Lecture 27

CSE 331

Nov 4, 2022

Coding P2 due TODAY

Fri, Nov 4	Kickass Property Lemma F21 F19 F18 F17 x²	[KT, Sec 5.4] (**Project (Problem 2 `Coding`) in**)
Mon, Nov 7	Weighted Interval Scheduling F21 F19 F17 x²	[KT, Sec 6.1] (**Project (Problem 2 `Reflection` in**)
Tue, Nov 8		(**HW 6 out**)
Wed, Nov 9	Recursive algorithm for weighted interval scheduling problem F21 F19 F17 x²	[KT, Sec 6.1]
Fri, Nov 11	Subset sum problem F21 F19 F18 F17 x²	[KT, Sec 6.1, 6.2, 6.4]
Mon, Nov 14	Dynamic program for subset sum F21 F19 F18 F17 x²	[KT, Sec 6.4]
Tue, Nov 15		(**HW 7 out, HW 6 in**)
Wed, Nov 16	Shortest path problem F21 F19 F18 F17 x²	[KT, Sec 6.8]
Fri, Nov 18	Bellman-Ford algorithm F21 F19 F18 F17 x²	[KT, Sec 6.8]
Mon, Nov 21	The P vs. NP problem F21 F19	[KT, Sec 8.1]
Wed, Nov 23	**No class**	Fall Recess
Fri, Nov 25	**No class**	Fall Recess
Mon, Nov 28	More on reductions F21 F19	[KT, Sec 8.1]
Tue, Nov 29		(**HW 8 out, HW 7 in**)
Wed, Nov 30	The SAT problem F21 F19	[KT, Sec 8.2]
Fri, Dec 2	NP-Completeness F21 F19	[KT, Sec. 8.3, 8.4] (**Project (Problem 3 `Coding`) in**)
Mon, Dec 5	*k*-coloring problem F21 F19	[KT, Sec 8.7] (**Quiz 2**) (**Project (Problem 3 `Reflection`) in**)

Group formation instructions

Autolab group submission for CSE 331 Project

The lowdown on submitting your project (especially the coding and reflection) problems as a group on Autolab.

Follow instructions **EXACTLY** as they are stated

The instruction below are for Coding Problem 1

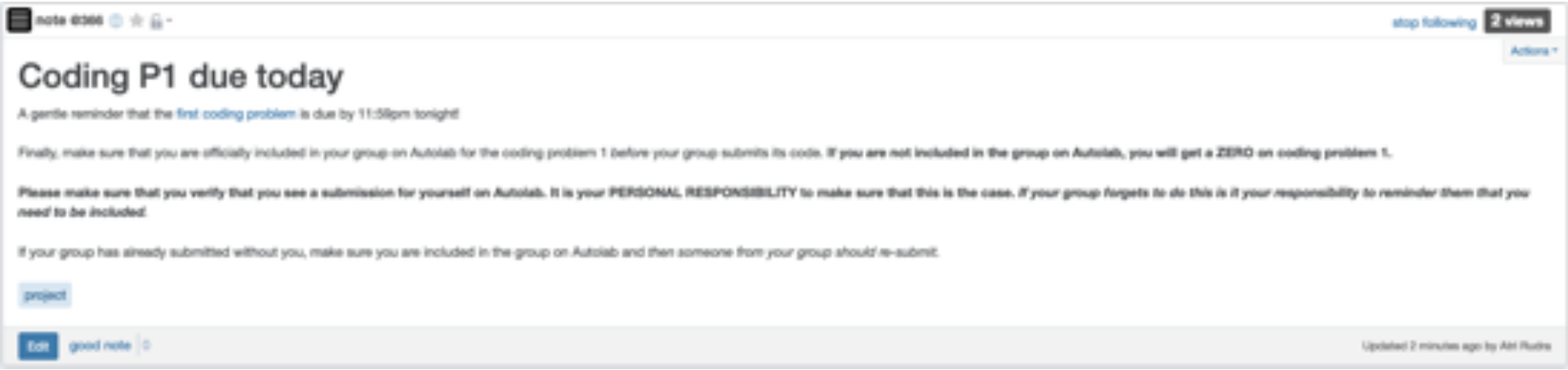
You will have to repeat the instructions below for EACH coding AND reflection problem on project on Autolab (with the appropriate changes to the actual problem).

Form your group on Autolab

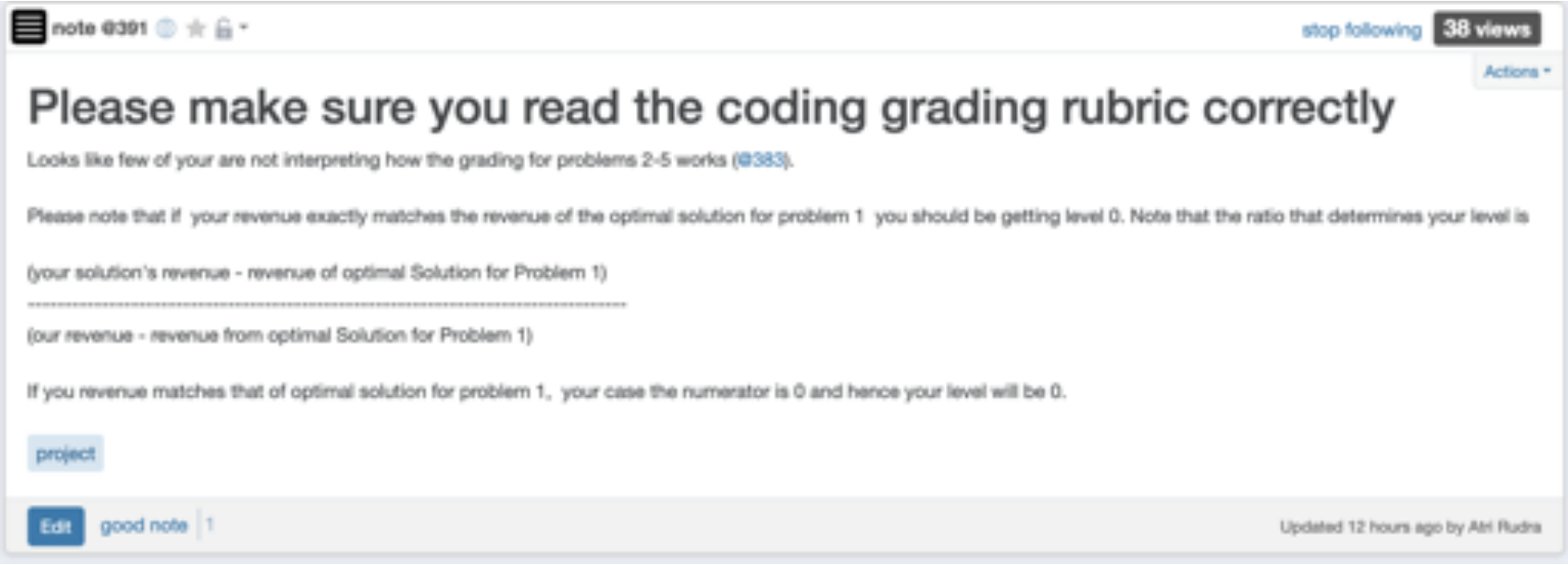
Groups on Autolab will NOT be automatically created

You will have to form a group on Autolab by yourself (as a group). Read on for instructions on how to go about this.

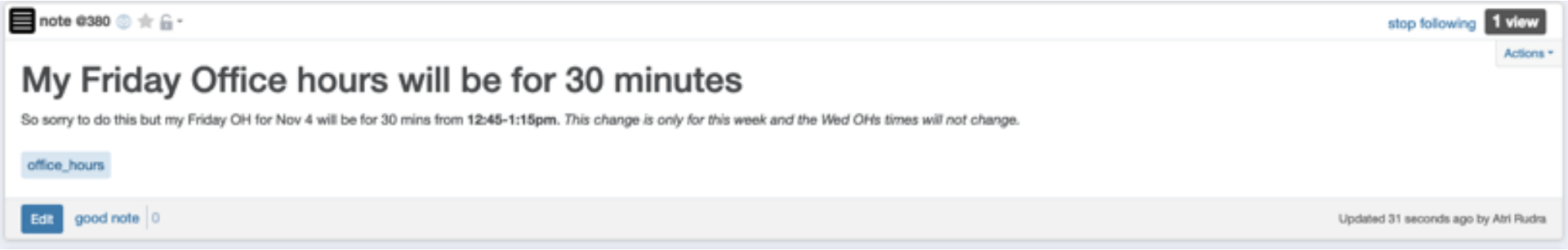
Make sure you are in your group



Coding problem grading

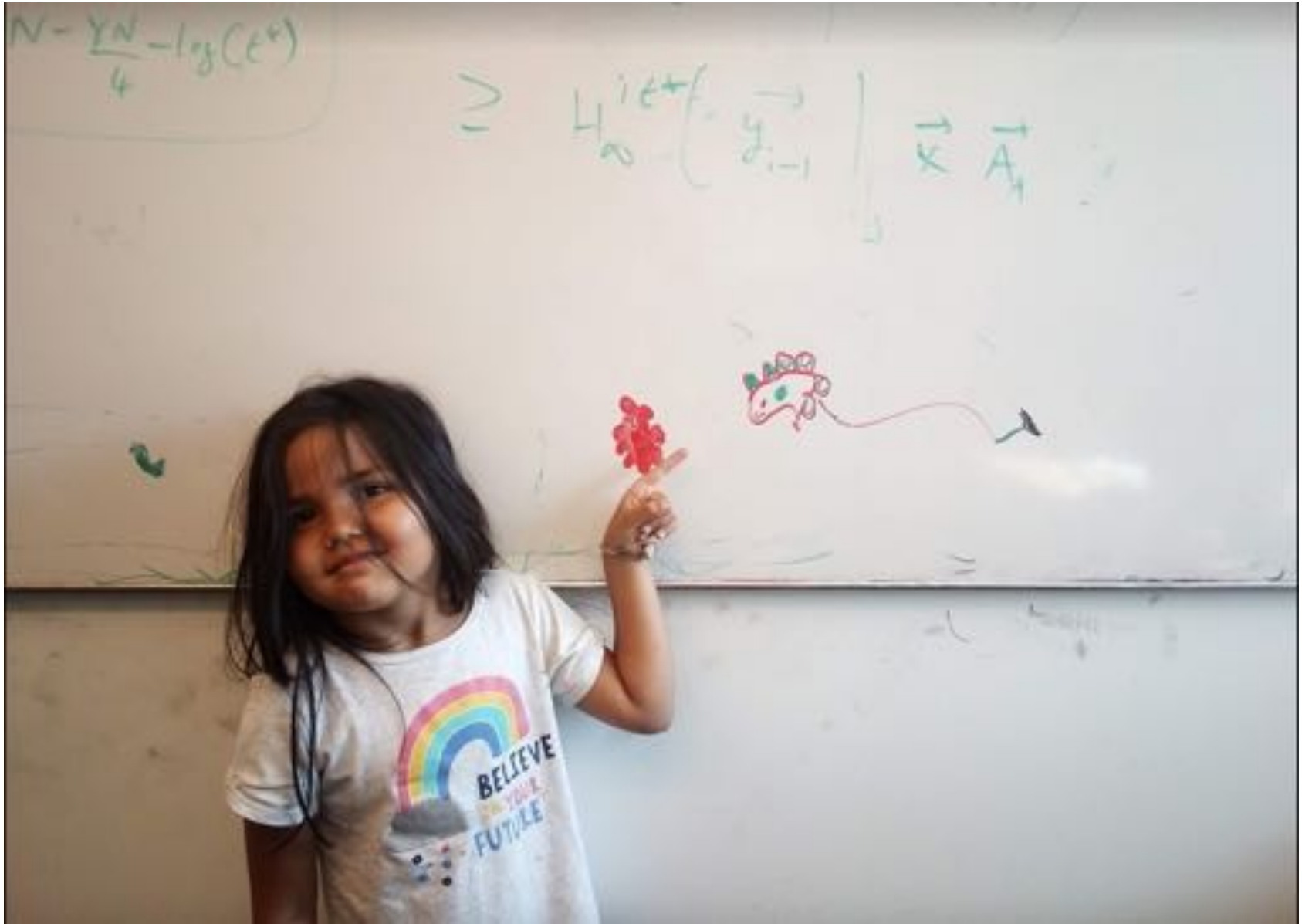


OH today shortened to 30 mins



Have fun @ UB Hacking!

Questions/Comments?

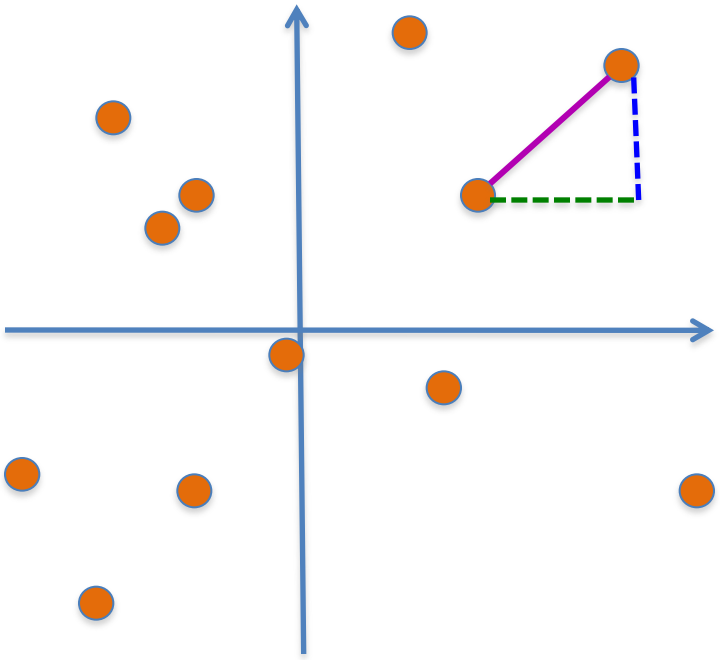


Closest pairs of points

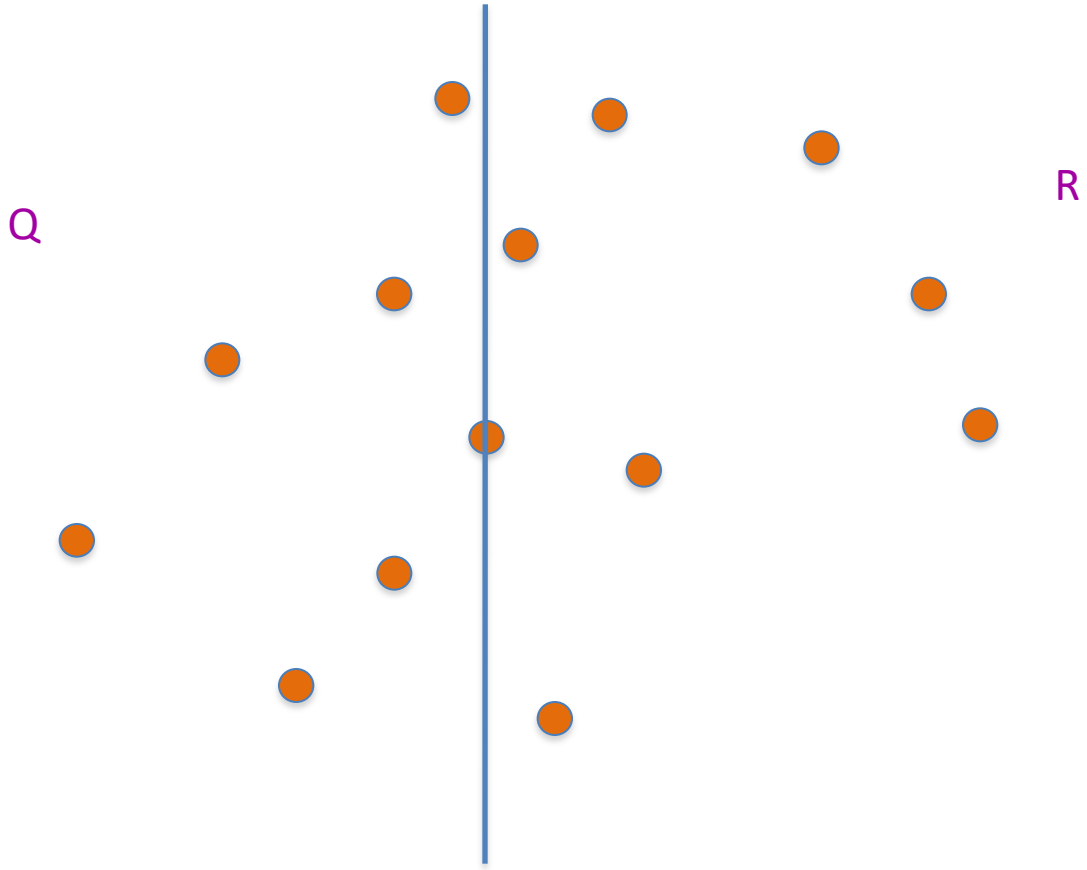
Input: n 2-D points $P = \{p_1,\ldots,p_n\}$; $p_i=(x_i,y_i)$

$d(p_i,p_j) = ((x_i-x_j)^2+(y_i-y_j)^2)^{1/2}$

Output: Points p and q that are closest

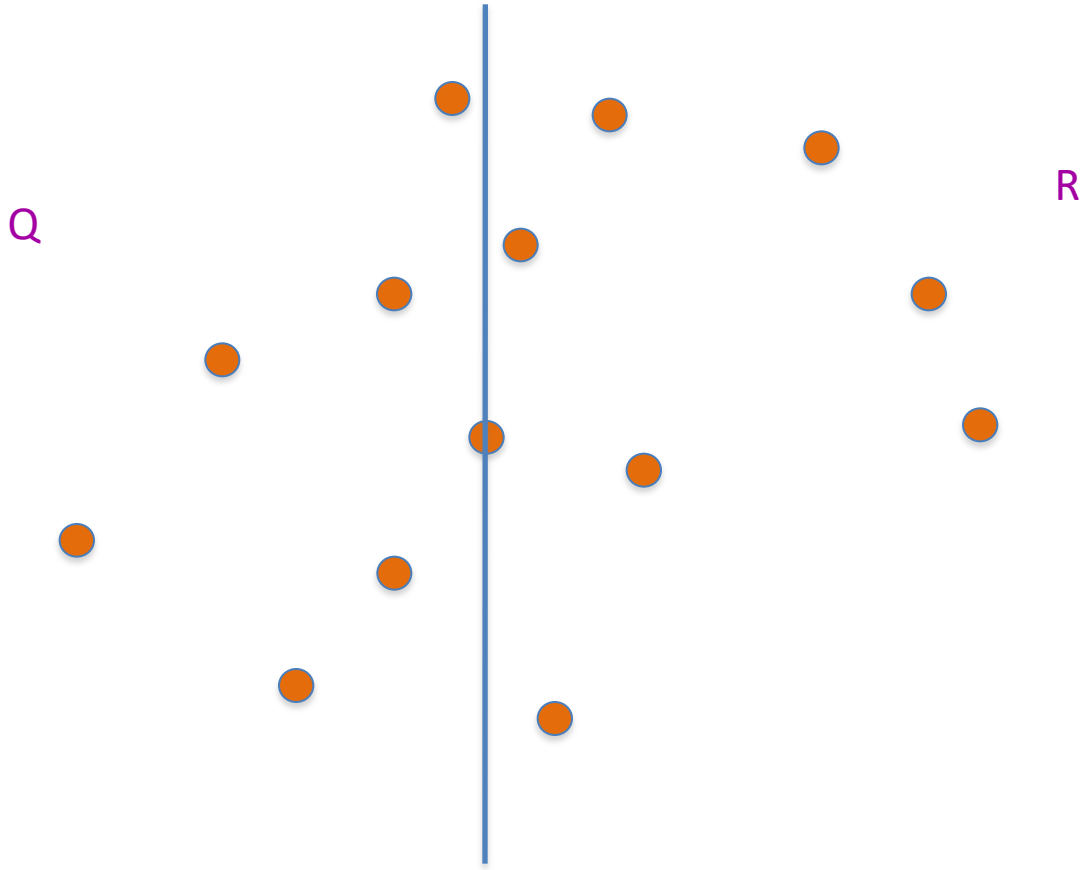


Dividing up P



First n/2 points according to the x-coord

Recursively find closest pairs



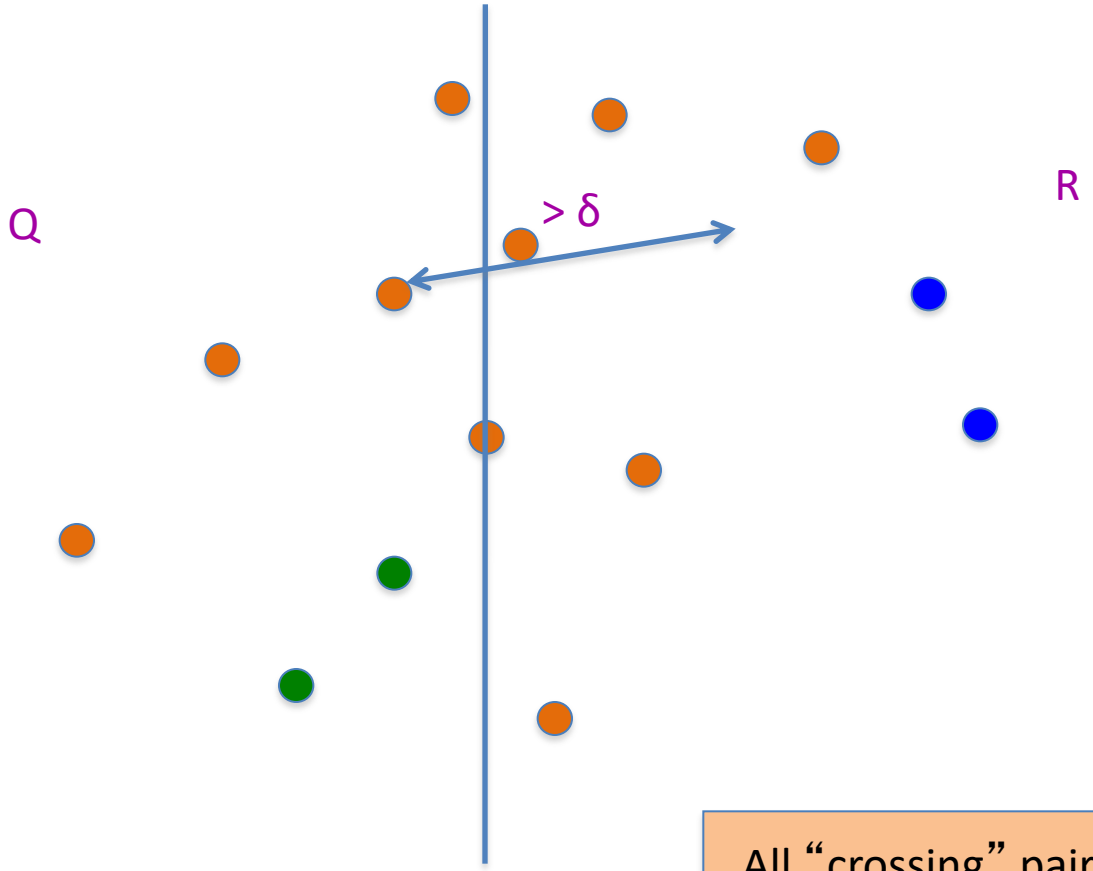
δ = min (**blue**, **green**)

An aside: maintain sorted lists

P_x and P_y are P sorted by x-coord and y-coord

Q_x, Q_y, R_x, R_y can be computed from P_x and P_y in $O(n)$ time

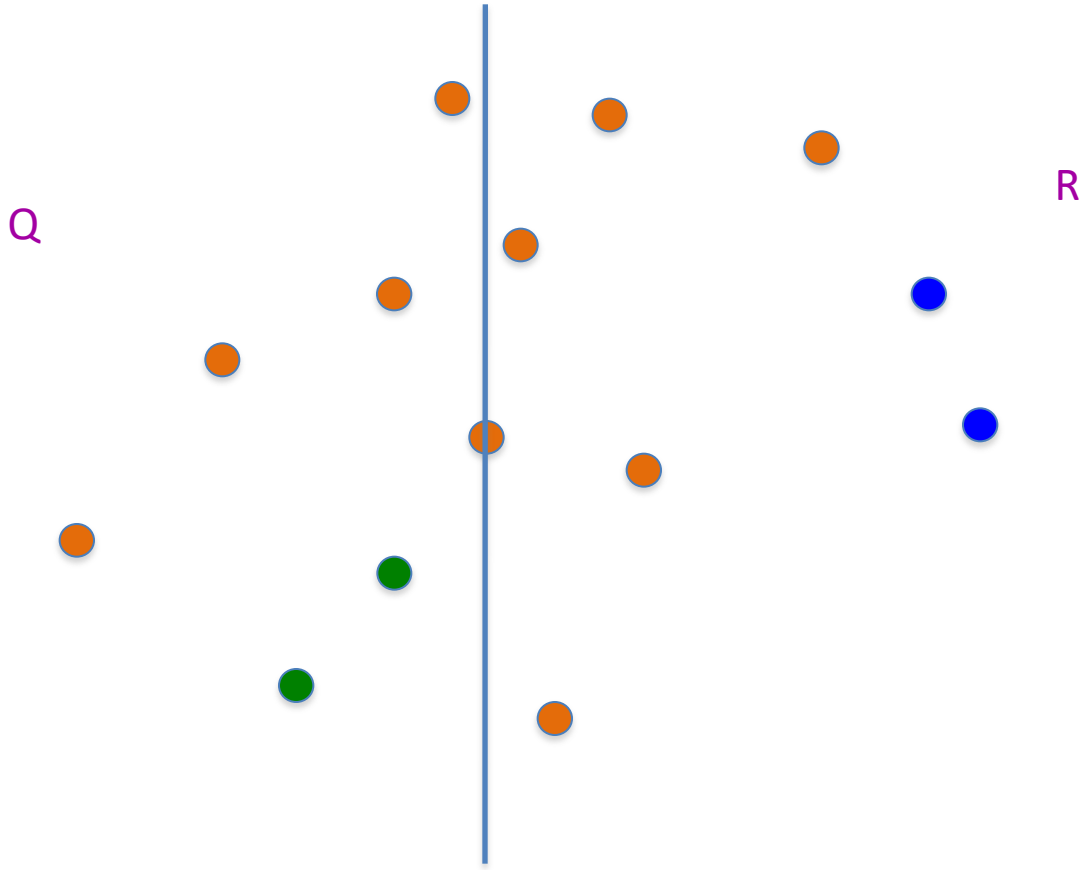
An easy case



All "crossing" pairs have distance > δ

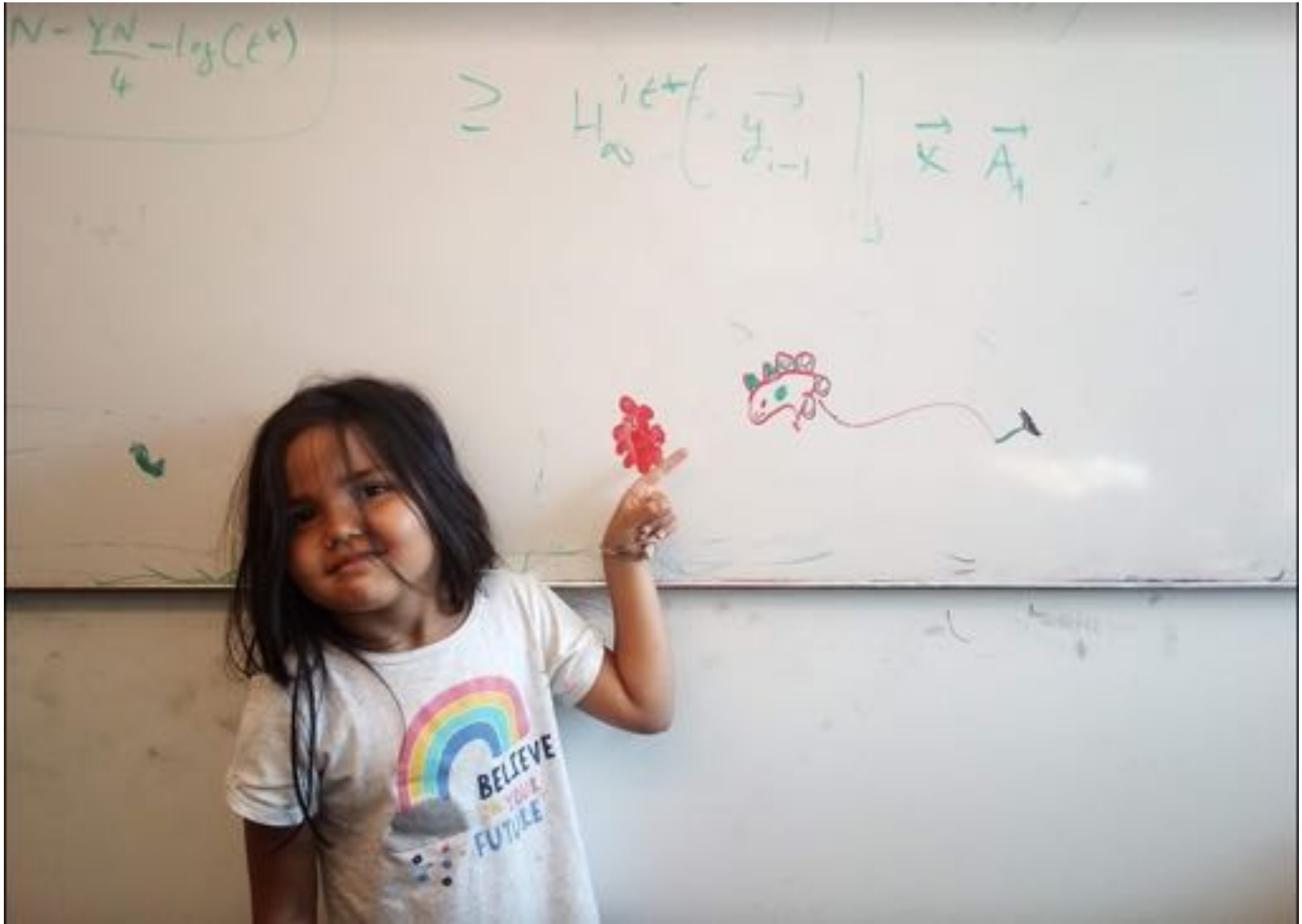
δ = min (**blue**, **green**)

Life is not so easy though



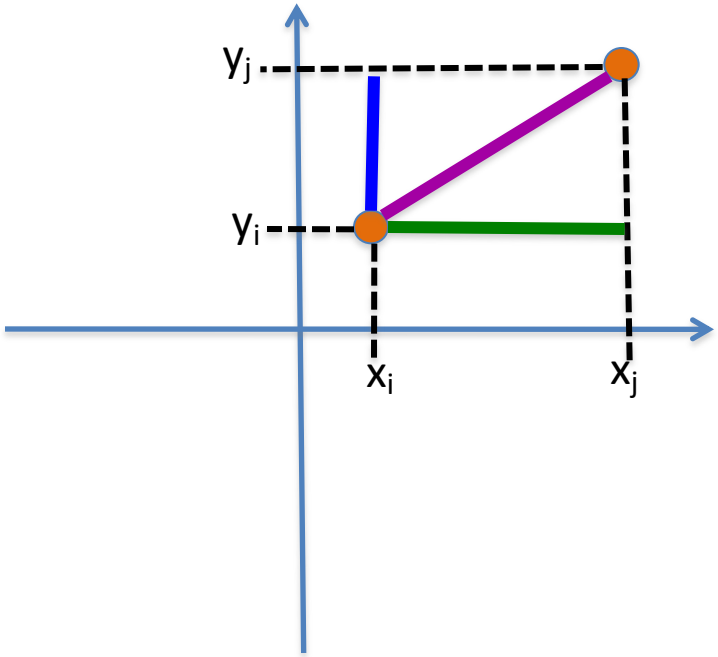
δ = min (**blue**, **green**)

Questions/Comments?



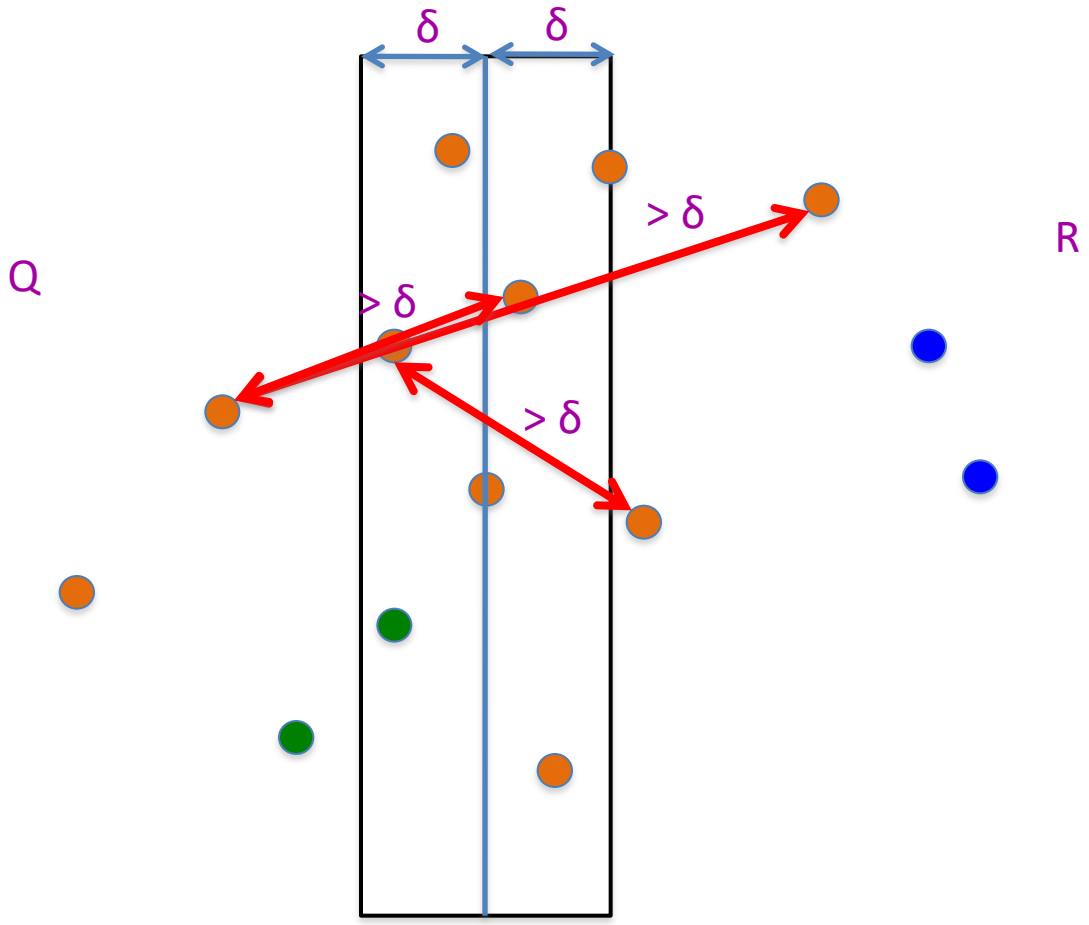
Euclid to the rescue (?)

$$d(p_i,p_j) = ((x_i-x_j)^2+(y_i-y_j)^2)^{1/2}$$



The distance is larger than the **x** or **y**-coord difference

Life is not so easy though



δ = min (**blue**, **green**)

All we have to do now

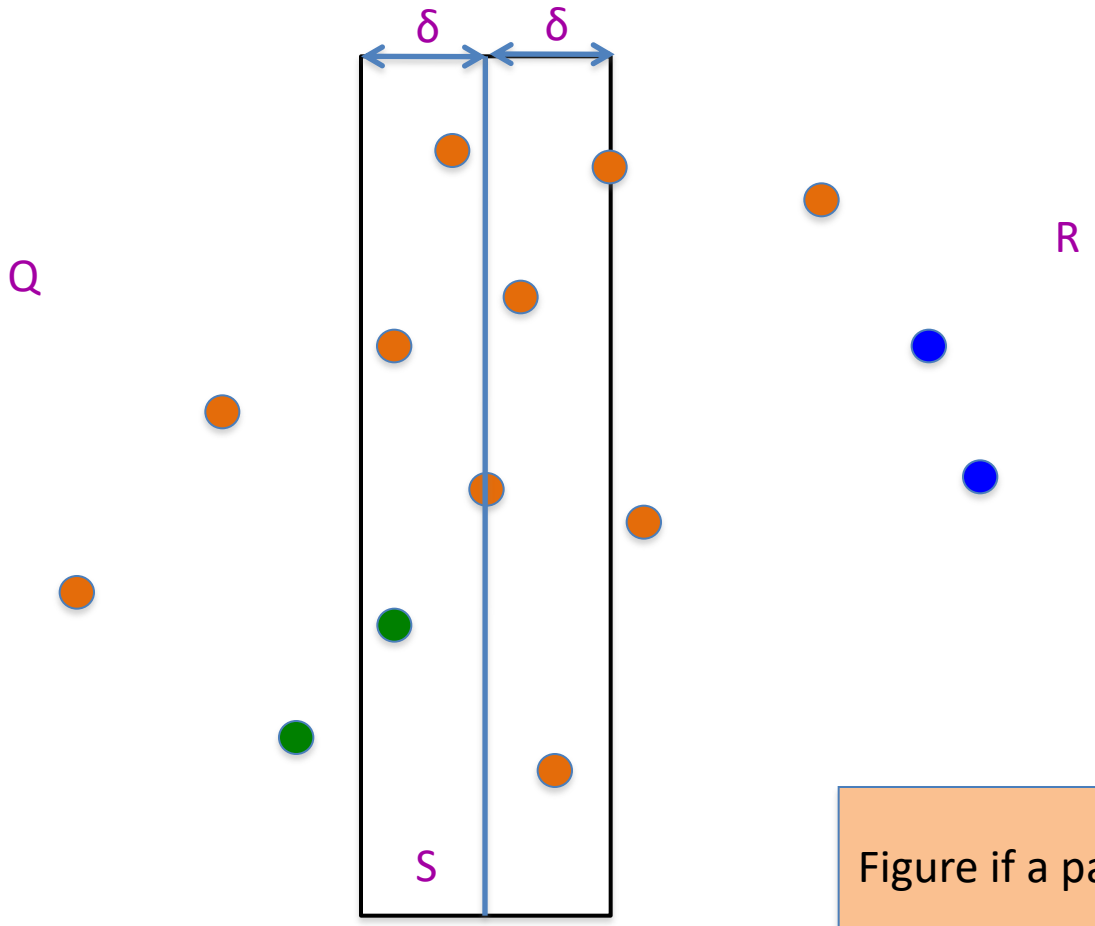


Figure if a pair in S has distance < δ

δ = min (**blue**, **green**)

The algorithm so far…

$O(n \log n) + T(n)$

Input: n 2-D points $P = \{p_1,\ldots,p_n\}$; $p_i=(x_i,y_i)$

Sort P to get P_x and P_y — $O(n \log n)$

Closest-Pair (P_x, P_y)

If $n < 4$ then find closest point by brute-force

$T(< 4) = c$

$T(n) = 2T(n/2) + cn$

Q is first half of P_x and R is the rest — $O(n)$

Compute Q_x, Q_y, R_x and R_y — $O(n)$

(q_0,q_1) = Closest-Pair (Q_x, Q_y)

(r_0,r_1) = Closest-Pair (R_x, R_y)

$O(n \log n)$ overall

$\delta = \min (d(q_0,q_1), d(r_0,r_1))$ — $O(n)$

S = points (x,y) in P s.t. $|x - x^*| < \delta$ — $O(n)$

return Closest-in-box (S, (q_0,q_1), (r_0,r_1)) — Assume can be done in $O(n)$

Rest of today's agenda

Implement Closest-in-box in O(n) time