Lecture 12

CSE 331 Sep 23, 2024

Quiz 1 next Monday

Mon, Sep 23	Breadth First Search $\square^{F23} \square^{F22} \square^{F21} \square^{F19} \square^{F18} \square^{F17} x^2$	[KT, Sec 3.2]
Tue, Sep 24		(HW 3 out, HW 2 in)
Wed, Sep 25	Explore algorithm $\square^{F23} \square^{F22} \square^{F21} \square^{F19} \square^{F118} \square^{F17} x^2$	[KT, Sec 3.3]
Fri, Sep 27	Runtime Analysis of BFS algorithm D ^{F23} D ^{F22} D ^{F21} D ^{F19} D ^{F18} D ^{F17} x ²	[KT, Sec 3.3, 3.6] <i>Reading Assignment:</i> [KT, Sec 3.3, 3.4, 3.5, 3.6] <i>Reading Assignment:</i> Care package on topological ordering
Mon, Sep 30	More graph stuff DF23 DF22 DF21 DF19 DF18 DF17 x ²	[KT, Sec 3.3, 3.6] (Quiz 1) (Group Registration on Autolab due)
Tue, Oct 1		(HW 3 in)
Tue, Oct 1 Wed, Oct 2	Interval Scheduling Problem ^{F23} ^{F22} ^{F22} ^{F19} ^{F19} ^{F18} ^{F17} ^{x²}	(HW 3 in) [KT, Sec 4.1] (Project out) Reading Assignment: [KT, Sec 4.1, 4.2]
Tue, Oct 1 Wed, Oct 2 Fri, Oct 4	Interval Scheduling Problem $F^{23} P^{22} P^{21} P^{19} P^{18} P^{17} x^2$ Greedy Algorithm for Interval Scheduling $P^{F23} P^{F22} P^{F19} P^{18} P^{17} x^2$	(HW 3 in) [KT, Sec 4.1] (Project out) Reading Assignment: [KT, Sec 4.1, 4.2] [KT, Sec 4.4] Reading Assignment: Care package on minimizing maximum lateness
Tue, Oct 1 Wed, Oct 2 Fri, Oct 4 Mon, Oct 7	Interval Scheduling Problem $\[Delta F^{23}\] \[Delta F^{22}\] \[Delta F^{19}\] \[Delta F^{18}\] \[Delta F^{17}\] x^2$ Greedy Algorithm for Interval Scheduling $\[Delta F^{23}\] \[Delta F^{22}\] \[Delta F^{19}\] \[Delta F^{18}\] \[Delta F^{17}\] x^2$ Mid-term exam: I	<pre>(HW 3 in) [KT, Sec 4.1] (Project out) Reading Assignment: [KT, Sec 4.1, 4.2] [KT, Sec 4.4] Reading Assignment: Care package on minimizing maximum lateness</pre>

Autolab Project Group Registration

Also due next Monday

note @110 💿 ★ 🔓 -	stop following	27 views		
Register your project groups on Autolab		Actions -		
Now that all the project groups have been assigned (@109), it is time for the next project deadline your group needs to register on Autolab by 11:30pm on Monday, September 30.				
(I was originally planning to have Autolab start t Make sure to EXACTLY follow the instruction If you miss this deadline then you will get a ZERO on the ENTIRE project				
MISSING THIS DEADLINE ALSO MEANS TH by the deadline of Sep 30. Note that completing	hake sure you get vell in advance of	: this done [;] the		
A request: This is the first time we are doing project group registration on Autolab so I would appreciate it is some of you could test out the instructions and confirm that you were successfully able to follow the instruction on Autolab. Thanks!				



Read the instruction carefully



Fill in the Group Composition form FIRST

Make sure you fill in this Google form C^A to submit your group composition. Please see the project overview page for more details on this.

You HAVE to submit the Google form

You have to submit this Google form C by 11:30pm on Friday, September 20. If you do not fill in the form on time, then you will not be able to register your group on Autolab and will get a ZERO (0) on the ENTIRE project.

Register your group on Autolab

Groups on Autolab will NOT be automatically registered

Even after filling the Google form for your group composition you will have to register your group on Autolab by yourself (as a group). Read on for instructions on how to go about this.

However, once you register your group on Autolab you will not have to form your group for the coding and reflections submissions.

Connectivity

u and w are connected iff there is a path between them

A graph is connected iff all pairs of vertices are connected

Questions/Comments?



Rest of Today's agenda

Algorithms for checking connectivity

Checking by inspection



What about large graphs?



Are s and t connected?

Brute-force algorithm?

List all possible vertex sequences between s and t $n^n \text{ such sequences}$ Check if any is a path between s and t

Algorithm motivation



Questions/Comments?



Breadth First Search (BFS)

BFS via examples

In which we derive the breadth first search (BFS) algorithm via a sequence of examples.

Expected background

These notes assume that you are familiar with the following:

- · Graphs and their representation. In particular,
 - · Notion of connectivity of nodes and connected components of graphs
 - Adjacency list representation of graphs
 - Notation:
 - G = (V, E)
 - n = |V| and m = |E|
 - CC(s) denotes the connected component of s
- · Trees and their basic properties

The problem

In these notes we will solve the following problem:

Connectivity Problem

Input: Graph G = (V,E) and s in V

Output: All t connected to s in G

Connected component of s

Breadth First Search (BFS)

Build layers of vertices connected to s

 $L_0 = \{s\}$

Assume $L_0,..,L_j$ have been constructed

 L_{i+1} set of vertices not chosen yet but are connected by an edge to L_i

Stop when new layer is empty

BFS Tree

BFS naturally defines a tree rooted at s

L_i forms the jth "level" in the tree

u in L_{j+1} is child of v in L_j from which it was "discovered"



Argue on the board...



Two facts about BFS trees

(1) All non-tree edges are in the same or consecutive layer

(2) If u is in L_i then dist(s,u) = i

Rest of today's agenda

Computing Connected component

Computing Connected Component



Explore(s)

Start with $R = \{s\}$

While exists (u,w) edge w not in R and u in R

Add w to R

Output $R^* = R$

BFS (Build layers of vertices)

 $\mathbf{L}_0 = \{\mathbf{s}\}$

Assume $L_0,..,L_i$ have been constructed

 L_{j+1} set of vertices not chosen yet but are connected to L_{j} Stop when new layer is empty

Argue correctness on the board...

