

# CSE 191

## Introduction to Discrete Structures

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**Logical Equivalence**

# Outline

## Logical Equivalence

- **Tautologies and Contradictions**
- Logical Equivalence
- Equivalence Laws
- Proving Logical Equivalence

# Tautologies and Contradictions

A **tautology** is a compound proposition that is ALWAYS TRUE, no matter what the truth values of the propositional variables that occur in it.

A **contradiction** is a compound proposition that is ALWAYS FALSE.

A **contingency** is a compound proposition that is neither a contradiction or a tautology. There at least one assignment of truth values to the atomics that can result in TRUE, and at least one that can result in FALSE.

# Tautologies and Contradictions Examples

$p$	$\neg p$	$p \vee \neg p$	$p \wedge \neg p$
F	T		
T	F		

# Tautologies and Contradictions Examples

$p$	$\neg p$	$p \vee \neg p$	$p \wedge \neg p$
F	T	T	
T	F	T	

$p \vee \neg p$  is a tautology

# Tautologies and Contradictions Examples

$p$	$\neg p$	$p \vee \neg p$	$p \wedge \neg p$
F	T	T	F
T	F	T	F

$p \vee \neg p$  is a tautology

$p \wedge \neg p$  is a contradiction

# Tautologies and Contradictions Examples

Construct the truth table for the proposition  $f_1: p \vee \neg(q \wedge p)$

Is  $f_1$  a tautology, contradiction, or contingency?

$p$	$q$	$p \vee \neg(q \wedge p)$
F	F	
F	T	
T	F	
T	T	

# Tautologies and Contradictions Examples

Construct the truth table for the proposition  $f_1: p \vee \neg(q \wedge p)$

Is  $f_1$  a tautology, contradiction, or contingency?

$p$	$q$	$(q \wedge p)$	$\neg(q \wedge p)$	$p \vee \neg(q \wedge p)$
F	F			
F	T			
T	F			
T	T			



# Tautologies and Contradictions Examples

Construct the truth table for the proposition  $f_1: p \vee \neg(q \wedge p)$

Is  $f_1$  a tautology, contradiction, or contingency?

$p$	$q$	$(q \wedge p)$	$\neg(q \wedge p)$	$p \vee \neg(q \wedge p)$
F	F	F		
F	T	F		
T	F	F		
T	T	T		

# Tautologies and Contradictions Examples

Construct the truth table for the proposition  $f_1: p \vee \neg(q \wedge p)$

Is  $f_1$  a tautology, contradiction, or contingency?

$p$	$q$	$(q \wedge p)$	$\neg(q \wedge p)$	$p \vee \neg(q \wedge p)$
F	F	F	T	
F	T	F	T	
T	F	F	T	
T	T	T	F	

# Tautologies and Contradictions Examples

Construct the truth table for the proposition  $f_1: p \vee \neg(q \wedge p)$

Is  $f_1$  a tautology, contradiction, or contingency?

$p$	$q$	$(q \wedge p)$	$\neg(q \wedge p)$	$p \vee \neg(q \wedge p)$
F	F	F	T	T
F	T	F	T	T
T	F	F	T	T
T	T	T	F	T

# Tautologies and Contradictions Examples

$f_1$  is therefore a tautology

We can also say that  $f_1$  is equivalent to TRUE,  $f_1 \equiv T$

$p$	$q$	$(q \wedge p)$	$\neg(q \wedge p)$	$p \vee \neg(q \wedge p)$
F	F	F	T	T
F	T	F	T	T
T	F	F	T	T
T	T	T	F	T

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# Logical Equivalence

Two propositions,  $p$  and  $q$  are logically equivalent if  $p \Leftrightarrow q$  is a tautology

# Logical Equivalence

Two propositions,  $p$  and  $q$  are logically equivalent if  $p \Leftrightarrow q$  is a tautology

- In other words,  $p$  and  $q$  are logically equivalent if their truth values in their truth table are all the same
- Two compound propositions are logically equivalent if their truth values agree for all combinations of the truth values of their atomics
- We write equivalence as  $p \equiv q$ 
  - $\equiv$  is NOT a logical operator
  - $p \equiv q$  is NOT a compound proposition

# Logical Equivalence Example

Are the propositions  $p \rightarrow q$  and  $q \vee \neg p$  logically equivalent?

$p$	$q$	$p \rightarrow q$	$\neg p$	$q \vee \neg p$	$(p \rightarrow q) \Leftrightarrow (q \vee \neg p)$
F	F				
F	T				
T	F				
T	T				



# Logical Equivalence Example

Are the propositions  $p \rightarrow q$  and  $q \vee \neg p$  logically equivalent?

$p$	$q$	$p \rightarrow q$	$\neg p$	$q \vee \neg p$	$(p \rightarrow q) \Leftrightarrow (q \vee \neg p)$
F	F	T	T	T	T
F	T	T	T	T	T
T	F	F	F	F	T
T	T	T	F	T	T

# Logical Equivalence Example

Are the propositions  $p \rightarrow q$  and  $q \vee \neg p$  logically equivalent?

$p$	$q$	$p \rightarrow q$	$\neg p$	$q \vee \neg p$	$(p \rightarrow q) \Leftrightarrow (q \vee \neg p)$
F	F	T	T		
F	T	T	T		
T	F	F	F		
T	T	T	F		

# Logical Equivalence Example

Are the propositions  $p \rightarrow q$  and  $q \vee \neg p$  logically equivalent?

$p$	$q$	$p \rightarrow q$	$\neg p$	$q \vee \neg p$	$(p \rightarrow q) \Leftrightarrow (q \vee \neg p)$
F	F	T	T	T	
F	T	T	T	T	
T	F	F	F	F	
T	T	T	F	T	

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Are the propositions  $p \rightarrow q$  and  $q \vee \neg p$  logically equivalent?

$p$	$q$	$p \rightarrow q$	$\neg p$	$q \vee \neg p$	$(p \rightarrow q) \Leftrightarrow (q \vee \neg p)$
F	F	T	T	T	T
F	T	T	T	T	T
T	F	F	F	F	T
T	T	T	F	T	T

# Logical Equivalence Example

$p$	$q$	$p \rightarrow q$	$\neg p$	$q \vee \neg p$	$(p \rightarrow q) \Leftrightarrow (q \vee \neg p)$
F	F	T	T	T	T
F	T	T	T	T	T
T	F	F	F	F	T
T	T	T	F	T	T

Columns 3 and 5 are identical, therefore the bidirectional implication of the two is always TRUE. Therefore  $p \rightarrow q \equiv q \vee \neg p$ .

# Logical Equivalence Example

Are the propositions  $p \oplus q$  and  $f_2: (p \vee q) \wedge (p \rightarrow q)$  logically equivalent?

$p$	$q$	$p \oplus q$	$p \vee q$	$p \rightarrow q$	$f_2$	$(p \oplus q) \Leftrightarrow f_2$
F	F					
F	T					
T	F					
T	T					

# Logical Equivalence Example

Are the propositions  $p \oplus q$  and  $f_2: (p \vee q) \wedge (p \rightarrow q)$  logically equivalent?

$p$	$q$	$p \oplus q$	$p \vee q$	$p \rightarrow q$	$f_2$	$(p \oplus q) \Leftrightarrow f_2$
F	F	F	F	T		
F	T	T	T	T		
T	F	T	T	F		
T	T	F	T	T		

# Logical Equivalence Example

Are the propositions  $p \oplus q$  and  $f_2: (p \vee q) \wedge (p \rightarrow q)$  logically equivalent?

$p$	$q$	$p \oplus q$	$p \vee q$	$p \rightarrow q$	$f_2$	$(p \oplus q) \Leftrightarrow f_2$
F	F	F	F	T	F	
F	T	T	T	T	T	
T	F	T	T	F	F	
T	T	F	T	T	T	



# Logical Equivalence Example

Are the propositions  $p \oplus q$  and  $f_2: (p \vee q) \wedge (p \rightarrow q)$  logically equivalent?

$p$	$q$	$p \oplus q$	$p \vee q$	$p \rightarrow q$	$f_2$	$(p \oplus q) \Leftrightarrow f_2$
F	F	F	F	T	F	T
F	T	T	T	T	T	T
T	F	T	T	F	F	F
T	T	F	T	T	T	F

# Logical Equivalence Example

Are the propositions  $p \oplus q$  and  $f_2: (p \vee q) \wedge (p \rightarrow q)$  logically equivalent?

**No, they are not equivalent**

$p$	$q$	$p \oplus q$	$p \vee q$	$p \rightarrow q$	$f_2$	$(p \oplus q) \Leftrightarrow f_2$
F	F	F	F	T	F	T
F	T	T	T	T	T	T
T	F	T	T	F	F	F
T	T	F	T	T	T	F

# Logical Equivalence Examples

Using truth tables we have proven that:

- $(p \rightarrow q) \equiv (q \vee \neg p)$
- $(p \oplus q) \neq f_2$

**Exercise**

- Prove  $(p \oplus q) \equiv \neg(p \leftrightarrow q)$
- Prove  $(p \oplus q) \equiv (p \vee q) \wedge (p \rightarrow \neg q)$

# Outline

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- **Equivalence Laws**
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# De Morgan's Law

$$(1) \quad \neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$(2) \quad \neg(p \vee q) \equiv \neg p \wedge \neg q$$

$p$	$q$	$p \wedge q$	$\neg(p \wedge q)$	$\neg p$	$\neg q$	$\neg p \vee \neg q$
F	F					
F	T					
T	F					
T	T					

Proof for (1). The proof for (2) is similar.

# De Morgan's Law

$$(1) \quad \neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$(2) \quad \neg(p \vee q) \equiv \neg p \wedge \neg q$$

$p$	$q$	$p \wedge q$	$\neg(p \wedge q)$	$\neg p$	$\neg q$	$\neg p \vee \neg q$
F	F	F				
F	T	F				
T	F	F				
T	T	T				

Proof for (1). The proof for (2) is similar.

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$$(1) \quad \neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$(2) \quad \neg(p \vee q) \equiv \neg p \wedge \neg q$$

$p$	$q$	$p \wedge q$	$\neg(p \wedge q)$	$\neg p$	$\neg q$	$\neg p \vee \neg q$
F	F	F	T			
F	T	F	T			
T	F	F	T			
T	T	T	F			

Proof for (1). The proof for (2) is similar.

# De Morgan's Law

$$(1) \quad \neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$(2) \quad \neg(p \vee q) \equiv \neg p \wedge \neg q$$

$p$	$q$	$p \wedge q$	$\neg(p \wedge q)$	$\neg p$	$\neg q$	$\neg p \vee \neg q$
F	F	F	T	T		
F	T	F	T	T		
T	F	F	T	F		
T	T	T	F	F		

Proof for (1). The proof for (2) is similar.



# De Morgan's Law

$$(1) \quad \neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$(2) \quad \neg(p \vee q) \equiv \neg p \wedge \neg q$$

$p$	$q$	$p \wedge q$	$\neg(p \wedge q)$	$\neg p$	$\neg q$	$\neg p \vee \neg q$
F	F	F	T	T	T	
F	T	F	T	T	F	
T	F	F	T	F	T	
T	T	T	F	F	F	

Proof for (1). The proof for (2) is similar.

# De Morgan's Law

$$(1) \quad \neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$(2) \quad \neg(p \vee q) \equiv \neg p \wedge \neg q$$

$p$	$q$	$p \wedge q$	$\neg(p \wedge q)$	$\neg p$	$\neg q$	$\neg p \vee \neg q$
F	F	F	T	T	T	T
F	T	F	T	T	F	T
T	F	F	T	F	T	T
T	T	T	F	F	F	F

Proof for (1). The proof for (2) is similar.

# De Morgan's Law

$$(1) \quad \neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$(2) \quad \neg(p \vee q) \equiv \neg p \wedge \neg q$$

$p$	$q$	$p \wedge q$	$\neg(p \wedge q)$	$\neg p$	$\neg q$	$\neg p \vee \neg q$
F	F	F	T	T	T	T
F	T	F	T	T	F	T
T	F	F	T	F	T	T
T	T	T	F	F	F	F

Proof for (1). The proof for (2) is similar.

# Law of Distributivity

(1)  $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$   
(2)  $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$

$p$	$q$	$r$	$(q \wedge r)$	$p \vee (q \wedge r)$	$(p \vee q)$	$(p \vee r)$	$(p \vee q) \wedge (p \vee r)$
F	F	F	F	F	F	F	F
F	F	T	F	F	F	T	F
F	T	F	F	F	T	F	F
F	T	T	T	T	T	T	T
T	F	F	F	T	T	T	T
T	F	T	F	T	T	T	T
T	T	F	F	T	T	T	T
T	T	T	T	T	T	T	T

# Examples

## Exercise

Prove case (2) of De Morgan's Law the Law of Distributivity using truth tables

### De Morgan's Law

$$(2) \quad \neg(p \vee q) \equiv \neg p \wedge \neg q$$

### Law of Distributivity

$$(2) \quad p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$

# Law of Contraposition

**Law of Contraposition:** An implication is always equivalent to its contrapositive

- **Reminder:** the contrapositive of  $p \rightarrow q$  is  $\neg q \rightarrow \neg p$

$p$	$q$	$p \rightarrow q$	$\neg p$	$\neg q$	$\neg q \rightarrow \neg p$
F	F	T	T	T	T
F	T	T	T	F	T
T	F	F	F	T	F
T	T	T	F	F	T

**Note:** A common proof technique is called proof by contraposition. Prove that the contrapositive is true, therefore the implication itself is true.

# Converse and Inverse

**Converse**: The converse of  $p \rightarrow q$  is  $q \rightarrow p$

**Inverse**: The inverse of  $p \rightarrow q$  is  $\neg p \rightarrow \neg q$

## Exercises

- Prove that the converse of  $p \rightarrow q$  is NOT equivalent to  $p \rightarrow q$
- Prove that the inverse of  $p \rightarrow q$  is NOT equivalent to  $p \rightarrow q$

# Logical Equivalence Rules

Equivalence	Name
$p \wedge T \equiv p$ $p \vee F \equiv p$	Identity laws
$p \vee T \equiv T$ $p \wedge F \equiv F$	Domination laws
$p \vee p \equiv p$ $p \wedge p \equiv p$	Idempotent laws
$\neg(\neg p) \equiv p$	Double negation law
$p \vee q \equiv q \vee p$ $p \wedge q \equiv q \wedge p$	Commutative laws
$(p \vee q) \vee r \equiv p \vee (q \vee r)$ $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$	Associative laws
$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$ $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	Distributive laws
$\neg(p \vee q) \equiv \neg p \wedge \neg q$ $\neg(p \wedge q) \equiv \neg p \vee \neg q$	De Morgan's laws
$p \vee (p \wedge q) \equiv p$ $p \wedge (p \vee q) \equiv p$	Absorption laws
$p \vee \neg p \equiv T$ $p \wedge \neg p \equiv F$	Negation laws



# Logical Equivalence Rules

## Equivalences with Implication

$$p \rightarrow q \equiv \neg p \vee q$$

$$p \rightarrow q \equiv \neg q \rightarrow \neg p$$

$$p \vee q \equiv \neg p \rightarrow q$$

$$p \wedge q \equiv \neg(p \rightarrow \neg q)$$

$$\neg(p \rightarrow q) \equiv p \wedge \neg q$$

$$(p \rightarrow q) \wedge (p \rightarrow r) \equiv p \rightarrow (q \wedge r)$$

$$(p \rightarrow q) \wedge (q \rightarrow r) \equiv (p \vee q) \rightarrow r$$

$$(p \rightarrow q) \vee (p \rightarrow r) \equiv p \rightarrow (q \vee r)$$

$$(p \rightarrow q) \vee (q \rightarrow r) \equiv (p \wedge q) \rightarrow r$$

## Equivalences with Bidirectional Implication

$$p \Leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$$

$$p \Leftrightarrow q \equiv q \Leftrightarrow p$$

$$p \Leftrightarrow q \equiv \neg p \Leftrightarrow \neg q$$

$$p \Leftrightarrow q \equiv (p \wedge q) \vee (\neg p \wedge \neg q)$$

$$\neg(p \Leftrightarrow q) \equiv p \Leftrightarrow \neg q$$

# Proving Logical Equivalence

- By using equivalence laws, we can prove two propositions are logically equivalent without having to construct large truth tables
- The logical equivalences shown in the tables can be used to construct additional logical equivalences

# Proving Logical Equivalence

**Example 1:** Prove  $\neg(p \vee (\neg p \wedge q)) \equiv \neg p \wedge \neg q$

Equivalence	Name
$p \wedge \top \equiv p$ $p \vee \text{F} \equiv p$	Identity laws
$p \vee \top \equiv \top$ $p \wedge \text{F} \equiv \text{F}$	Domination laws
$p \vee p \equiv p$ $p \wedge p \equiv p$	Idempotent laws
$\neg(\neg p) \equiv p$	Double negation law
$p \vee q \equiv q \vee p$ $p \wedge q \equiv q \wedge p$	Commutative laws
$(p \vee q) \vee r \equiv p \vee (q \vee r)$ $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$	Associative laws
$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$ $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	Distributive laws
$\neg(p \vee q) \equiv \neg p \wedge \neg q$ $\neg(p \wedge q) \equiv \neg p \vee \neg q$	De Morgan's laws
$p \vee (p \wedge q) \equiv p$ $p \wedge (p \vee q) \equiv p$	Absorption laws
$p \vee \neg p \equiv \top$ $p \wedge \neg p \equiv \text{F}$	Negation laws

# Proving Logical Equivalence

Example 1: Prove  $\neg(p \vee (\neg p \wedge q)) \equiv \neg p \wedge \neg q$

$\neg(p \vee (\neg p \wedge q)) \equiv \neg p \wedge \neg(\neg p \wedge q)$  by De Morgan's Law

Equivalence	Name
...	
$\neg(p \vee q) \equiv \neg p \wedge \neg q$ $\neg(p \wedge q) \equiv \neg p \vee \neg q$	De Morgan's laws
...	

# Proving Logical Equivalence

**Example 1:** Prove  $\neg(p \vee (\neg p \wedge q)) \equiv \neg p \wedge \neg q$

$$\begin{aligned}\neg(p \vee (\neg p \wedge q)) &\equiv \neg p \wedge \neg(\neg p \wedge q) && \text{by De Morgan's Law} \\ &\equiv \neg p \wedge (\neg\neg p \vee \neg q) && \text{by De Morgan's Law}\end{aligned}$$

Equivalence	Name
...	
$\neg(p \vee q) \equiv \neg p \wedge \neg q$ $\neg(p \wedge q) \equiv \neg p \vee \neg q$	De Morgan's laws
...	

# Proving Logical Equivalence

**Example 1:** Prove  $\neg(p \vee (\neg p \wedge q)) \equiv \neg p \wedge \neg q$

$$\begin{aligned}\neg(p \vee (\neg p \wedge q)) &\equiv \neg p \wedge \neg(\neg p \wedge q) && \text{by De Morgan's Law} \\ &\equiv \neg p \wedge (\neg\neg p \vee \neg q) && \text{by De Morgan's Law} \\ &\equiv \neg p \wedge (p \vee \neg q) && \text{by Double Negation Law}\end{aligned}$$

Equivalence	Name
...	
$\neg(\neg p) \equiv p$	Double negation law
...	

# Proving Logical Equivalence

**Example 1:** Prove  $\neg(p \vee (\neg p \wedge q)) \equiv \neg p \wedge \neg q$

$$\begin{aligned}\neg(p \vee (\neg p \wedge q)) &\equiv \neg p \wedge \neg(\neg p \wedge q) && \text{by De Morgan's Law} \\ &\equiv \neg p \wedge (\neg\neg p \vee \neg q) && \text{by De Morgan's Law} \\ &\equiv \neg p \wedge (p \vee \neg q) && \text{by Double Negation Law} \\ &\equiv (\neg p \wedge p) \vee (\neg p \wedge \neg q) && \text{by Distributive Law}\end{aligned}$$

Equivalence	Name
...	
$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$ $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	Distributive laws

# Proving Logical Equivalence

**Example 1:** Prove  $\neg(p \vee (\neg p \wedge q)) \equiv \neg p \wedge \neg q$

$$\begin{aligned}\neg(p \vee (\neg p \wedge q)) &\equiv \neg p \wedge \neg(\neg p \wedge q) && \text{by De Morgan's Law} \\ &\equiv \neg p \wedge (\neg\neg p \vee \neg q) && \text{by De Morgan's Law} \\ &\equiv \neg p \wedge (p \vee \neg q) && \text{by Double Negation Law} \\ &\equiv (\neg p \wedge p) \vee (\neg p \wedge \neg q) && \text{by Distributive Law} \\ &\equiv (p \wedge \neg p) \vee (\neg p \wedge \neg q) && \text{by Commutative Law}\end{aligned}$$

Equivalence	Name
...	
$p \vee q \equiv q \vee p$ $p \wedge q \equiv q \wedge p$	Commutative laws
...	



# Proving Logical Equivalence

**Example 1:** Prove  $\neg(p \vee (\neg p \wedge q)) \equiv \neg p \wedge \neg q$

$$\begin{aligned}\neg(p \vee (\neg p \wedge q)) &\equiv \neg p \wedge \neg(\neg p \wedge q) && \text{by De Morgan's Law} \\ &\equiv \neg p \wedge (\neg\neg p \vee \neg q) && \text{by De Morgan's Law} \\ &\equiv \neg p \wedge (p \vee \neg q) && \text{by Double Negation Law} \\ &\equiv (\neg p \wedge p) \vee (\neg p \wedge \neg q) && \text{by Distributive Law} \\ &\equiv (p \wedge \neg p) \vee (\neg p \wedge \neg q) && \text{by Commutative Law} \\ &\equiv F \vee (\neg p \wedge \neg q) && \text{by Negation Law}\end{aligned}$$

Equivalence	Name
...	
$p \vee \neg p \equiv \top$ $p \wedge \neg p \equiv F$	Negation laws
...	

# Proving Logical Equivalence

Example 1: Prove  $\neg(p \vee (\neg p \wedge q)) \equiv \neg p \wedge \neg q$

$$\begin{aligned}\neg(p \vee (\neg p \wedge q)) &\equiv \neg p \wedge \neg(\neg p \wedge q) && \text{by De Morgan's Law} \\ &\equiv \neg p \wedge (\neg\neg p \vee \neg q) && \text{by De Morgan's Law} \\ &\equiv \neg p \wedge (p \vee \neg q) && \text{by Double Negation Law} \\ &\equiv (\neg p \wedge p) \vee (\neg p \wedge \neg q) && \text{by Distributive Law} \\ &\equiv (p \wedge \neg p) \vee (\neg p \wedge \neg q) && \text{by Commutative Law} \\ &\equiv F \vee (\neg p \wedge \neg q) && \text{by Negation Law} \\ &\equiv (\neg p \wedge \neg q) \vee F && \text{by Commutative Law}\end{aligned}$$

# Proving Logical Equivalence

Example 1: Prove  $\neg(p \vee (\neg p \wedge q)) \equiv \neg p \wedge \neg q$

$$\begin{aligned}\neg(p \vee (\neg p \wedge q)) &\equiv \neg p \wedge \neg(\neg p \wedge q) && \text{by De Morgan's Law} \\ &\equiv \neg p \wedge (\neg\neg p \vee \neg q) && \text{by De Morgan's Law} \\ &\equiv \neg p \wedge (p \vee \neg q) && \text{by Double Negation Law} \\ &\equiv (\neg p \wedge p) \vee (\neg p \wedge \neg q) && \text{by Distributive Law} \\ &\equiv (p \wedge \neg p) \vee (\neg p \wedge \neg q) && \text{by Commutative Law} \\ &\equiv F \vee (\neg p \wedge \neg q) && \text{by Negation Law} \\ &\equiv (\neg p \wedge \neg q) \vee F && \text{by Commutative Law} \\ &\equiv (\neg p \wedge \neg q) && \text{by Identity Law}\end{aligned}$$

# Proving Logical Equivalence

## In General:

- Each line should be equivalent to the previous
- Each line should list the law that led to it
  - Exactly one law applied per line
- Start with LHS and go until you reach the RHS

**Note:** Logical equivalence proofs are very exact. Later proofs will be less restrictive.

# Proving Logical Equivalence

**Exercise:** Prove  $(p \wedge q) \rightarrow (p \vee q) \equiv \top$

Equivalence	Name
$p \wedge \top \equiv p$ $p \vee \text{F} \equiv p$	Identity laws
$p \vee \top \equiv \top$ $p \wedge \text{F} \equiv \text{F}$	Domination laws
$p \vee p \equiv p$ $p \wedge p \equiv p$	Idempotent laws
$\neg(\neg p) \equiv p$	Double negation law
$p \vee q \equiv q \vee p$ $p \wedge q \equiv q \wedge p$	Commutative laws
$(p \vee q) \vee r \equiv p \vee (q \vee r)$ $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$	Associative laws
$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$ $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	Distributive laws
$\neg(p \vee q) \equiv \neg p \wedge \neg q$ $\neg(p \wedge q) \equiv \neg p \vee \neg q$	De Morgan's laws
$p \vee (p \wedge q) \equiv p$ $p \wedge (p \vee q) \equiv p$	Absorption laws
$p \vee \neg p \equiv \top$ $p \wedge \neg p \equiv \text{F}$	Negation laws

Equivalences with Implication
$p \rightarrow q \equiv \neg p \vee q$
$p \rightarrow q \equiv \neg q \rightarrow \neg p$
$p \vee q \equiv \neg p \rightarrow q$
$p \wedge q \equiv \neg(p \rightarrow \neg q)$
$\neg(p \rightarrow q) \equiv p \wedge \neg q$
$(p \rightarrow q) \wedge (p \rightarrow r) \equiv p \rightarrow (q \wedge r)$
$(p \rightarrow q) \wedge (q \rightarrow r) \equiv (p \vee q) \rightarrow r$
$(p \rightarrow q) \vee (p \rightarrow r) \equiv p \rightarrow (q \vee r)$
$(p \rightarrow q) \vee (q \rightarrow r) \equiv (p \wedge q) \rightarrow r$

# Proving Logical Equivalence

Exercise: Prove  $(p \wedge q) \rightarrow (p \vee q) \equiv \top$

$(p \wedge q) \rightarrow (p \vee q) \equiv \neg(p \wedge q) \vee (p \vee q)$       by Conditional Law

# Proving Logical Equivalence

Exercise: Prove  $(p \wedge q) \rightarrow (p \vee q) \equiv \top$

$$\begin{aligned}(p \wedge q) \rightarrow (p \vee q) &\equiv \neg(p \wedge q) \vee (p \vee q) && \text{by Conditional Law} \\ &\equiv (\neg p \vee \neg q) \vee (p \vee q) && \text{by De Morgan's Law}\end{aligned}$$

# Proving Logical Equivalence

Exercise: Prove  $(p \wedge q) \rightarrow (p \vee q) \equiv \top$

$$\begin{aligned}(p \wedge q) \rightarrow (p \vee q) &\equiv \neg(p \wedge q) \vee (p \vee q) && \text{by Conditional Law} \\ &\equiv (\neg p \vee \neg q) \vee (p \vee q) && \text{by De Morgan's Law} \\ &\equiv \neg p \vee (\neg q \vee (p \vee q)) && \text{by Associative Law}\end{aligned}$$



# Proving Logical Equivalence

Exercise: Prove  $(p \wedge q) \rightarrow (p \vee q) \equiv \top$

$$\begin{aligned}(p \wedge q) \rightarrow (p \vee q) &\equiv \neg(p \wedge q) \vee (p \vee q) && \text{by Conditional Law} \\ &\equiv (\neg p \vee \neg q) \vee (p \vee q) && \text{by De Morgan's Law} \\ &\equiv \neg p \vee (\neg q \vee (p \vee q)) && \text{by Associative Law} \\ &\equiv \neg p \vee ((\neg q \vee p) \vee q) && \text{by Associative Law}\end{aligned}$$

# Proving Logical Equivalence

Exercise: Prove  $(p \wedge q) \rightarrow (p \vee q) \equiv \top$

$$\begin{aligned}(p \wedge q) \rightarrow (p \vee q) &\equiv \neg(p \wedge q) \vee (p \vee q) && \text{by Conditional Law} \\ &\equiv (\neg p \vee \neg q) \vee (p \vee q) && \text{by De Morgan's Law} \\ &\equiv \neg p \vee (\neg q \vee (p \vee q)) && \text{by Associative Law} \\ &\equiv \neg p \vee ((\neg q \vee p) \vee q) && \text{by Associative Law} \\ &\equiv \neg p \vee ((p \vee \neg q) \vee q) && \text{by Commutative Law}\end{aligned}$$

# Proving Logical Equivalence

Exercise: Prove  $(p \wedge q) \rightarrow (p \vee q) \equiv \top$

$$\begin{aligned}(p \wedge q) \rightarrow (p \vee q) &\equiv \neg(p \wedge q) \vee (p \vee q) && \text{by Conditional Law} \\ &\equiv (\neg p \vee \neg q) \vee (p \vee q) && \text{by De Morgan's Law} \\ &\equiv \neg p \vee (\neg q \vee (p \vee q)) && \text{by Associative Law} \\ &\equiv \neg p \vee ((\neg q \vee p) \vee q) && \text{by Associative Law} \\ &\equiv \neg p \vee ((p \vee \neg q) \vee q) && \text{by Commutative Law} \\ &\equiv \neg p \vee (p \vee (\neg q \vee q)) && \text{by Associative Law}\end{aligned}$$

# Proving Logical Equivalence

Exercise: Prove  $(p \wedge q) \rightarrow (p \vee q) \equiv \top$

$$\begin{aligned}(p \wedge q) \rightarrow (p \vee q) &\equiv \neg(p \wedge q) \vee (p \vee q) && \text{by Conditional Law} \\ &\equiv (\neg p \vee \neg q) \vee (p \vee q) && \text{by De Morgan's Law} \\ &\equiv \neg p \vee (\neg q \vee (p \vee q)) && \text{by Associative Law} \\ &\equiv \neg p \vee ((\neg q \vee p) \vee q) && \text{by Associative Law} \\ &\equiv \neg p \vee ((p \vee \neg q) \vee q) && \text{by Commutative Law} \\ &\equiv \neg p \vee (p \vee (\neg q \vee q)) && \text{by Associative Law} \\ &\equiv (\neg p \vee p) \vee (\neg q \vee q) && \text{by Associative Law}\end{aligned}$$

# Proving Logical Equivalence

**Exercise:** Prove  $(p \wedge q) \rightarrow (p \vee q) \equiv \top$

$$(p \wedge q) \rightarrow (p \vee q) \equiv \neg(p \wedge q) \vee (p \vee q) \quad \text{by Conditional Law}$$

...

$$\equiv (\neg p \vee p) \vee (\neg q \vee q) \quad \text{by Associative Law}$$

# Proving Logical Equivalence

**Exercise:** Prove  $(p \wedge q) \rightarrow (p \vee q) \equiv \top$

$$(p \wedge q) \rightarrow (p \vee q) \equiv \neg(p \wedge q) \vee (p \vee q) \quad \text{by Conditional Law}$$

...

...

$$\equiv (\neg p \vee p) \vee (\neg q \vee q) \quad \text{by Associative Law}$$

$$\equiv (p \vee \neg p) \vee (\neg q \vee q) \quad \text{by Commutative Law}$$

# Proving Logical Equivalence

Exercise: Prove  $(p \wedge q) \rightarrow (p \vee q) \equiv \top$

$$(p \wedge q) \rightarrow (p \vee q) \equiv \neg(p \wedge q) \vee (p \vee q) \quad \text{by Conditional Law}$$

...

...

$$\equiv (\neg p \vee p) \vee (\neg q \vee q) \quad \text{by Associative Law}$$

$$\equiv (p \vee \neg p) \vee (\neg q \vee q) \quad \text{by Commutative Law}$$

$$\equiv (p \vee \neg p) \vee (q \vee \neg q) \quad \text{by Commutative Law}$$

# Proving Logical Equivalence

**Exercise:** Prove  $(p \wedge q) \rightarrow (p \vee q) \equiv \top$

$$(p \wedge q) \rightarrow (p \vee q) \equiv \neg(p \wedge q) \vee (p \vee q) \quad \text{by Conditional Law}$$

...

...

$$\equiv (\neg p \vee p) \vee (\neg q \vee q) \quad \text{by Associative Law}$$

$$\equiv (p \vee \neg p) \vee (\neg q \vee q) \quad \text{by Commutative Law}$$

$$\equiv (p \vee \neg p) \vee (q \vee \neg q) \quad \text{by Commutative Law}$$

$$\equiv \top \vee (q \vee \neg q) \quad \text{by Negation Law}$$



# Proving Logical Equivalence

**Exercise:** Prove  $(p \wedge q) \rightarrow (p \vee q) \equiv \top$

$$\begin{aligned}(p \wedge q) \rightarrow (p \vee q) &\equiv \neg(p \wedge q) \vee (p \vee q) && \text{by Conditional Law} \\ &\dots && \dots \\ &\equiv (\neg p \vee p) \vee (\neg q \vee q) && \text{by Associative Law} \\ &\equiv (p \vee \neg p) \vee (\neg q \vee q) && \text{by Commutative Law} \\ &\equiv (p \vee \neg p) \vee (q \vee \neg q) && \text{by Commutative Law} \\ &\equiv \top \vee (q \vee \neg q) && \text{by Negation Law} \\ &\equiv \top \vee \top && \text{by Negation Law}\end{aligned}$$

# Proving Logical Equivalence

**Exercise:** Prove  $(p \wedge q) \rightarrow (p \vee q) \equiv \top$

$$\begin{aligned}(p \wedge q) \rightarrow (p \vee q) &\equiv \neg(p \wedge q) \vee (p \vee q) && \text{by Conditional Law} \\ &\dots && \dots \\ &\equiv (\neg p \vee p) \vee (\neg q \vee q) && \text{by Associative Law} \\ &\equiv (p \vee \neg p) \vee (\neg q \vee q) && \text{by Commutative Law} \\ &\equiv (p \vee \neg p) \vee (q \vee \neg q) && \text{by Commutative Law} \\ &\equiv \top \vee (q \vee \neg q) && \text{by Negation Law} \\ &\equiv \top \vee \top && \text{by Negation Law} \\ &\equiv \top && \text{by Domination Law}\end{aligned}$$

# Proving Logical Equivalence

**Exercise:** Prove  $(p \rightarrow q) \wedge (p \rightarrow r) \equiv p \rightarrow (q \wedge r)$

Note that by the conditional laws, we can rewrite the RHS as:

$$\neg p \vee (q \wedge r)$$

And now, we can start from the LHS and try to get to the new RHS

# Proving Logical Equivalence

**Exercise:** Prove  $(p \rightarrow q) \wedge (p \rightarrow r) \equiv \neg p \vee (q \wedge r)$

Equivalence	Name
$p \wedge \top \equiv p$ $p \vee \text{F} \equiv p$	Identity laws
$p \vee \top \equiv \top$ $p \wedge \text{F} \equiv \text{F}$	Domination laws
$p \vee p \equiv p$ $p \wedge p \equiv p$	Idempotent laws
$\neg(\neg p) \equiv p$	Double negation law
$p \vee q \equiv q \vee p$ $p \wedge q \equiv q \wedge p$	Commutative laws
$(p \vee q) \vee r \equiv p \vee (q \vee r)$ $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$	Associative laws
$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$ $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	Distributive laws
$\neg(p \vee q) \equiv \neg p \wedge \neg q$ $\neg(p \wedge q) \equiv \neg p \vee \neg q$	De Morgan's laws
$p \vee (p \wedge q) \equiv p$ $p \wedge (p \vee q) \equiv p$	Absorption laws
$p \vee \neg p \equiv \top$ $p \wedge \neg p \equiv \text{F}$	Negation laws

Equivalences with Implication
$p \rightarrow q \equiv \neg p \vee q$
$p \rightarrow q \equiv \neg q \rightarrow \neg p$
$p \vee q \equiv \neg p \rightarrow q$
$p \wedge q \equiv \neg(p \rightarrow \neg q)$
$\neg(p \rightarrow q) \equiv p \wedge \neg q$
$(p \rightarrow q) \wedge (p \rightarrow r) \equiv p \rightarrow (q \wedge r)$
$(p \rightarrow q) \wedge (q \rightarrow r) \equiv (p \vee q) \rightarrow r$
$(p \rightarrow q) \vee (p \rightarrow r) \equiv p \rightarrow (q \vee r)$
$(p \rightarrow q) \vee (q \rightarrow r) \equiv (p \wedge q) \rightarrow r$

# Proving Logical Equivalence

**Exercise:** Prove  $(p \rightarrow q) \wedge (p \rightarrow r) \equiv \neg p \vee (q \wedge r)$

$(p \rightarrow q) \wedge (p \rightarrow r) \equiv (\neg p \vee q) \wedge (p \rightarrow r)$  by Conditional Law

# Proving Logical Equivalence

**Exercise:** Prove  $(p \rightarrow q) \wedge (p \rightarrow r) \equiv \neg p \vee (q \wedge r)$

$$\begin{aligned}(p \rightarrow q) \wedge (p \rightarrow r) &\equiv (\neg p \vee q) \wedge (p \rightarrow r) && \text{by Conditional Law} \\ &\equiv (\neg p \vee q) \wedge (\neg p \vee r) && \text{by Conditional Law}\end{aligned}$$

# Proving Logical Equivalence

**Exercise:** Prove  $(p \rightarrow q) \wedge (p \rightarrow r) \equiv \neg p \vee (q \wedge r)$

$$\begin{aligned}(p \rightarrow q) \wedge (p \rightarrow r) &\equiv (\neg p \vee q) \wedge (p \rightarrow r) && \text{by Conditional Law} \\ &\equiv (\neg p \vee q) \wedge (\neg p \vee r) && \text{by Conditional Law} \\ &\equiv \neg p \vee (q \wedge r) && \text{by Distributive Law}\end{aligned}$$

# Propositional Satisfiability

A compound proposition is **satisfiable** if there is an assignment of truth values to its variables that makes it true.

A compound proposition is **unsatisfiable** when no such assignment exists

- A compound proposition is unsatisfiable iff its negation is a tautology
- An assignment of truth values that make a compound proposition true is called a solution to that satisfiability problem



# Propositional Satisfiability

Example:  $(p \vee \neg q) \wedge (q \vee \neg r) \wedge (r \vee \neg p)$  is true when  $p, q, r$  have the same truth value.

$p = T, q = T, r = T$  is a solution to this satisfiability problem

$p = F, q = F, r = F$  is also a solution

# A Few Tips...

- Two aspects to getting comfortable with logical equivalence proofs
  - Recognizing the patterns required to be able to apply a particular law
  - Knowing which laws to actually apply
- Both require PRACTICE