

Sep 19

Implementing GS

Overall runtime

Initialization $\leftarrow \leq T_0$

While (...) $\leftarrow \#iterations = T_1 \leq n^2$

Body $\} \leq T_2$

Output S $\leftarrow \leq T_3$

If we show:
 $T_0, T_3 \leq O(n^2)$
 $T_2 \leq O(1)$

$$\begin{aligned} &\leq T_0 + T_1 \cdot T_2 + T_3 \\ &\leq T_0 + n^2 \cdot T_2 + T_3 \\ &\leq O(n^2) + n^2 \cdot O(1) + O(n) \\ &\leq O(n^2) + O(n^2) + O(n) \\ &\leq O(n^2) \end{aligned}$$

Goal: $T_0, T_3 \leq O(n^2), T_2 \leq O(1)$ & choose the data structures

Notation change

$W = \{w_1, \dots, w_n\} \mapsto \{1, \dots, n\}$

$M = \{m_1, \dots, m_n\}$

$M = \{m_1, \dots, m_n\} \mapsto \{1, \dots, n\}$

def. $[n]$

Our solution to Q1(a) on HWO

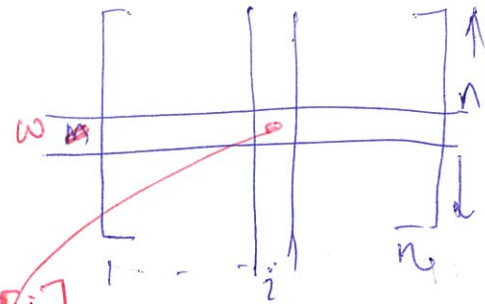
$W = [n] \rightarrow$ Array indices start 1

(Q0) How is the input represented?

2D arrays: $WomanPref, ManPref$
 $n \times n$

$ManPref[M][i] =$ ID of the i th most preferred woman for m in L_m

$WomanPref[W][i] =$ man for w in L_w



Init: n/a

Query: Read $O(1)$

Update: n/a

(Q1) How do we find a free woman?

(AD) Maintain a linked list called free (of free women)

Init: Add all n women to free $O(n)$

Query: Pick (say) the 1st element in free w (+ delete w from free) $O(1)$

Update: w proposes to m

Case 1: m is free \Rightarrow do nothing \checkmark

Case 2.1 (m, w') remain engaged \Rightarrow Add w to free } $O(1)$

Case 2.2 (m, w) get engaged \Rightarrow Add w' to free } $O(1)$

(Q2) How would w pick her best unproposed man?

(A2) Maintain an array of size n called next

$\text{Next}[w] = \text{rank}$ of the man m that w should propose to next

(Q): What is the id of the next man w should propose to?

Read Woman Pref $[w][\text{next}[w]]$ } $O(1)$ time

Init $\text{next}[w] = 1 \quad \forall w \in [n]$ } $O(n)$ time

Query: $O(1)$ time

Update: $\text{next}[w] = \text{next}[w] + 1$ } $O(1)$ time.

(Q3) How do we know who m is currently engaged to?

(A3) Array of size n call current

$\text{current}[m] = \begin{cases} -1 & \text{if } m \text{ is free} \\ w & \text{if } (m, w) \text{ are engaged} \end{cases}$

Init: $\text{current}[m] = -1 \quad \forall m$ } $O(n)$

Query: Read $\text{current}[m]$ } $O(1)$

Update: If after w 's proposal, (m, w) get engaged } $O(1)$
 $\rightarrow \text{current}[m] = w$

Init (4) + (3) $O(n) + O(n) + O(n) = O(n) \leq O(n^2)$

Query/Update $O(1) + O(1) + O(1) \leq O(1)$.