





































Lecture 28

CSE 331

Nov 7, 2022

Reflection P2 due TODAY

Fri, Nov 4	Kickass Property Lemma     x^2	[KT, Sec 5.4] (Project (Problem 2 Coding) in)
Mon, Nov 7	Weighted Interval Scheduling    x^2	[KT, Sec 6.1] (Project (Problem 2 Reflection) in)
Tue, Nov 8		(HW 6 out)
Wed, Nov 9	Recursive algorithm for weighted interval scheduling problem    x^2	[KT, Sec 6.1]
Fri, Nov 11	Subset sum problem     x^2	[KT, Sec 6.1, 6.2, 6.4]
Mon, Nov 14	Dynamic program for subset sum     x^2	[KT, Sec 6.4]
Tue, Nov 15		(HW 7 out, HW 6 in)
Wed, Nov 16	Shortest path problem     x^2	[KT, Sec 6.8]
Fri, Nov 18	Bellman-Ford algorithm     x^2	[KT, Sec 6.8]
Mon, Nov 21	The P vs. NP problem  	[KT, Sec 8.1]
Wed, Nov 23	No class	Fall Recess
Fri, Nov 25	No class	Fall Recess
Mon, Nov 28	More on reductions  	[KT, Sec 8.1]
Tue, Nov 29		(HW 8 out, HW 7 in)
Wed, Nov 30	The SAT problem  	[KT, Sec 8.2]
Fri, Dec 2	NP-Completeness  	[KT, Sec. 8.3, 8.4] (Project (Problem 3 Coding) in)
Mon, Dec 5	k -coloring problem  	[KT, Sec 8.7] (Quiz 2) (Project (Problem 3 Reflection) in)

Group formation instructions

Autolab group submission for CSE 331 Project

The lowdown on submitting your [project](#) (especially the [coding](#) and [reflection](#)) problems as a group on Autolab.

Follow instructions **EXACTLY** as they are stated

The instructions below are for Coding Problem 1

You will have to repeat the instructions below for EACH coding AND reflection problem on project on Autolab (with the appropriate changes to the actual problem).

Form your group on Autolab

Groups on Autolab will NOT be automatically created

You will have to form a group on Autolab by yourself (as a group). Read on for instructions on how to go about this.

Make sure you are in your group

note @366

stop following 2 views

Actions

Coding P1 due today

A gentle reminder that the [first coding problem](#) is due by 11:59pm tonight!

Finally, make sure that you are officially included in your group on Autolab for the coding problem 1 *before* your group submits its code. **If you are not included in the group on Autolab, you will get a ZERO on coding problem 1.**

Please make sure that you verify that you see a submission for yourself on Autolab. It is your PERSONAL RESPONSIBILITY to make sure that this is the case. *If your group forgets to do this is it your responsibility to remind them that you need to be included.*

If your group has already submitted without you, make sure you are included in the group on Autolab and *then someone from your group should re-submit.*

project

Edit good note | 0

Updated 2 minutes ago by Atri Rudra

Reflection P1 graded

Reflection 1 graded

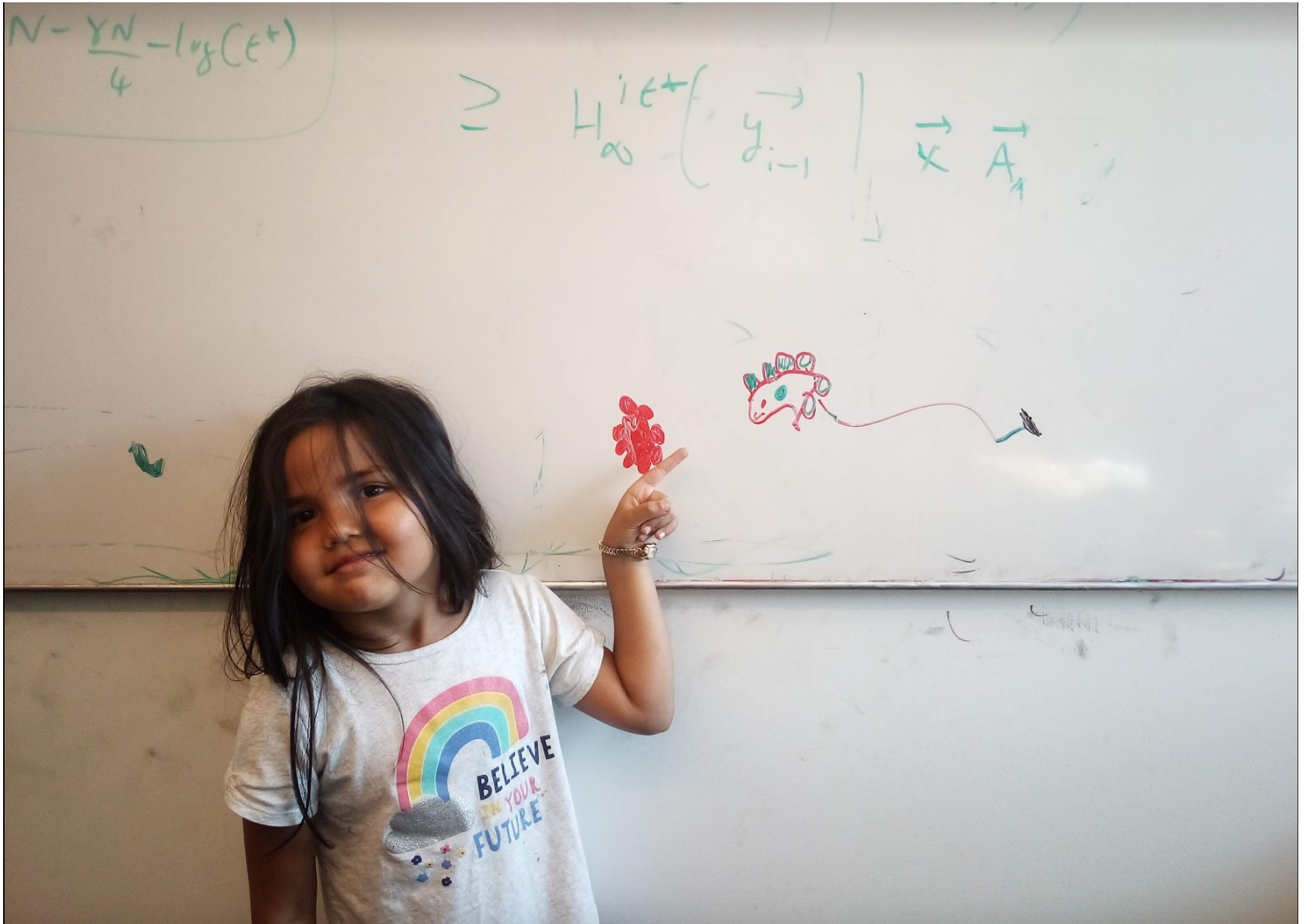
Reflection 1 grades have been released on Autolab. Hope the feedback helps as y'all prepare your reflection 2.

(Please see the re-grade policy as well as the grading rubric below before contacting us with questions on grading.)

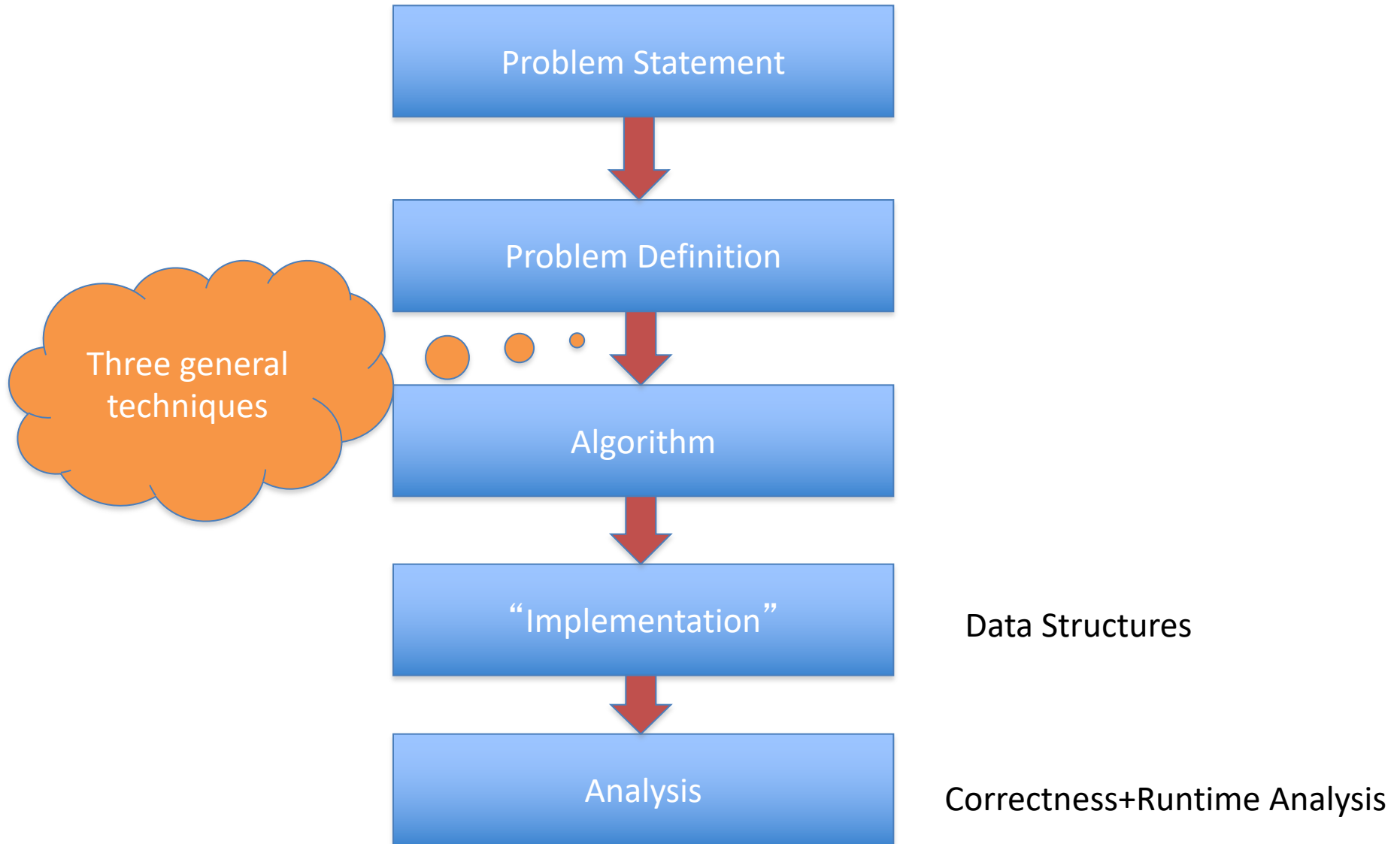
Here are the stats:

Problem	Mean	Median	StdDev	Max	Min
Stakeholder 1	0.9	1.0	0.2	1.0	0.5
Stakeholder 2	0.9	1.0	0.2	1.0	0.0
Stakeholder 3	0.9	1.0	0.2	1.0	0.0
Stakeholder 4	0.8	1.0	0.3	1.0	0.0

Questions/Comments?

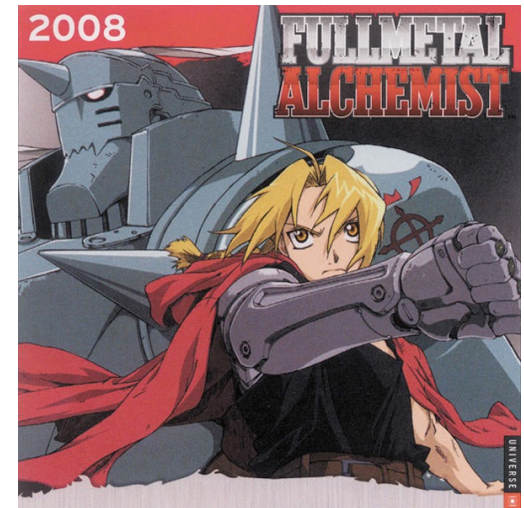
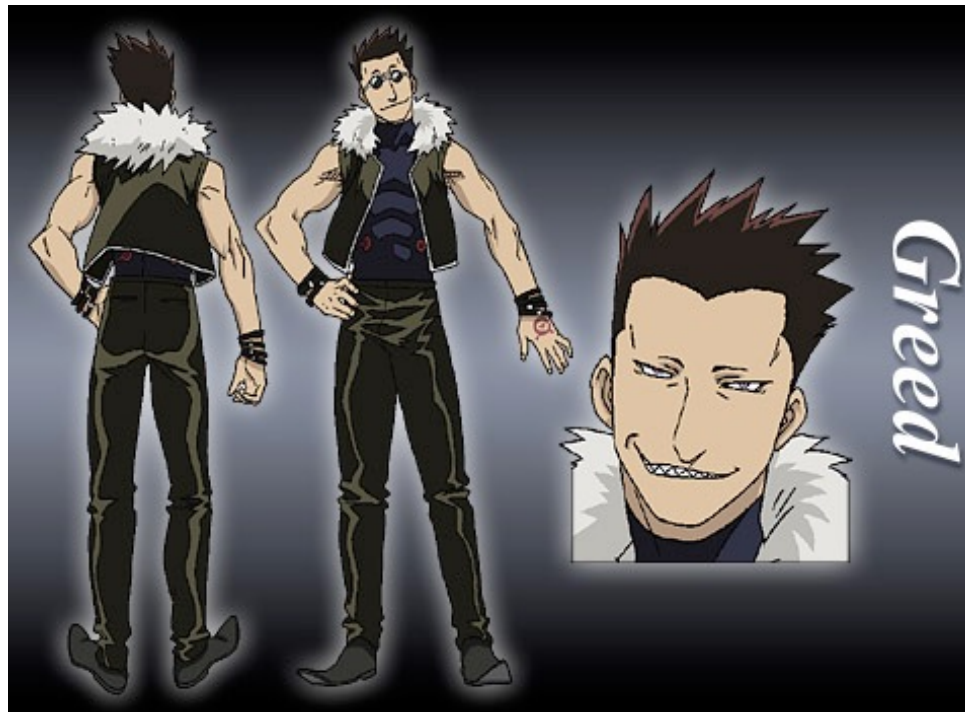


High level view of CSE 331



Greedy Algorithms

Natural algorithms



Reduced exponential running time to polynomial

Divide and Conquer

Recursive algorithmic paradigm

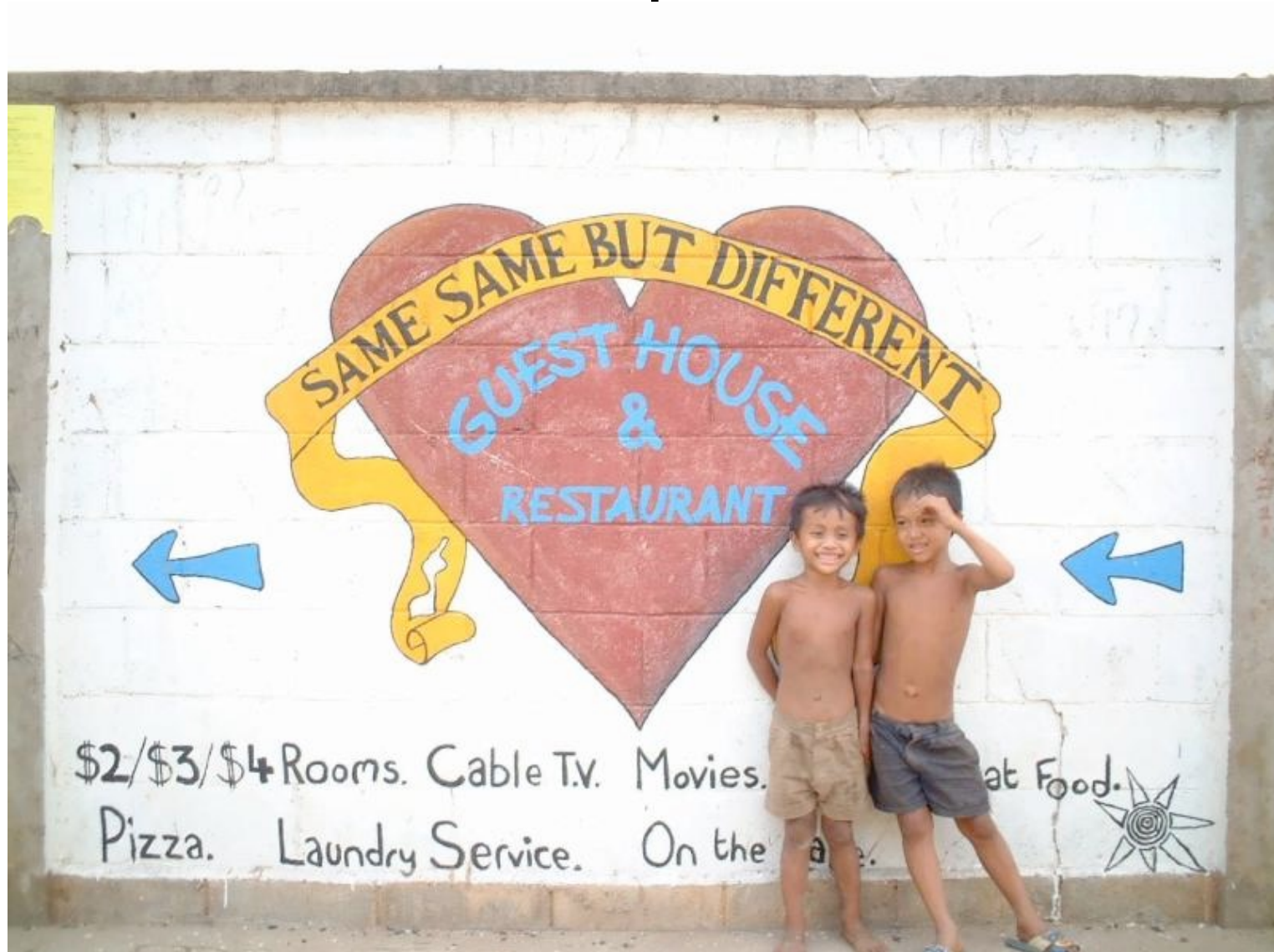


Reduced large polynomial time to smaller polynomial time

A new algorithmic technique

Dynamic Programming

Dynamic programming vs. Divide & Conquer



Same same because

Both design recursive algorithms



Different because

Dynamic programming is smarter about solving recursive sub-problems



End of Semester blues

Can only do one thing at any day: what is the optimal schedule to obtain maximum value?



Write up a term paper (10)

Party! (2)

Exam study (5)

331 HW (3)

Project (30)

Friday

Saturday

Sunday

Monday

Tuesday

Previous Greedy algorithm

Order by end time and pick jobs greedily

Greedy value = $5+2+3=10$

Write up a term paper (10)

Party! (2)

Exam study (5)

331 HW (3)

Project (30)

OPT = 30



Friday

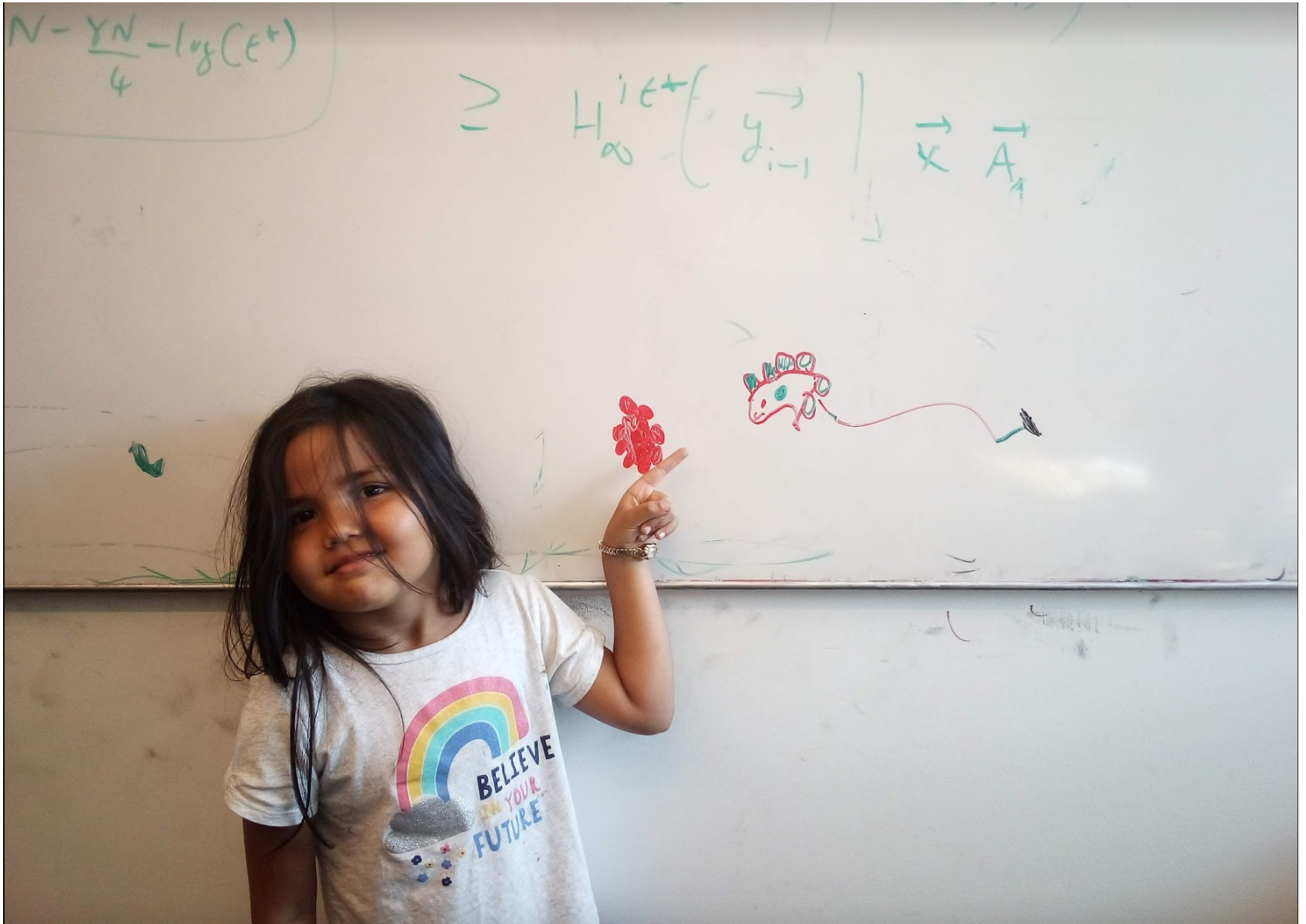
Saturday

Sunday

Monday

Tuesday

Questions/Comments?



Today's agenda

Formal definition of the problem

Start designing a recursive algorithm for the problem



Weighted Interval Scheduling

Input: n jobs/intervals. Interval i is triple (s_i, f_i, v_i)

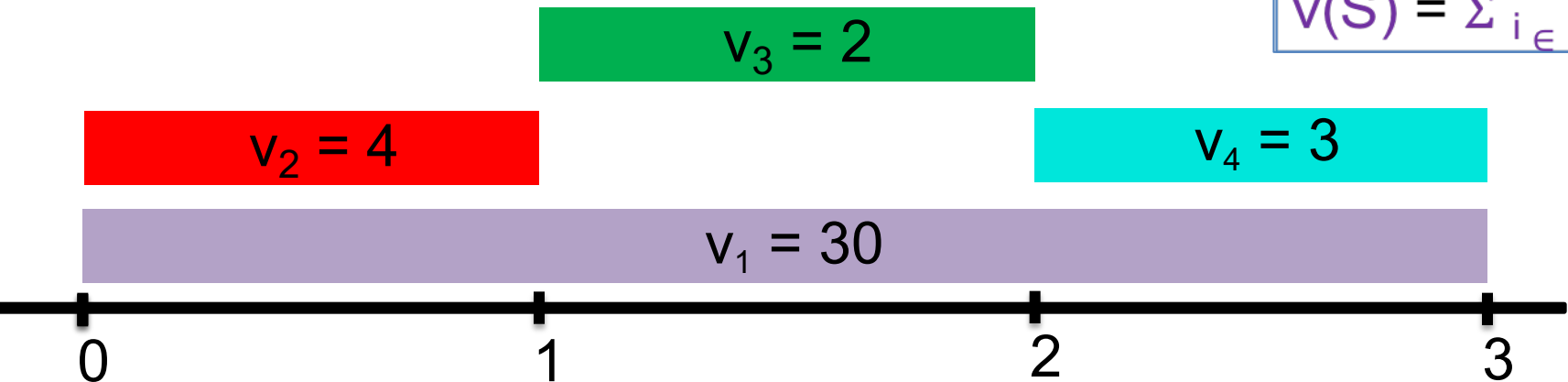
start time

finish time

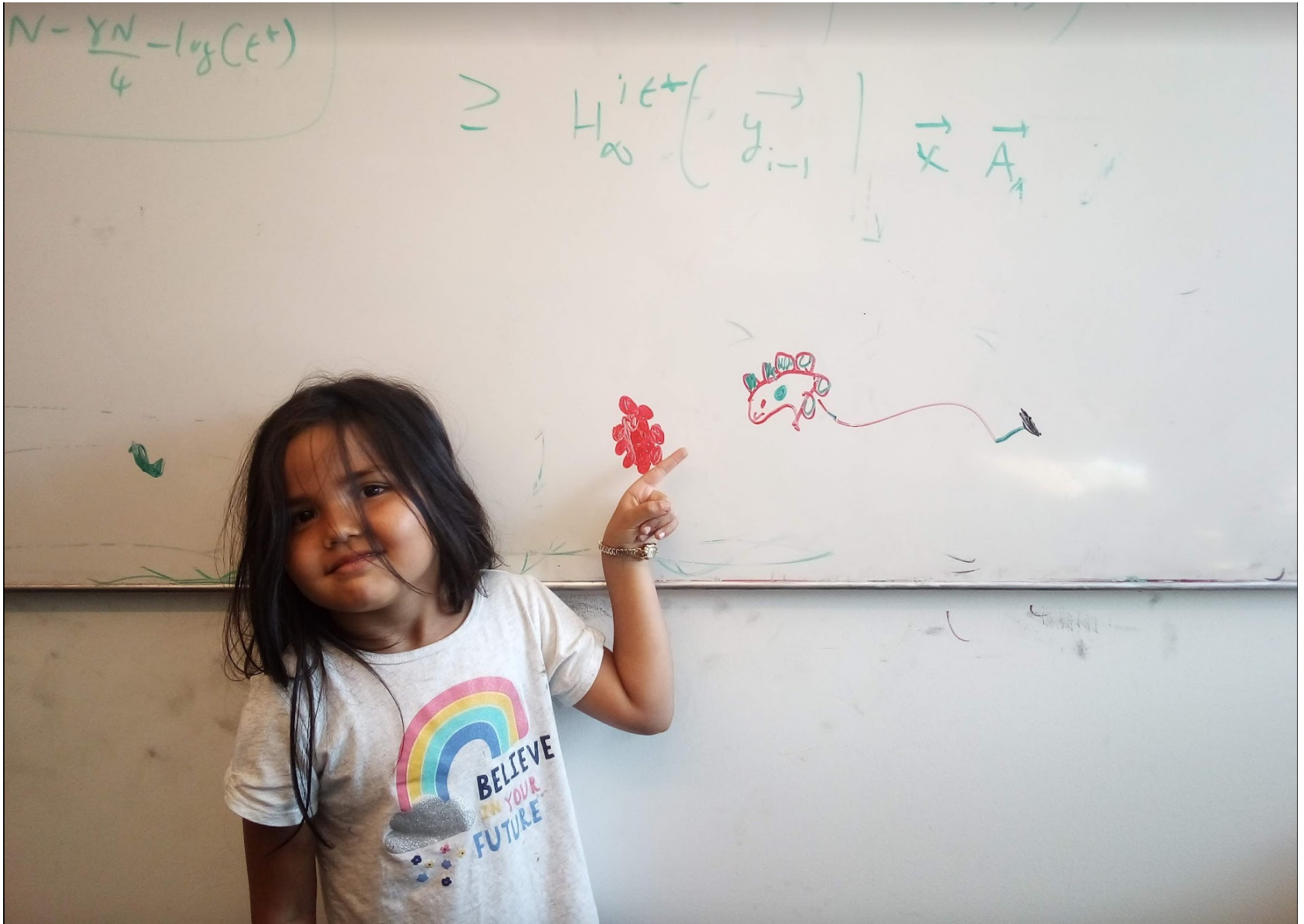
value

Output: A valid schedule $S \subseteq [n]$ that maximizes $v(S)$

$$v(S) = \sum_{i \in S} v_i$$



Questions/Comments?



Previous Greedy Algorithm

R = original set of jobs

$S = \phi$

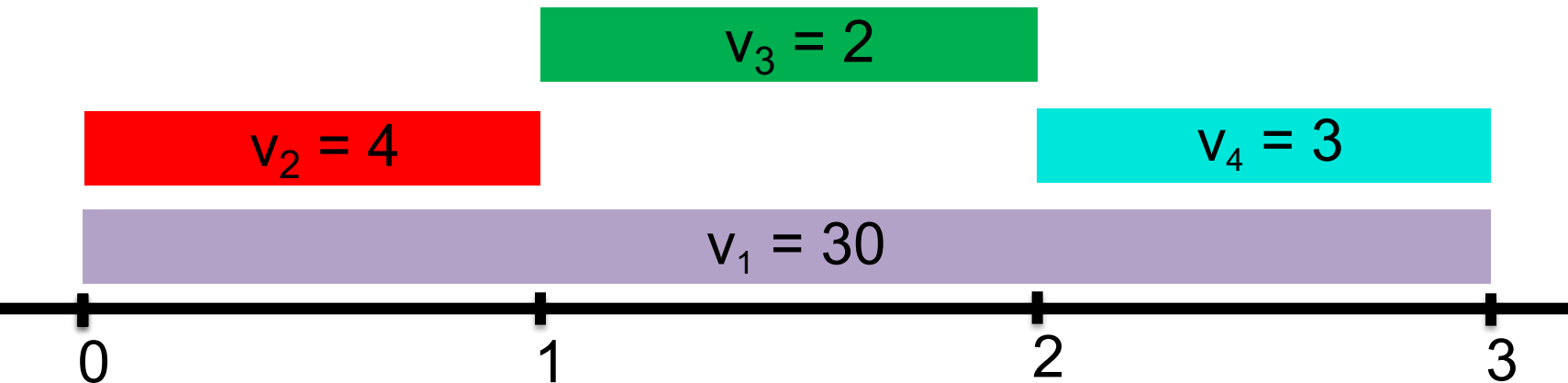
While R is not empty

 Choose i in R where f_i is the smallest

 Add i to S

 Remove all requests that conflict with i from R

Return $S^* = S$



Perhaps be greedy differently?

R = original set of jobs

$S = \phi$

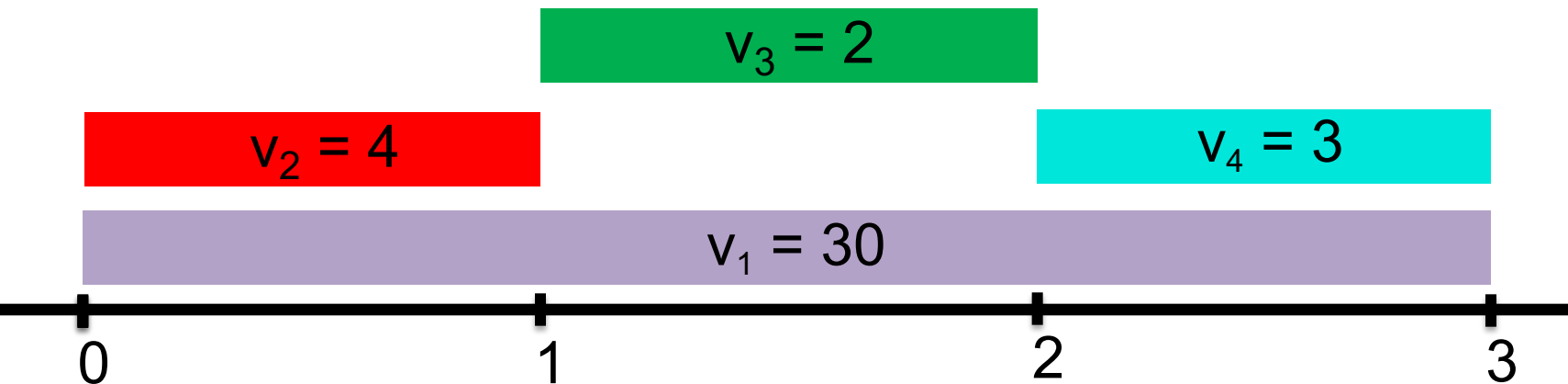
While R is not empty

Choose i in R where $v_i/(f_i - s_i)$ is the largest

Add i to S

Remove all requests that conflict with i from R

Return $S^* = S$



Can this work?

R = original set of jobs

$S = \phi$

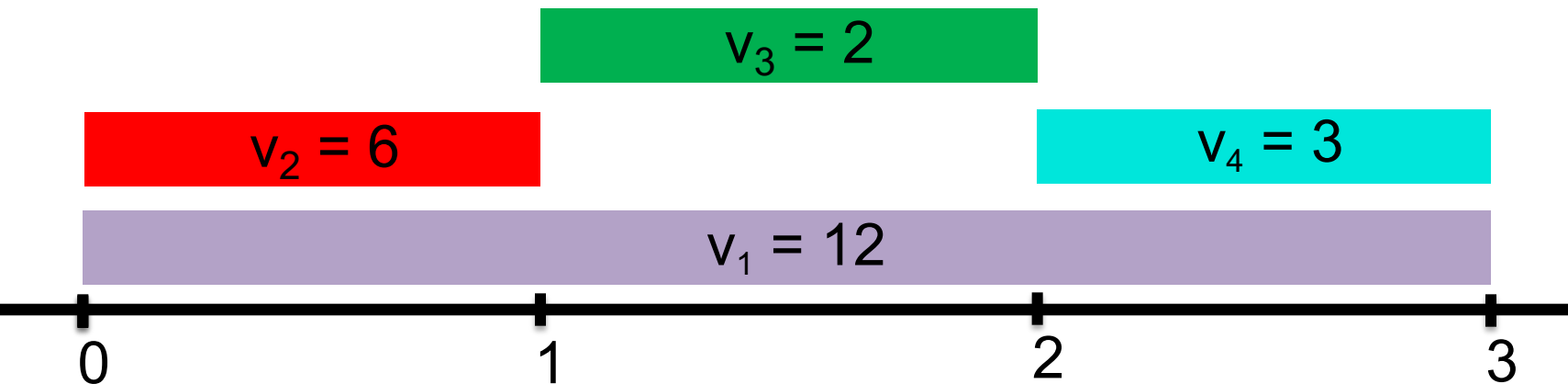
While R is not empty

Choose i in R where $v_i/(f_i - s_i)$ is the largest

Add i to S

Remove all requests that conflict with i from R

Return $S^* = S$



Avoiding the greedy rabbit hole

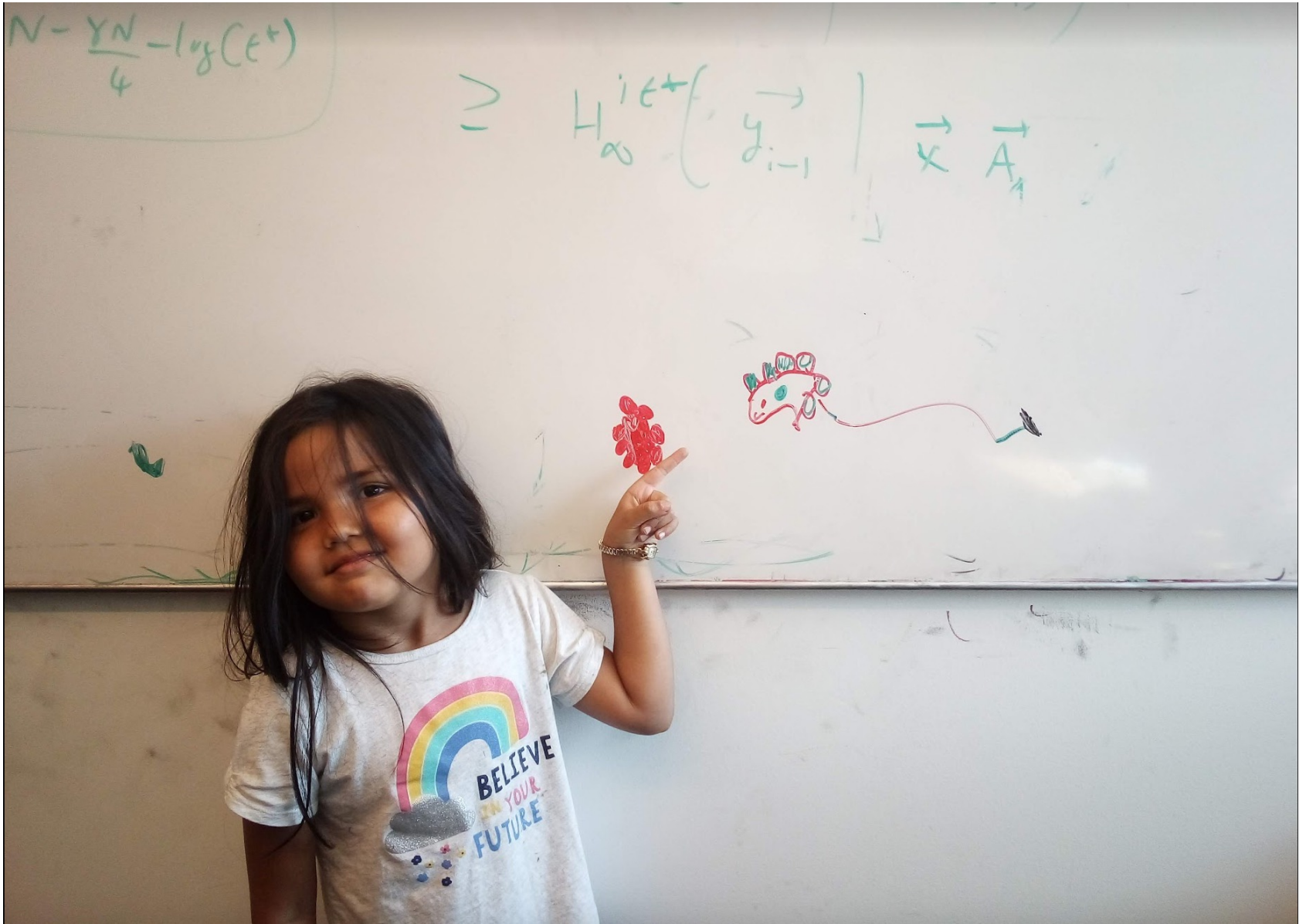


<https://www.wrightwords.com/down-the-rabbit-hole/>

Provably
IMPOSSIBLE
for a large
class of
greedy algos

There are no known greedy algorithm to solve this problem

Questions/Comments?



Perhaps a divide & conquer algo?

Divide the problem in 2 or more many EQUAL SIZED
INDEPENDENT problems

Recursively solve the sub-problems

Patchup the SOLUTIONS to the sub-problems

Perhaps a divide & conquer algo?

RecurWeightedInt([n])

if $n = 1$ return the only interval

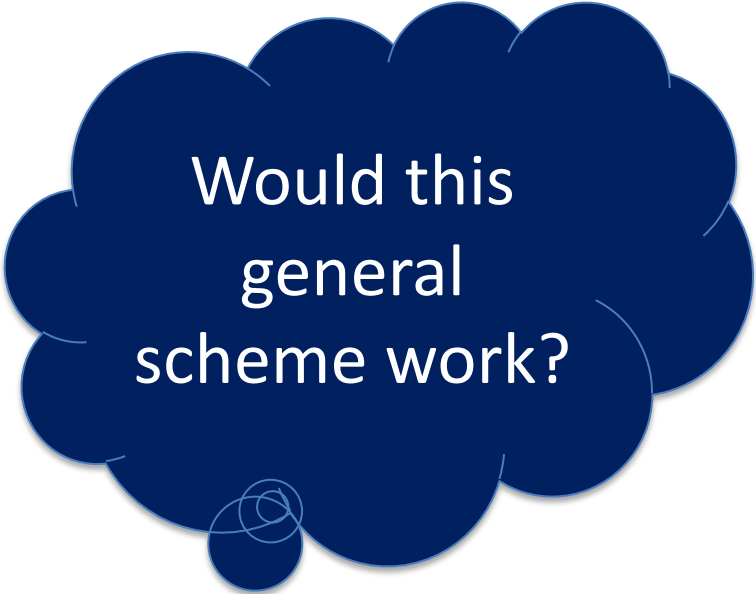
L = first $n/2$ intervals

R = last $n/2$ intervals

S_L = RecurWeightedInt(L)

S_R = RecurWeightedInt(R)

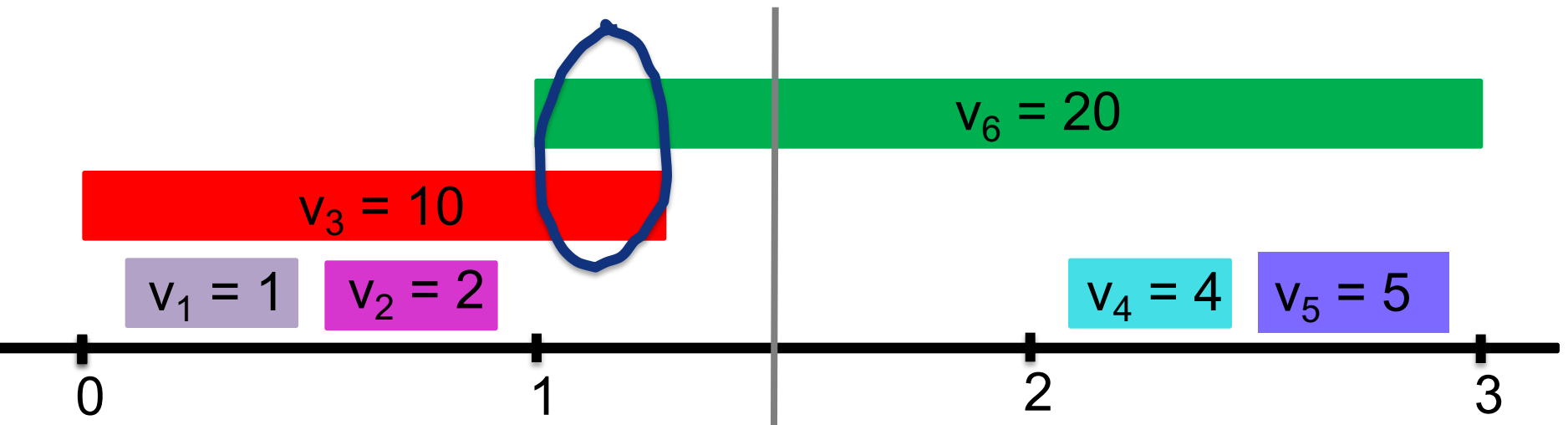
PatchUp(S_L , S_R)



Would this
general
scheme work?

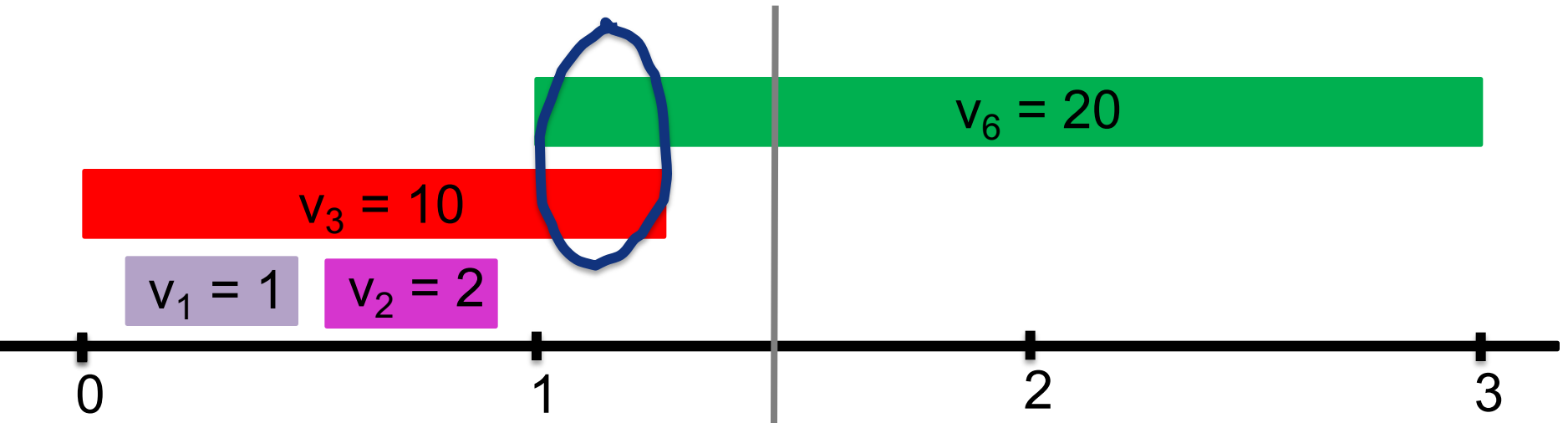
Divide the problem in 2 or more many EQUAL SIZED
INDEPENDENT problems

Sub-problems NOT independent!

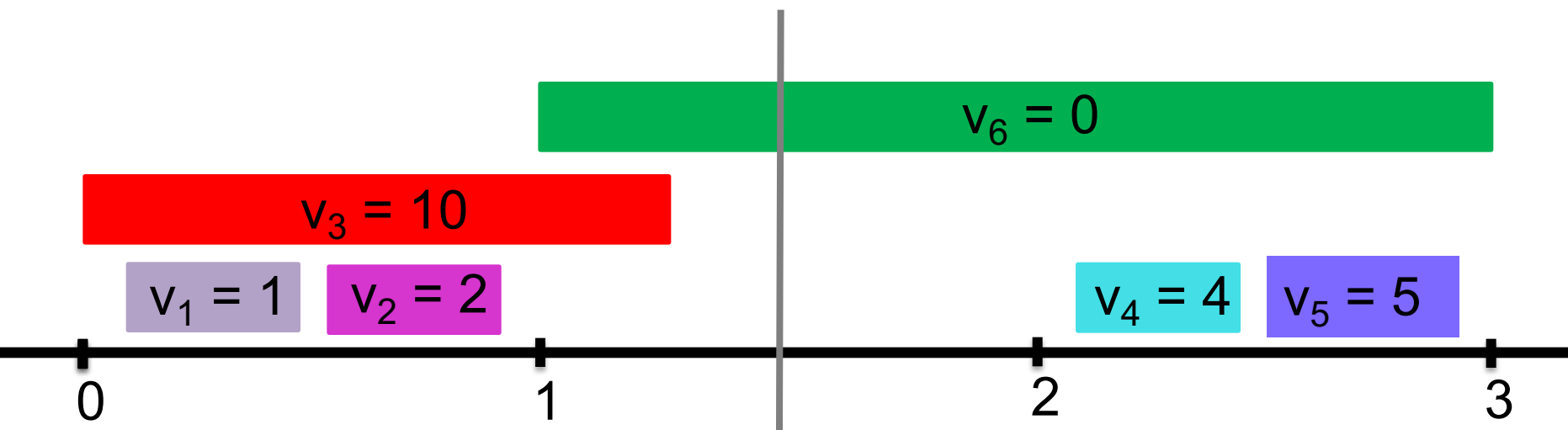


Perhaps patchup can help?

Patchup the SOLUTIONS to the sub-problems



Sometimes patchup NOT needed!



Check for two cases?

6 is in the optimal solution

$$v_6 = 20$$

$$v_3 = 10$$

$$v_1 = 1$$

$$v_2 = 2$$

$$v_4 = 4$$

$$v_5 = 5$$

6 is NOT in the optimal solution

$$v_6 = 0$$

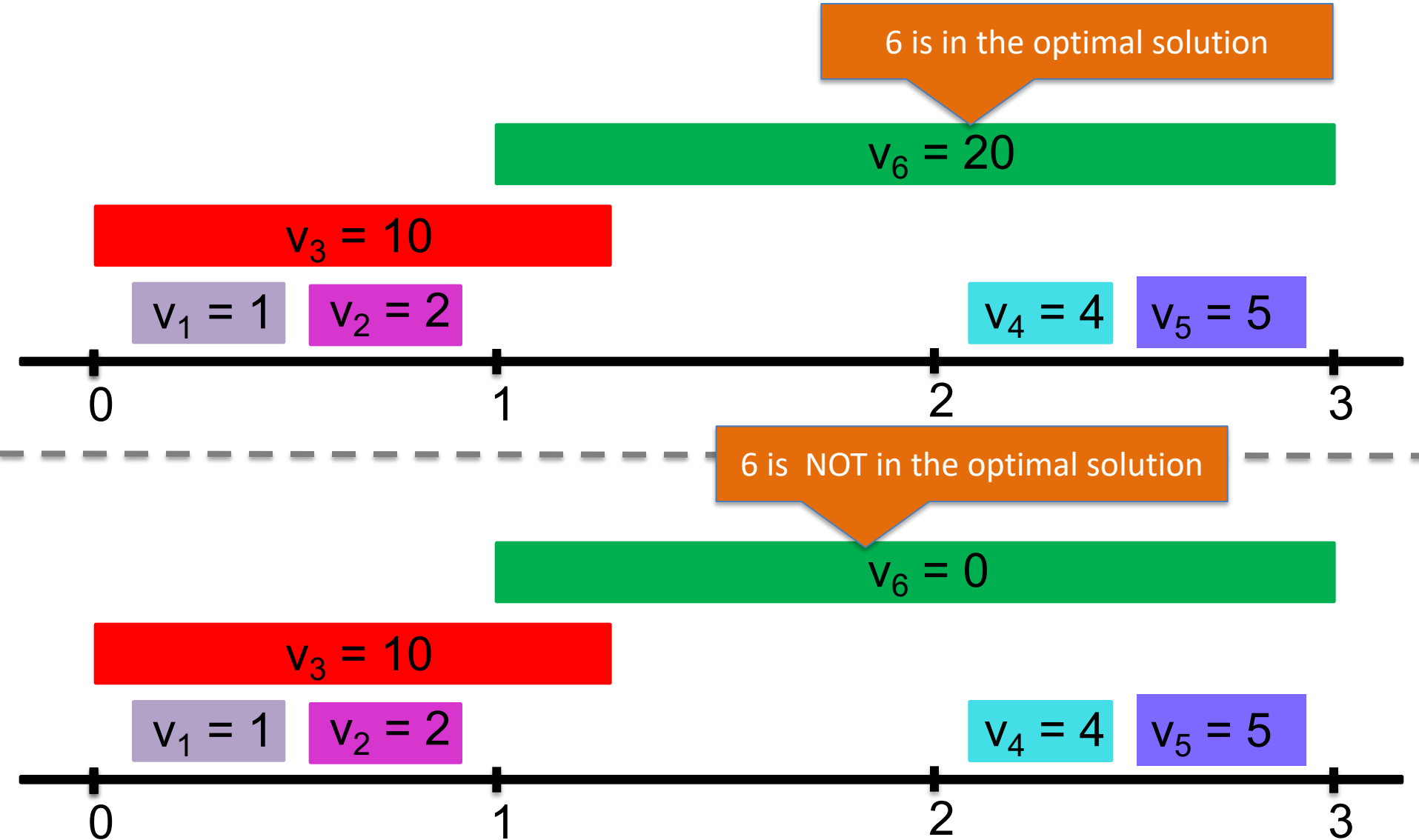
$$v_3 = 10$$

$$v_1 = 1$$

$$v_2 = 2$$

$$v_4 = 4$$

$$v_5 = 5$$



Check if v_6 is the largest value?

6 is in the optimal solution

$v_6 = 20$

$v_3 = 10$

$v_1 = 1$

$v_2 = 2$

$v_4 = 4$

$v_5 = 5$

Cannot decide this greedily. Need to have a global view!

al solution

$v_6 = 20$

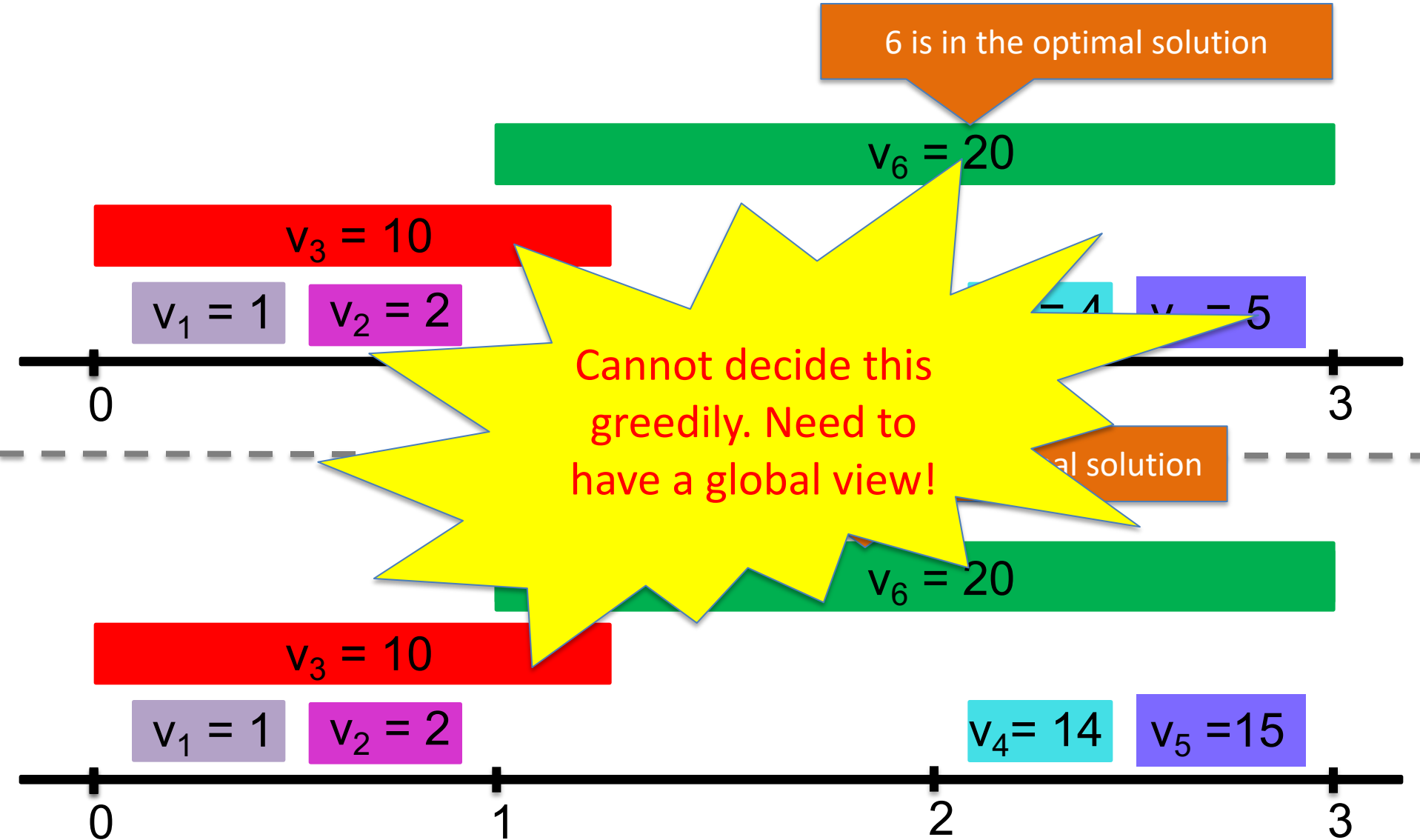
$v_3 = 10$

$v_1 = 1$

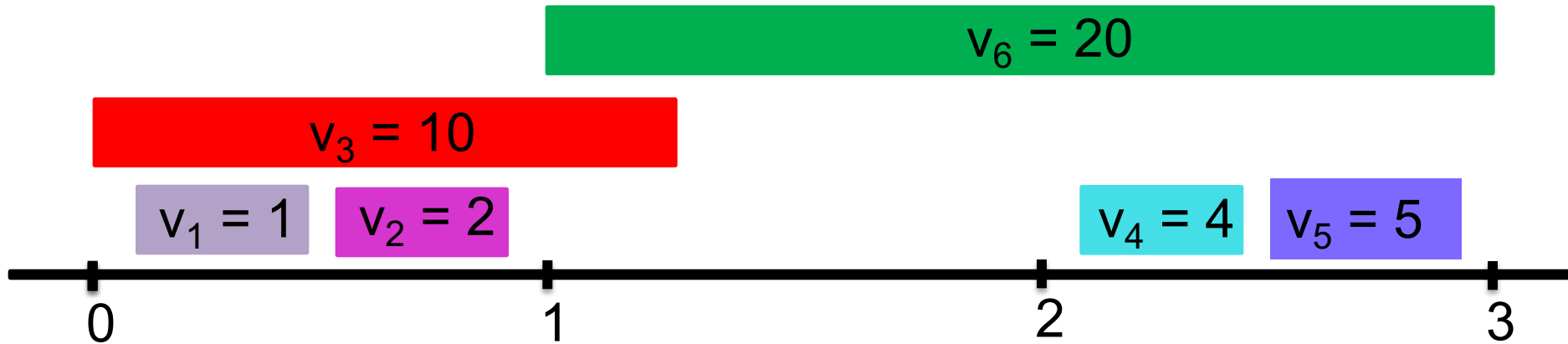
$v_2 = 2$

$v_4 = 14$

$v_5 = 15$



Check out both options!



Case 1: 6 is in the optimal solution

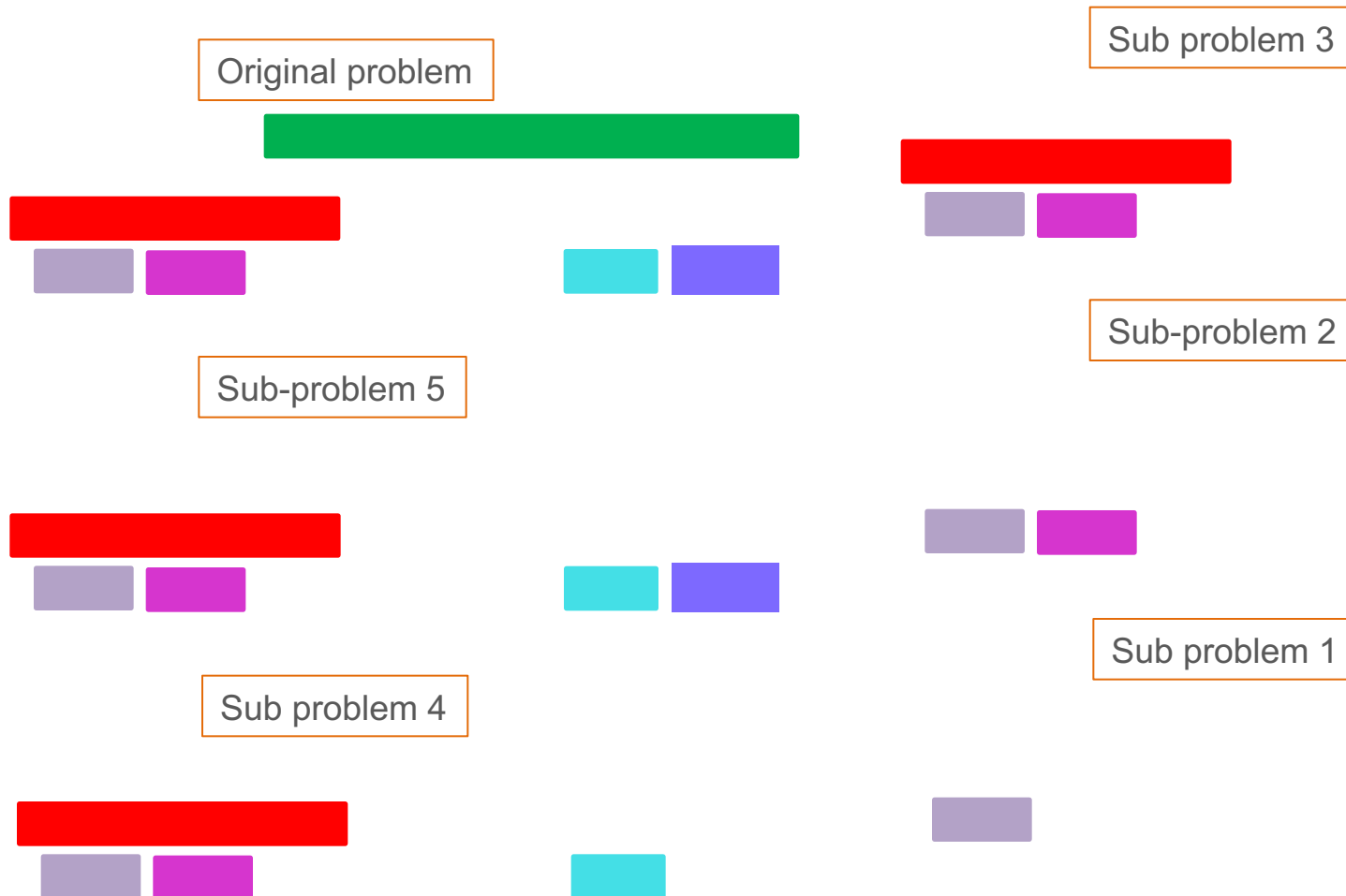
6 is not in optimal solution





So what sub-problems?

Divide the problem in 2 or more many ~~EQUAL SIZED~~
~~INDEPENDENT~~ problems





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