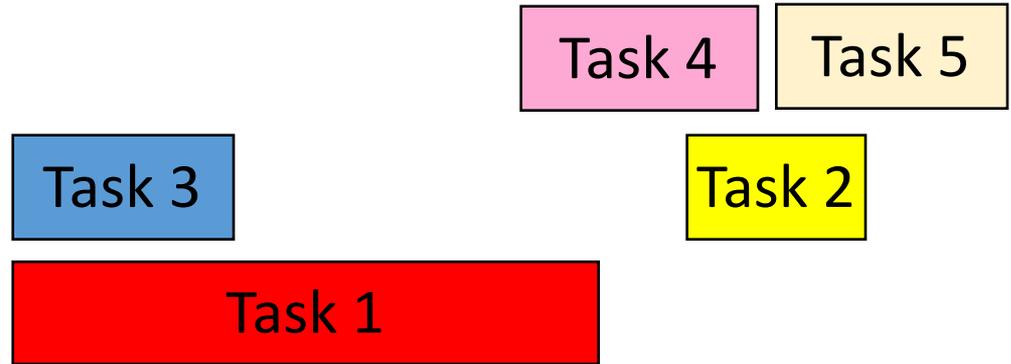
Associate a
value $v(i)$
with task i

Questions/Comments?



What is a good choice for $v(i)$?

More than one conflict



Set S to be the empty set

While R is not empty

 Choose i in R that minimizes $v(i)$

 Add i to S

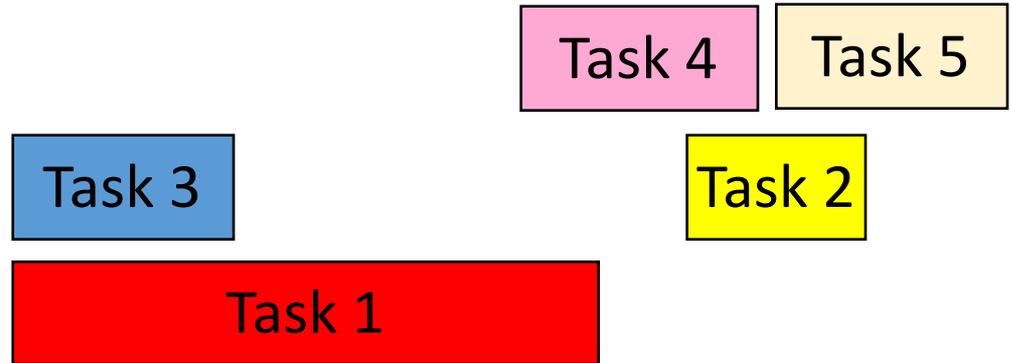
 Remove all tasks that conflict with i from R

Return $S^* = S$

Associate a
value $v(i)$
with task i

$$v(i) = f(i) - s(i)$$

Smallest duration first



Set S to be the empty set

While R is not empty

 Choose i in R that minimizes $f(i) - s(i)$

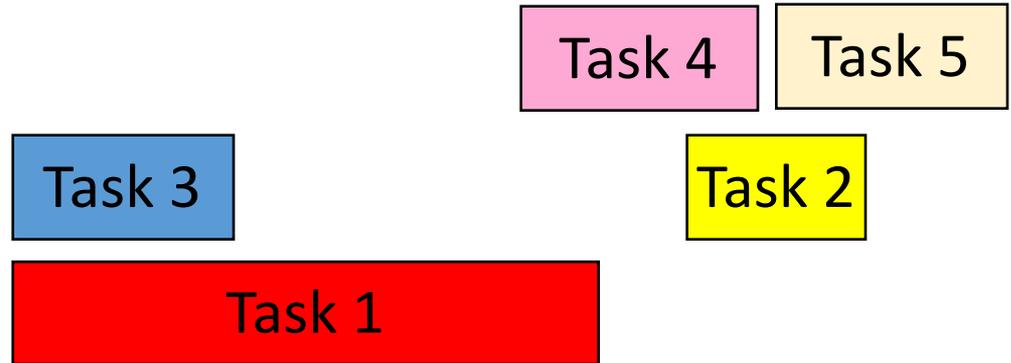
 Add i to S

 Remove all tasks that conflict with i from R

Return $S^* = S$

$$v(i) = s(i)$$

Earliest time first?



Set S to be the empty set

While R is not empty

 Choose i in R that minimizes $s(i)$

 Add i to S

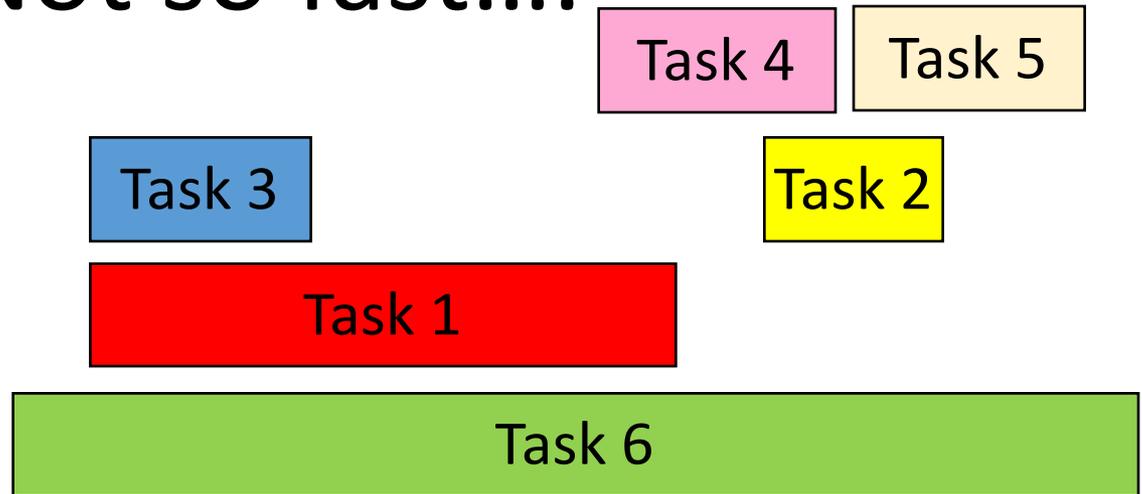
 Remove all tasks that conflict with i from R

Return $S^* = S$

So are we
done?

Not so fast....

Earliest time first?



Set S to be the empty set

While R is not empty

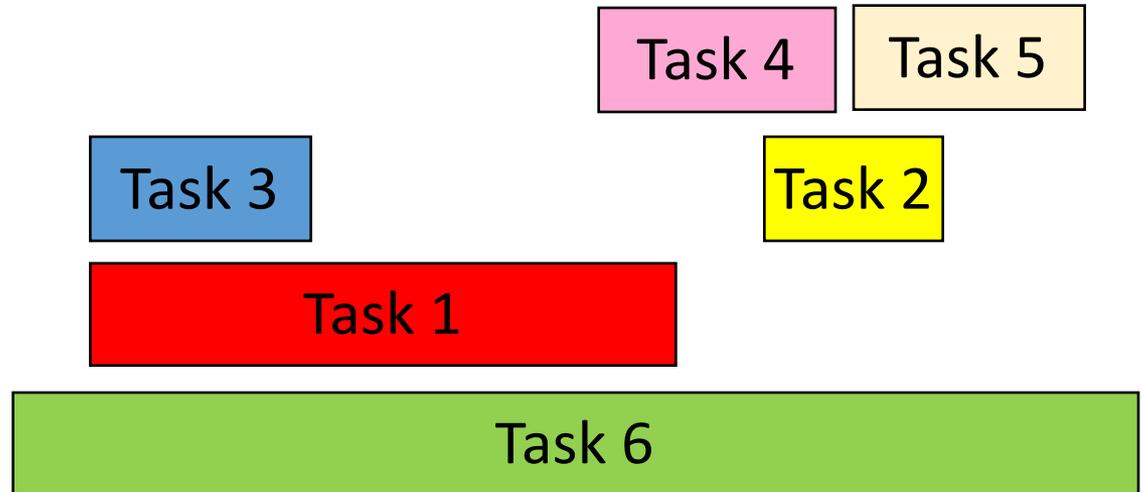
 Choose i in R that minimizes $s(i)$

 Add i to S

 Remove all tasks that conflict with i from R

Return $S^* = S$

Pick job with minimum conflicts



Set S to be the empty set

While R is not empty

 Choose i in R that has smallest number of conflicts

 Add i to S

 Remove all tasks that conflict with i from R

Return $S^* = S$

So are we
done?

Nope (but harder to show)

Set S to be the empty set

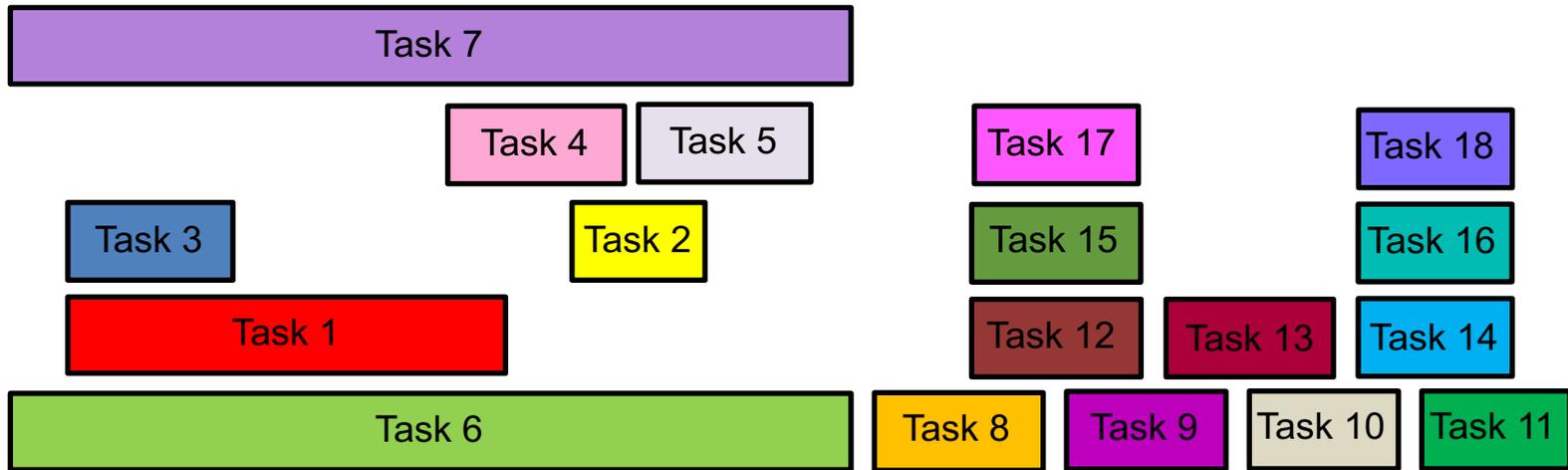
While R is not empty

 Choose i in R that has smallest number of conflicts

 Add i to S

 Remove all tasks that conflict with i from R

Return $S^* = S$



Set S to be the empty set

While R is not empty

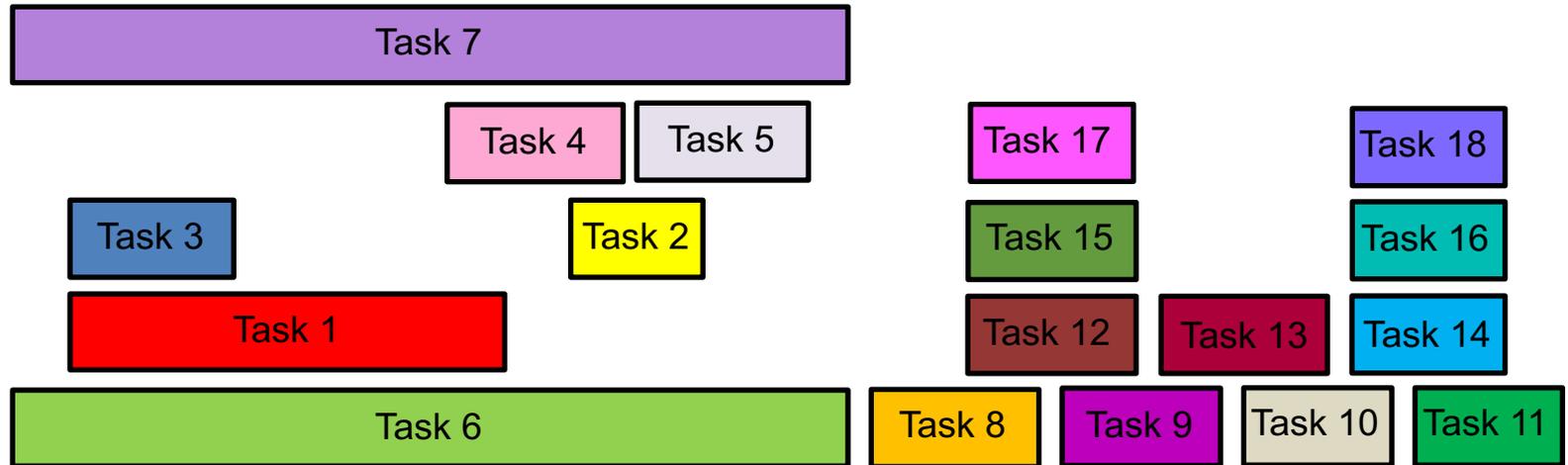
 Choose i in R that has smallest number of conflicts

 Add i to S

 Remove all tasks that conflict with i from R

Return $S^* = S$

Algorithm?



Set S to be the empty set

While R is not empty

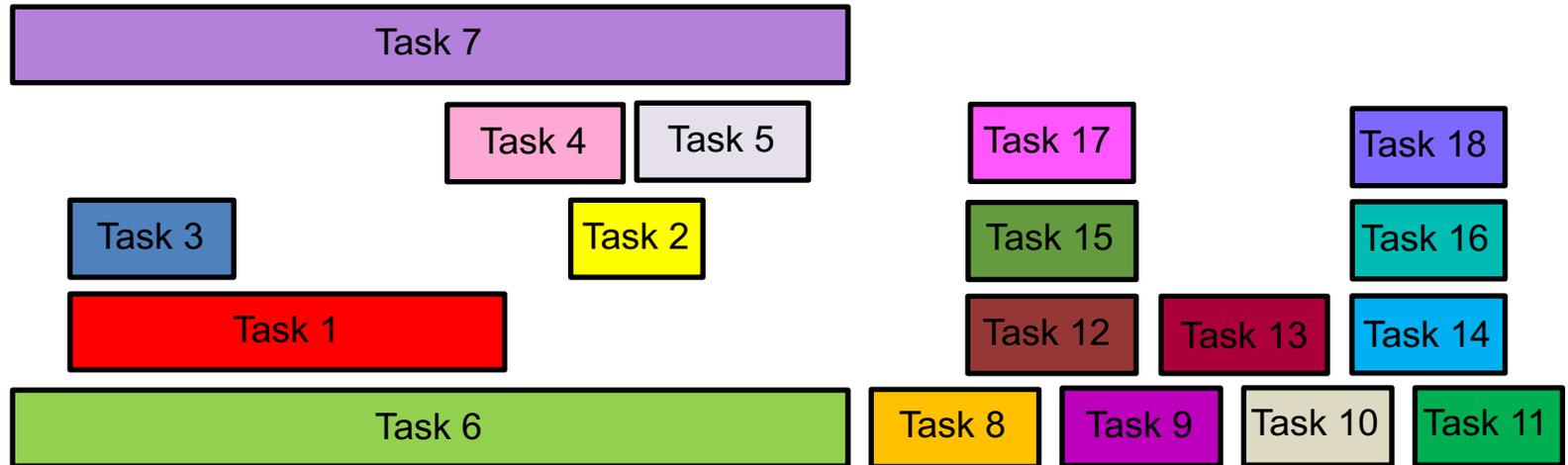
 Choose i in R that minimizes $v(i)$

 Add i to S

 Remove all tasks that conflict with i from R

Return $S^* = S$

Earliest finish time first



Set S to be the empty set

While R is not empty

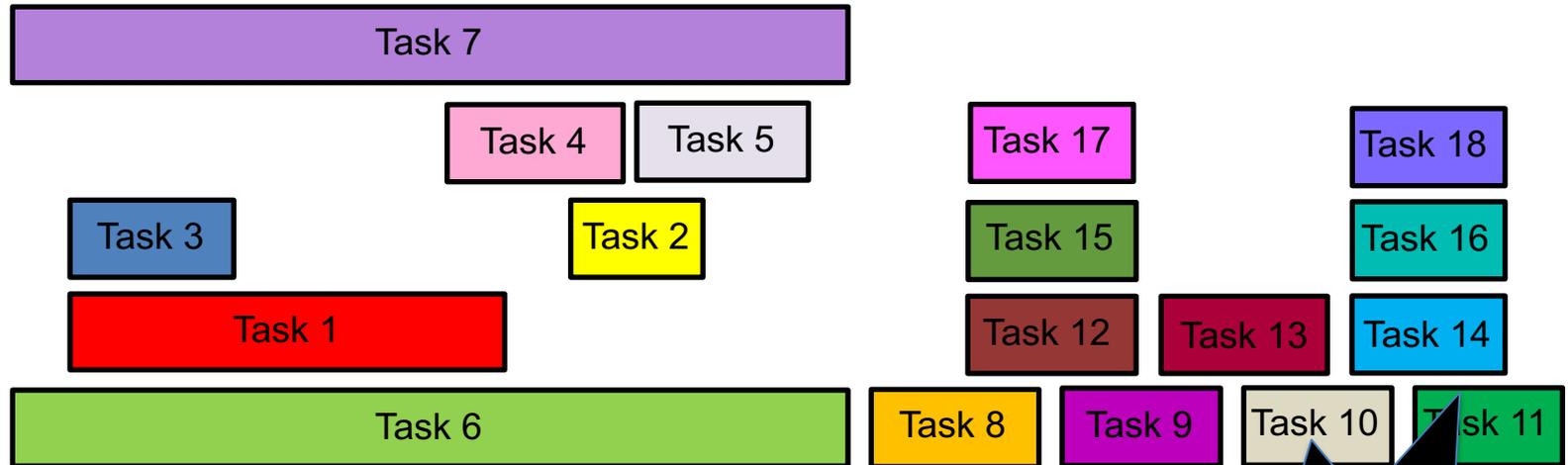
 Choose i in R that minimizes $f(i)$

 Add i to S

 Remove all tasks that conflict with i from R

Return $S^* = S$

Find a counter-example?



Set S to be the empty set

While R is not empty

 Choose i in R that minimizes $f(i)$

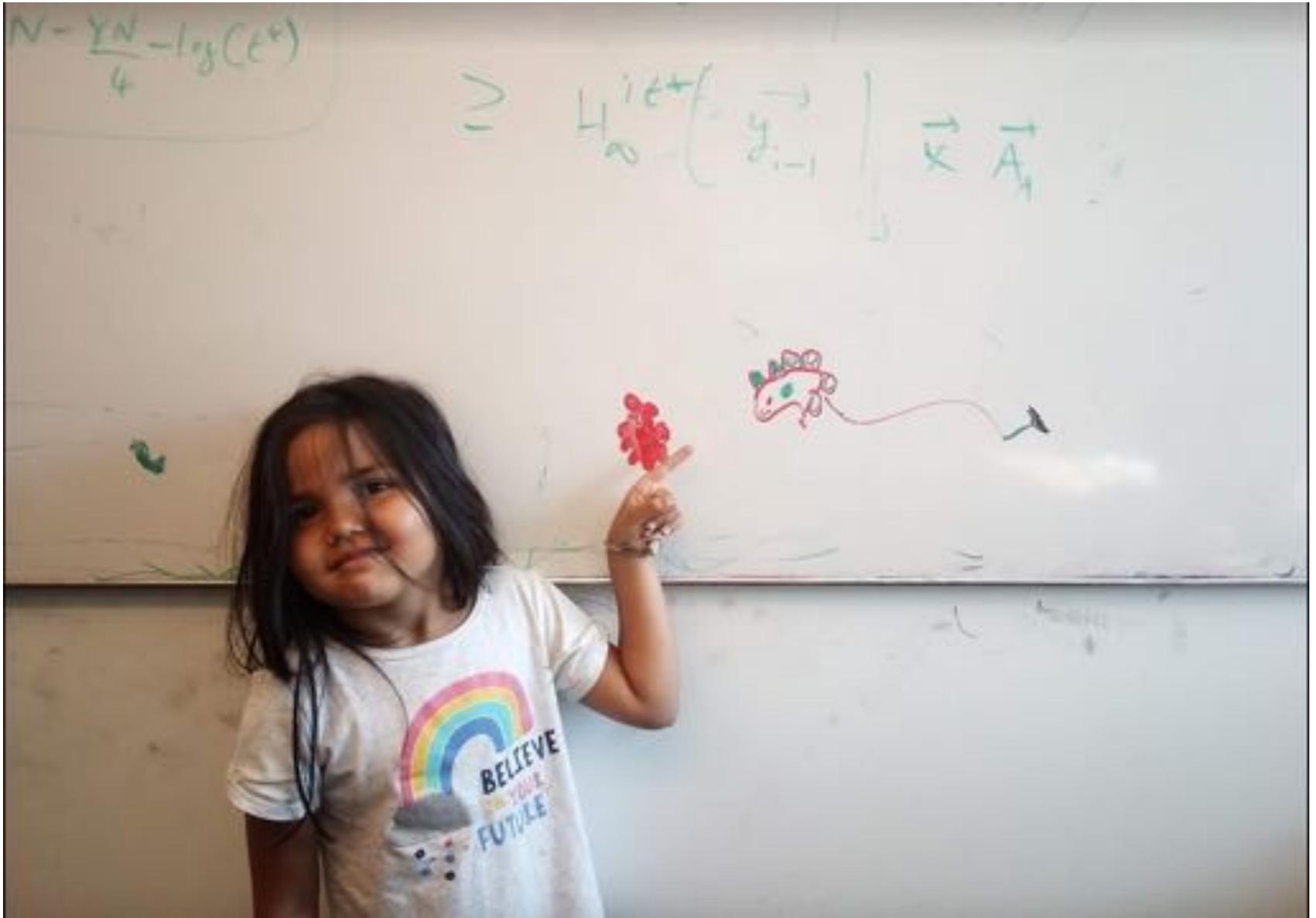
 Add i to S

 Remove all tasks that conflict with i from R

Return $S^* = S$

It
works!

Questions/Comments?



Today's agenda

Prove the correctness of the algorithm

Final Algorithm

R : set of requests

Set S to be the empty set

While R is not empty

 Choose i in R with the earliest finish time

 Add i to S

 Remove all requests that conflict with i from R

Return $S^* = S$

Argue correctness on the board...

