

Lecture 20

CSE 331

Oct 13, 2017

Mid-term-I Monday

In class

1:00pm-1:50pm sharp

Eight True/False with justification Qs

Graded Quiz 1

On Autolab by tonight

Graded HW3

Hopefully by tonight

Scheduling to minimize lateness

n jobs: i th job (t_i, d_i)

start time: s

Schedule the n jobs: i th job gets interval $[s(i), f(i)=s(i)+t_i)$

At most one job at any time

Not the sum

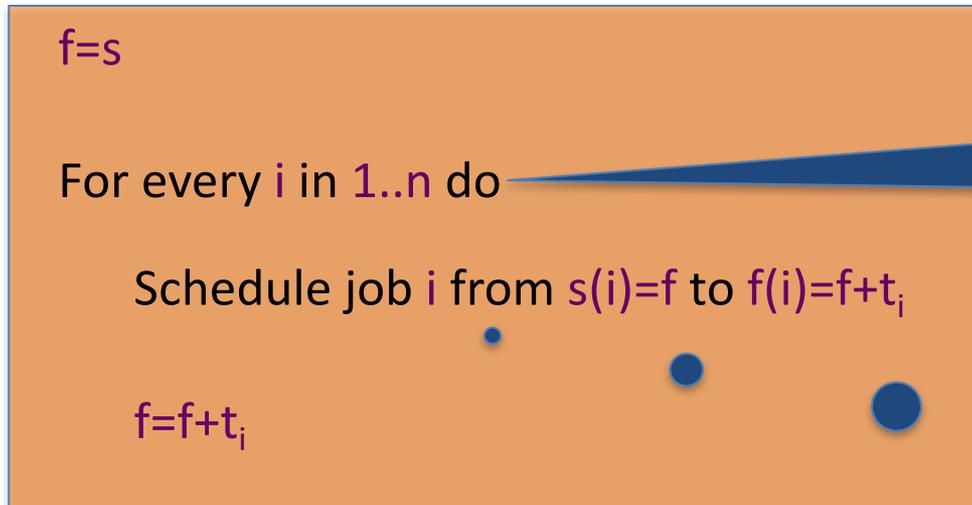
Algo picks $s(i)$

GOAL: Minimize MAXIMUM lateness

Lateness of job i , $l_i = \max(0, f(i) - d_i)$

The Greedy Algorithm

(Assume jobs sorted by deadline: $d_1 \leq d_2 \leq \dots \leq d_n$)

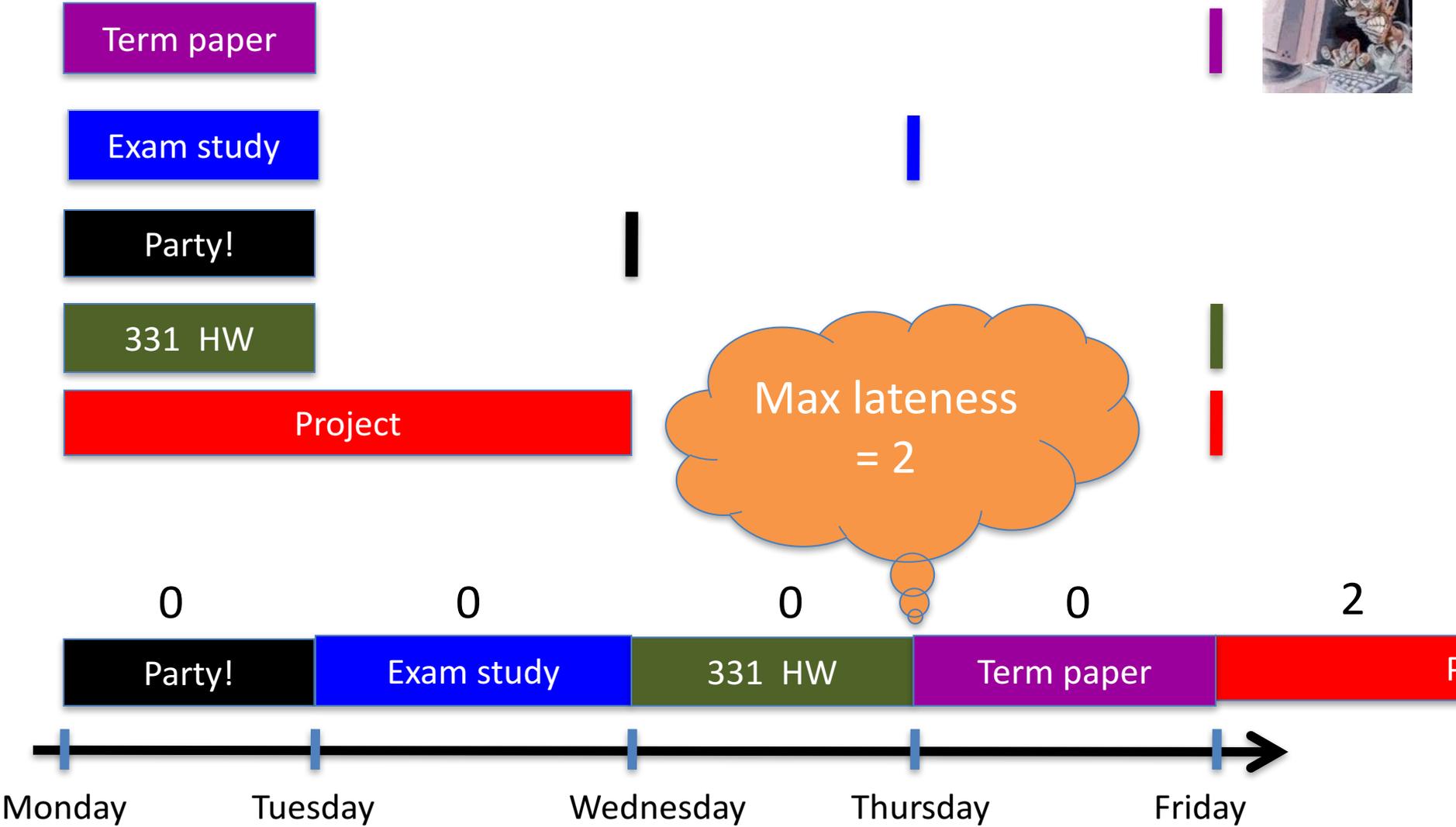


n iterations

$O(1)$ per iteration

$O(n)$ overall.
 $O(n \log n)$ if deadlines are not sorted.

Solving end of Semester blues



Two definitions for schedules

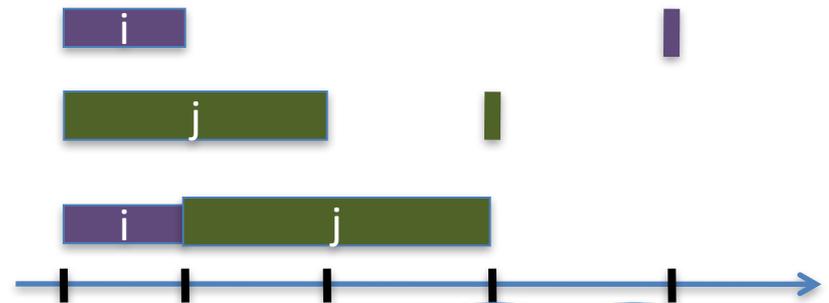
Idle time

Max “gap” between two consecutively scheduled tasks



Inversion

(i,j) is an inversion if i is scheduled before j but $d_i > d_j$



$f=1$

For every i in $1..n$ do

Schedule job i from $s_i=f$ to $f_i=f+t_i$

$f=f+t_i$

0 idle time and 0
inversions for greedy
schedule

We will prove

Any two schedules with 0 idle time and 0 inversions have the same max lateness

Proving greedy is optimal

Any two schedules with 0 idle time and 0 inversions have the same max lateness

Greedy schedule has 0 idle time and 0 inversions

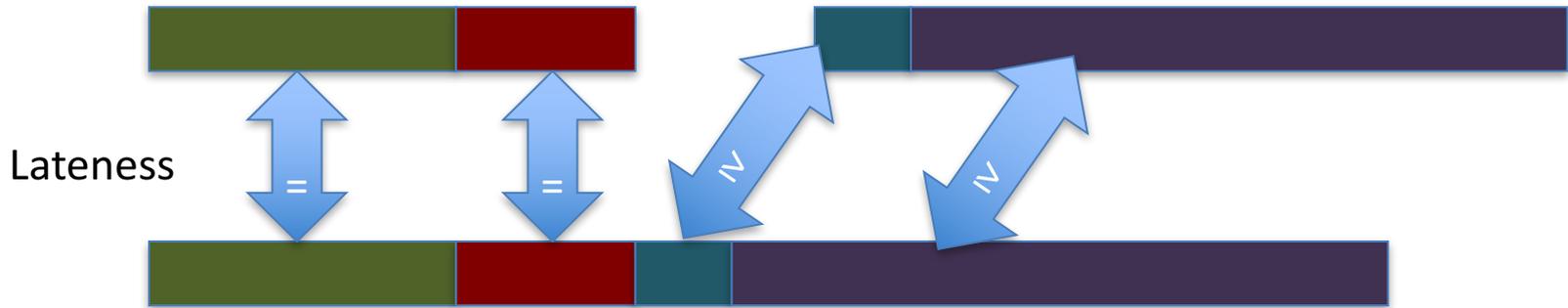
To prove

Any two schedules with 0 idle time and 0 inversions have the same max lateness

Greedy schedule has 0 idle time and 0 inversions

There is an optimal schedule with 0 idle time and 0 inversions

Optimal schedule with 0 idle time



“Only” need to convert a 0
idle optimal ordering to one
with 0 inversions (and 0 idle
time)

Today's agenda

Prove any schedules with 0 idle time and 0 inversions have the same L

“Exchange” argument to convert an optimal solution into a 0 inversion one