

Lecture 29

CSE 331

Nov 9, 2022

Homework 6 out

Homework 6

Due by **11:30pm, Tuesday, November 15, 2022.**

Make sure you follow all the [homework policies](#).

All submissions should be done via [Autolab](#).

Question 1 (Exponentiation) [50 points]

The Problem

We will consider the problem of exponentiating an integer to another. In particular, for non-negative integers a and n , define `Power` (a, n) be the number a^n . (For this problem assume that you can multiply two integers in $O(1)$ time.) Here are the two parts of the problem:

- Part **(a)**: Present a naive algorithm that given non-negative integers a and n computes `Power` (a, n) in time $O(n)$.

Note

For this part, there is no need to prove correctness of the naive algorithm but you do need a runtime analysis.

- Part **(b)**: Present a divide and conquer algorithm that given non-negative integers a and n computes `Power` (a, n) in $O(\log n)$ time.

Important Note

To get credit you must present a recursive divide and conquer algorithm and then analyze its running time by solving a recurrence relation. If you present an algorithm that is not a divide and conquer algorithm you will get a level 0 on this entire part.

Questions/Comments?



End of Semester blues

Can only do one thing at any day: what is the optimal schedule to obtain maximum value?



Write up a term paper (10)

Party! (2)

Exam study (5)

331 HW (3)

Project (30)

Friday

Saturday

Sunday

Monday

Tuesday

Previous Greedy algorithm

Order by end time and pick jobs greedily

Greedy value = $5+2+3=10$

Write up a term paper (10)

Party! (2)

Exam study (5)

331 HW (3)

Project (30)

OPT = 30



Friday

Saturday

Sunday

Monday

Tuesday

Weighted Interval Scheduling

Input: n jobs (s_i, f_i, v_i)

Output: A schedule S s.t. no two jobs in S have a conflict

Goal: $\max \sum_{i \in S} v_j$

Assume: jobs are sorted by their finish time

Today's agenda

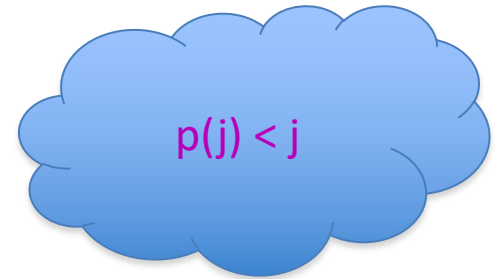
Finish designing a recursive algorithm for the problem



Couple more definitions

$p(j)$ = largest $i < j$ s.t. i does not conflict with j

= 0 if no such i exists



$OPT(j)$ = optimal value on instance $1, \dots, j$

Moving to the board...



Property of OPT

j in $\text{OPT}(j)$

j not in $\text{OPT}(j)$

$$\text{OPT}(j) = \max \{ v_j + \text{OPT}(p(j)), \text{OPT}(j-1) \}$$

Given $\text{OPT}(1), \dots, \text{OPT}(j-1)$,
how can one figure out if j
in optimal solution or not?



A recursive algorithm

Compute-Opt(j)

Correct for $j=0$

Proof of
correctness by
induction on j

If $j = 0$ then return 0

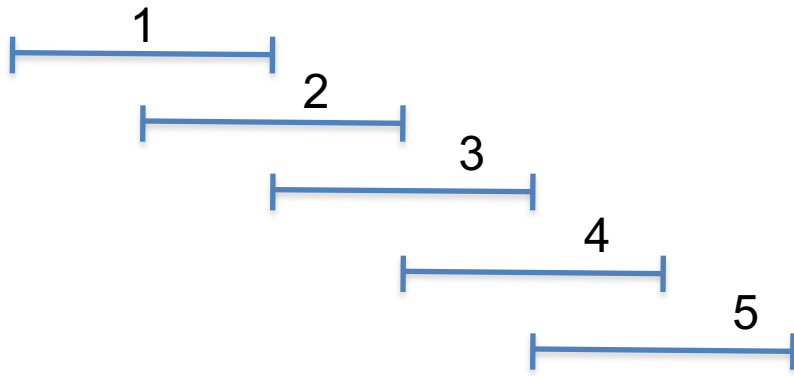
return max { $v_j + \text{Compute-Opt}(p(j))$, $\text{Compute-Opt}(j-1)$ }

= OPT($p(j)$)

= OPT($j-1$)

$$\text{OPT}(j) = \max \{ v_j + \text{OPT}(p(j)), \text{OPT}(j-1) \}$$

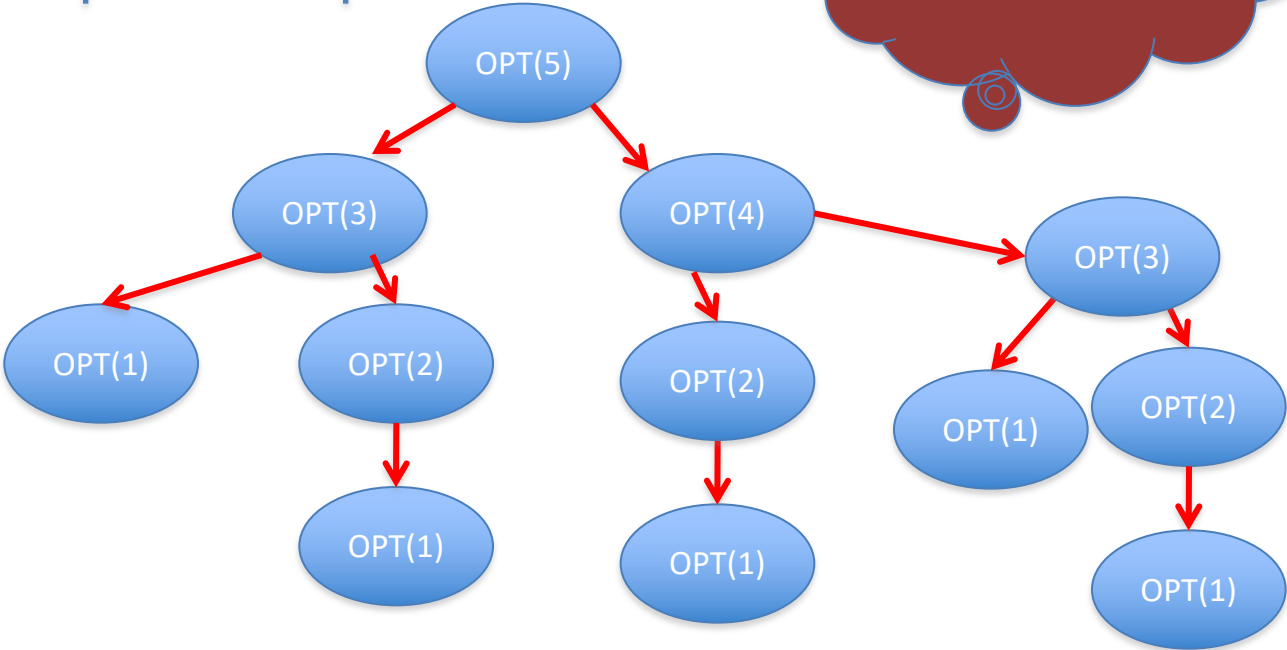
Exponential Running Time



$$p(j) = j - 2$$

Only 5 OPT values!

Formal proof: Ex.





Using Memory to be smarter

Using more space can reduce runtime!

How many distinct OPT values?

A recursive algorithm

M-Compute-Opt(j)

If $j = 0$ then return 0

If $M[j]$ is not null then return $M[j]$

$M[j] = \max \{ v_j + \text{M-Compute-Opt}(p(j)), \text{M-Compute-Opt}(j-1) \}$

return $M[j]$

M-Compute-Opt(j)
= OPT(j)

Run time = $O(\# \text{ recursive calls})$