

Sep 26

Explore (s) in G

$\Delta \equiv S$



- $R \leftarrow \{s\}$
- While $\exists (u, w) \in E$ s.t. $w \notin R, u \in R$
Add w to R
- Output $R^* \leftarrow R$

such that

THEOREM: For all $G=(V, E)$, for all $s \in V$, $R^* = CC(s)$
If R^* is the output of Explore (s) in G , then $R^* = CC(s)$

COROLLARY: BFS is correct

BFS is a special case of Explore

General idea: To show that two sets A & B are $A=B \iff (i) A \subseteq B \text{ (ii) } B \subseteq A$

Lemma 1: $R^* \subseteq CC(s)$ ← everything that is output by Explore is correct.

Lemma 2: $CC(s) \subseteq R^*$ ← everything that should be output by Explore (i.e. in $CC(s)$) is actually output by Explore

Lemmas 1+2 \implies THM

Pf of Lemma 1: Ex (by induction)

Pf (idea) of Lem 2: Pf by contradiction

Assume $CC(s) \not\subseteq R^*$

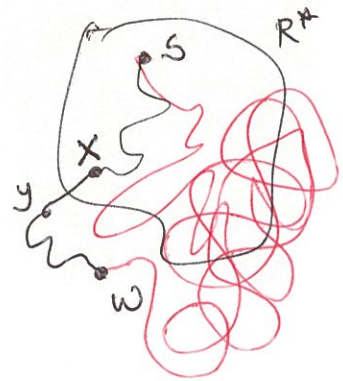
$\implies \exists w \in CC(s)$ s.t. $w \notin R^*$

$\iff \exists$ an $s-w$ path P in G but $w \notin R^*$

Since P starts inside of R^* ($s \in R^*$) but ends up outside of R^* ($w \notin R^*$)

$\implies P$ has to "cross" R^* at some point

Note: Since R^* is output \implies Explore has terminated (#)



$\Rightarrow \exists (x, y) \in E$ s.t. $x \in R^*$ but $y \notin R^*$

\Rightarrow
by algo
defs

y should have been added by Explore to R

\Rightarrow Explore should not have terminated with R^*
 \rightarrow contradiction (#)