

Lecture 26

CSE 331

Nov 3, 2023

Coding Project 1+2 due today

Tue, Oct 31		(HW 5 in)
Wed, Nov 1	Multiplying large integers  F22  F21  F19  F18  F17 x ²	[KT, Sec 5.5] <i>Reading Assignment:</i> Unraveling the mystery behind the identity
Fri, Nov 3	Closest Pair of Points  F22  F21  F19  F18  F17 x ²	[KT, Sec 5.4] (Project (Problems 1 & 2 Coding) in)
Mon, Nov 6	Kickass Property Lemma  F22  F21  F19  F18  F17 x ²	[KT, Sec 5.4] (Project (Problems 1 & 2 Reflection) in)

Group formation instructions

Autolab group submission for CSE 331 Project

The lowdown on submitting your [project](#) (especially the [coding](#) and [reflection](#)) problems as a group on Autolab.

Follow instructions **EXACTLY** as they are stated

The instructions below are for Coding Problem 1

You will have to repeat the instructions below for EACH coding AND reflection problem on project on Autolab (with the appropriate changes to the actual problem).

Form your group on Autolab

Groups on Autolab will NOT be automatically created

You will have to form a group on Autolab by yourself (as a group). Read on for instructions on how to go about this.

Do not use Join a Group “feature”

note @331 stop following 143 views Actions

When forming groups on Autolab

I forgot to add this warning earlier but I have updated the [Autolab project page](#) to add a warning to **not click on the "Join a Group" function when you are creating your group on Autolab:**

Do NOT click on `Join a Group`

Do NOT use the "Join a Group" feature. ONLY follow the instructions above EXACTLY.

This step can be un-done but needs intervention on our part BUT that'll cause delays on your side and we are not responsible if you miss your deadline due to this delay.

Here is what such an option looks like (the actual group name and group members would be different in your case):

Join a Group ← Do NOT use this option

TAs_(Java_testers)

Gitanjali Nandi, Thomas Sherwood

[Ask to Join Group](#)

project autolab

No project group => no submission

note @432

stop following 13 views

Actions

You can only submit if you have an official group

If for whatever reason you did not create or sign up to be on a random group by the deadline (and you did not get an email from me confirming your group), then you **cannot** submit anything for the project.

I will be double-checking the groups on Autolab with the official group after the deadline. If you formed a group even though you do not have an official group that would be considered an AI violation.

Also make sure that you are in a group before the *final* submission by your group. If your group member submits without adding you to a group on Autolab, you will get a **zero**.

project

Edit good note | 0

Updated 0 seconds ago by Atri Rudra

1-on-1 meeting slots

note @392

stop following 33 views

Actions

Meetings to discuss CSE 331 performance

By Sunday tonight, I will email those who have a D+ or below in their mid-term grade (for more details on the grade see @393) to setup a one-on-one meeting to talk with me but I figured I should post the information about meeting times now rather than tomorrow.

Of course you can also come and talk about your 331 performance even if you have a temp grade higher than D+ (though students with a D+ or below will get preference).

I have locked out certain times over next week or so for **15 mins** meetings. Please note that **these are NOT walk-ins**: if no one signs up for a slot, I will NOT be on zoom then. If you want to come and talk with me, **please EMAIL me with ALL the slots below that work for you**. (Private posts on piazza will not work: please email me!) *Slots will be assigned on a first-come-first-serve basis. Also I might only be able to confirm your time after 11pm on the day before your scheduled slot.*

Note: These are my current availabilities-- some of the slots might be used up in some other non-CSE 331 meetings. So please send multiple choices for when you can meet.

To make things easier, **ALL meeting will be on zoom** (<https://buffalo.zoom.us/j/96902087672?pwd=UXVxL21OQkdLYWd1VzdhdHFNbmIPdz09>)

Below are all the available slots (below the start times are listed: a slot that is already taken has a strike-through):

- **Monday (Oct 30):** 9:00am, 9:15, 9:30, 9:45, 10:00, 12:30pm, 2:30, 2:45, 5:15, 5:30, 5:45, 6:00, 6:15, 6:30, 6:45
- **Tuesday (Oct 31):** 9:45am, 10:00am, 2:45pm
- **Wednesday (Nov 1):** 9:00am, 9:15, 9:30, 9:45, 10:00, 1:00pm, 1:15, 5:15, 5:30, 5:45, 6:00, 6:15, 6:30
- **Thursday (Nov 2):** 9:00am, 9:15, 9:30, 9:45, 10:00, 10:15, 10:30, 10:45, 11:00, 11:15, 11:30, 11:45, 12:00pm, 12:15, 12:30, 12:45, 1:00, 1:15
- **Friday (Nov 3):** 9:00am, 9:15, 9:30, 9:45, 10:00, 12:30pm

(Apologies but my schedule for this semester is bit of a mess. If none of the times above work for you but you still want to meet, please email me and we can set up a time for the week of Nov 5.)

You can of course also stop by during my office hours (but students with Qs on the HWs will get higher priority) and you unfortunately cannot book a slot during my usual office hours.

Again, **please email me your (at least top 3) choices** (again note the ALL slots are virtual).

mid-term grading

Edit good note 0

Updated 21 minutes ago by Atri Rudra

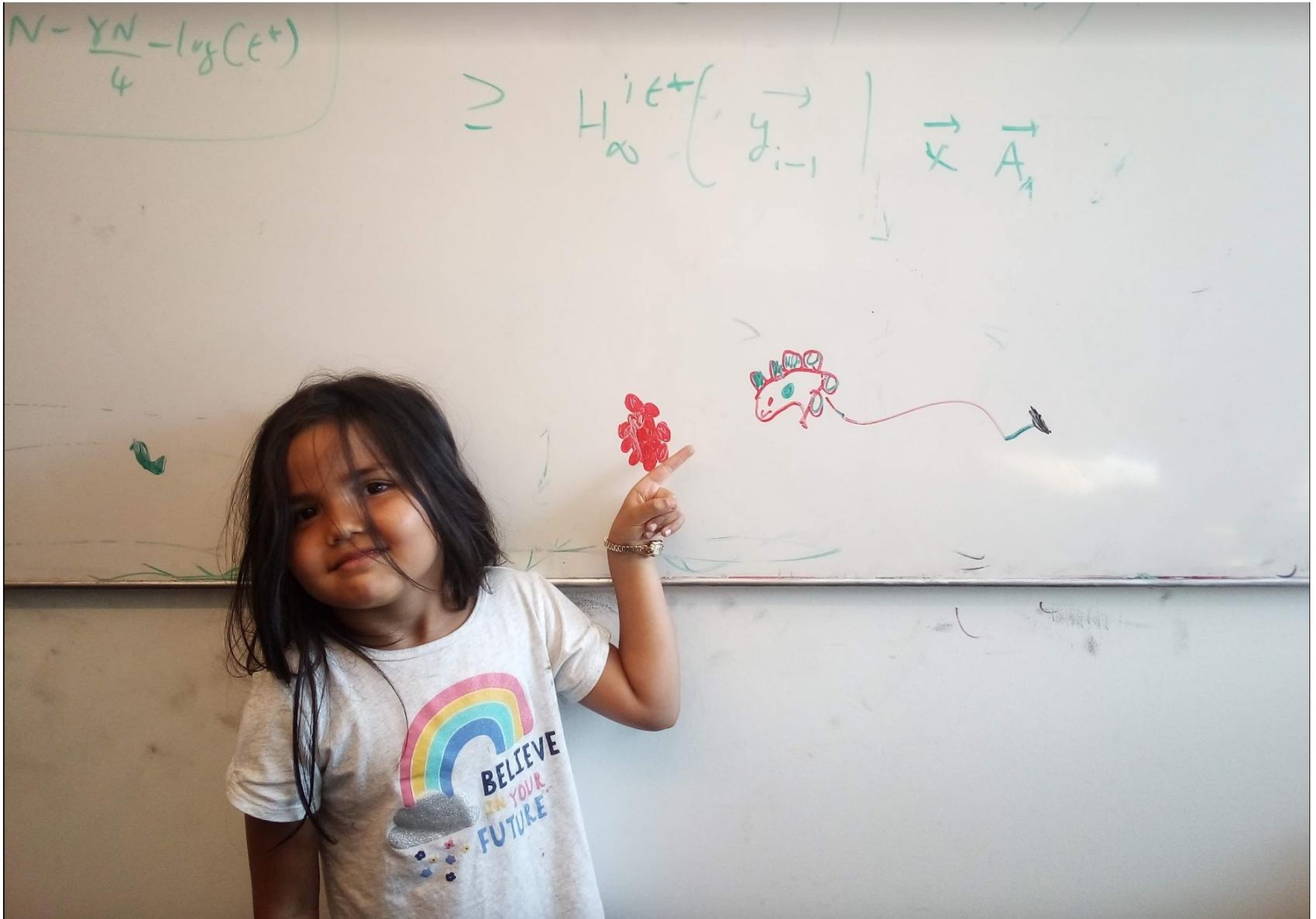
Resolved Unresolved @392_f3

Actions

Atri Rudra 18 hours ago

Since I had to cancel some of the slots today and Wed, I added some slots for next Monday, Nov 6. **However, you need to sign up for those slots by 5pm this Friday.**

Questions/Comments?



Divide and Conquer

Divide up the problem into at least two sub-problems

Recursively solve the sub-problems

“Patch up” the solutions to the sub-problems for the final solution

Improvements on a smaller scale

Greedy algorithms: exponential \rightarrow poly time

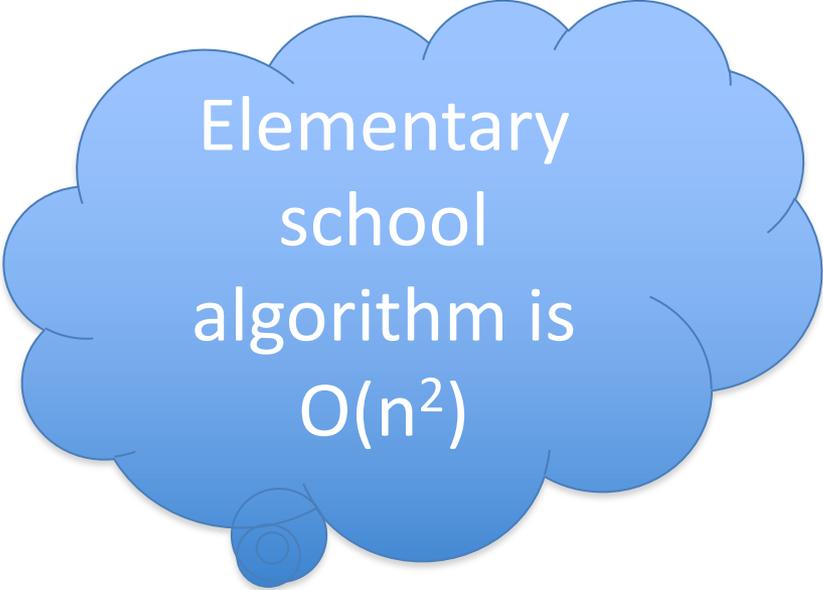
(Typical) Divide and Conquer: $O(n^2)$ \rightarrow asymptotically smaller running time

Multiplying two numbers

Given two numbers a and b in binary

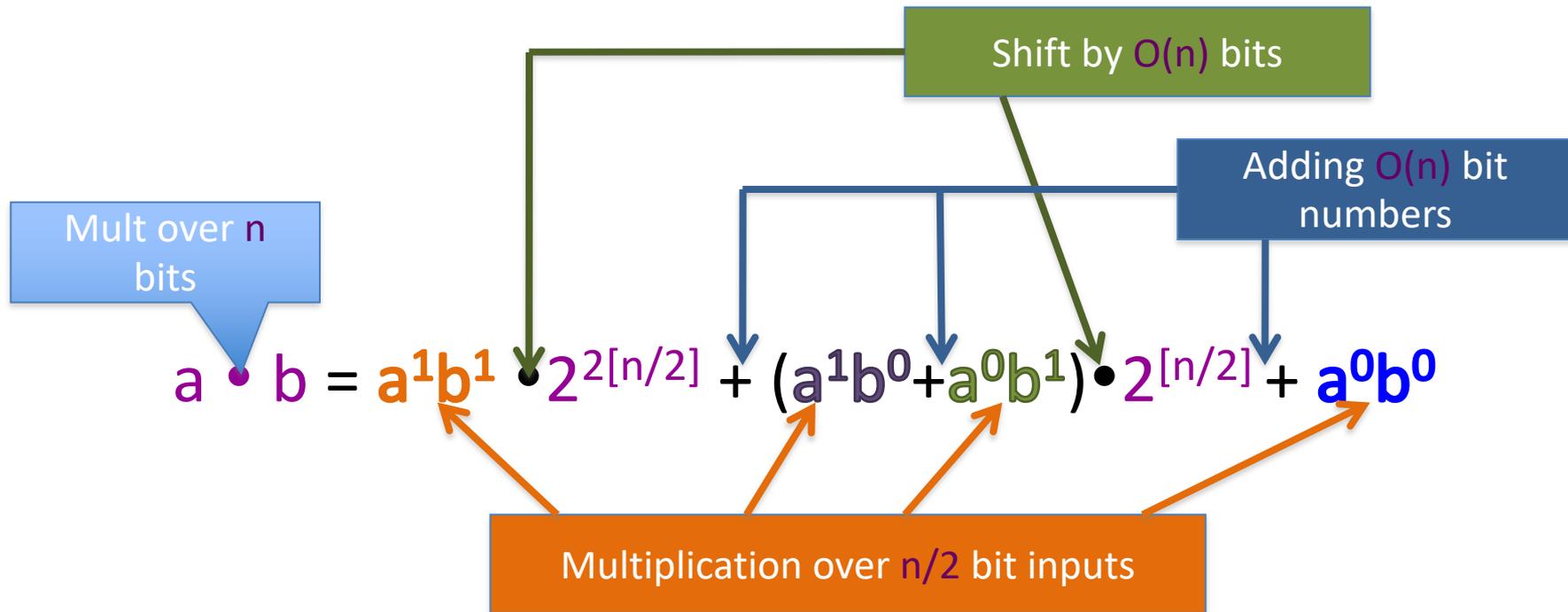
$$a = (a_{n-1}, \dots, a_0) \text{ and } b = (b_{n-1}, \dots, b_0)$$

Compute $c = a \times b$



Elementary
school
algorithm is
 $O(n^2)$

The current algorithm scheme



$$T(n) \leq 4T(n/2) + cn \dots$$

$$T(1) \leq c$$

$T(n)$ is $O(n^2)$

The key identity

$$a^1b^0 + a^0b^1 = (a^1 + a^0)(b^1 + b^0) - a^1b^1 - a^0b^0$$

Wait, how do you think of that?

De-Mystifying the Integer Multiplication Algorithm

In class, we saw an $O(n^{\log_2 3})$ time algorithm to multiply two n bit numbers that used an identity that seemed to be plucked out of thin air. In this note, we will try and de-mystify how one might come about thinking of this identity in the first place.

The setup

We first recall the problem that we are trying to solve:

Multiplying Integers

Given two n bit numbers $a = (a_{n-1}, \dots, a_0)$ and $b = (b_{n-1}, \dots, b_0)$, output their product $c = a \times b$.

Next, recall the following notation that we used:

$$a^0 = (a_{\lceil \frac{n}{2} \rceil - 1}, \dots, a_0),$$

$$a^1 = (a_{n-1}, \dots, a_{\lceil \frac{n}{2} \rceil}),$$

The final algorithm

Input: $a = (a_{n-1}, \dots, a_0)$ and $b = (b_{n-1}, \dots, b_0)$

Mult (a, b)

If $n = 1$ return a_0b_0

$a^1 = a_{n-1}, \dots, a_{\lfloor n/2 \rfloor}$ and $a^0 = a_{\lfloor n/2 \rfloor - 1}, \dots, a_0$

Compute b^1 and b^0 from b

$x = a^1 + a^0$ and $y = b^1 + b^0$

Let $p = \text{Mult}(x, y)$, $D = \text{Mult}(a^1, b^1)$, $E = \text{Mult}(a^0, b^0)$

$F = p - D - E$

return $D \cdot 2^{2\lfloor n/2 \rfloor} + F \cdot 2^{\lfloor n/2 \rfloor} + E$

$$T(1) \leq c$$

$$T(n) \leq 3T(n/2) + cn$$

$O(n^{\log_2 3}) = O(n^{1.59})$
run time

All **green** operations
are $O(n)$ time

$$a \cdot b = a^1 b^1 \cdot 2^{2\lfloor n/2 \rfloor} + ((a^1 + a^0)(b^1 + b^0) - a^1 b^1 - a^0 b^0) \cdot 2^{\lfloor n/2 \rfloor} + a^0 b^0$$