

CSE 191

Introduction to Discrete Structures

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Predicates and Quantifiers

Recap

- Propositions
 - Declarative statements – either TRUE or FALSE
 - Can represent them with propositional variables
 - Can be combined/modified with logical operators
- Truth tables
 - Lists all possible combinations of truth values for the atomics in a compound proposition and the resulting truth value
- Logical Equivalence
 - Two compound propositions are logically equivalent if they have the same truth value no matter the truth values of their atomics
 - Can be proven with truth tables, or by applying laws of equivalence

Outline

Predicates and Quantifiers

- **From Propositions to Predicates**
- Quantifiers

From Propositions to Predicates

Consider the statement: "X is even"

- Contains a variable, **X**, so it is not a proposition
 - Given a value for **X**, we can determine the truth value
 - Once we know the **X**, the statement is TRUE or FALSE, but not both
- Sentences whose truth value is based on variables are **predicates**

Predicates

A **predicate** is a *function* that takes some *variable(s)* as arguments; it returns either TRUE or FALSE, but never both, depending on the combination of the combination of values passed as arguments.

a proposition can be thought of as a function of 0 variables

Predicates

A **predicate** is a *function* that takes some *variable(s)* as arguments; it returns either TRUE or FALSE, but never both, depending on the combination of the combination of values passed as arguments.

Example: $P(x)$: x is an even number.

P is the function, x is the variable

$P(x)$ is the value of the predicate P at x

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P is the function, x is the variable

$P(x)$ is the value of the predicate P at x

What are the possible values of X ?

Domain of Discourse

Given a predicate, $P(x)$, the domain of discourse (often just called the domain) is the set of all possible values for the variable x .

- Predicates with multiple variables may have:
 - Multiple domains of discourse (one for each variable)
 - A single domain of discourse for all variables

Examples

Example 1

Consider the predicate $Q(x,y)$: y is enrolled in recitation x .

- We can define the domain of discourse of x as $\{C1, C2, C3\}$
- ...and the domain of discourse of y as $\{\text{all students in CSE 191}\}$

Example 2

Consider the predicate $R(x,y)$: x and y are friends.

- We can define the domain of discourse of R as $\{\text{all students at UB}\}$

Examples (Predicate vs Proposition)

Let the predicate $P(x,y)$ be defined by:

$P(x,y)$: $2x = y$, where the domain for x is $\{1,2,3\}$, y is $\{4,5,6\}$

Statement	Is Proposition?	Truth Value
$P(1,4): 2(1) = 4$		

Examples (Predicate vs Proposition)

Let the predicate $P(x,y)$ be defined by:

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Statement	Is Proposition?	Truth Value
$P(1,4): 2(1) = 4$	Yes	

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$P(2,4)$: $2(2) = 4$		

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What about $P(2,3)$: $2(2) = 3$?

Examples (Predicate vs Proposition)

Let the predicate $P(x,y)$ be defined by:

$P(x,y)$: $2x = y$, where the domain for x is $\{1,2,3\}$, y is $\{4,5,6\}$

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$P(2,4)$: $2(2) = 4$	Yes	TRUE
$P(x,4)$: $2(x) = 4$	No	—

$P(2,3)$ is meaningless (in this example). 3 is not in the domain of y .

Examples (Predicate vs Proposition)

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$P(x,y)$: $2x = y$, where the domain for x is $\{1,2,3\}$, y is $\{4,5,6\}$

Statement	Is Proposition?	Truth Value
$P(1,4) \vee P(3,6)$		
$P(1,4) \vee \neg P(3,6)$		
$P(2,4) \rightarrow P(2,5)$		
$P(2,4) \wedge P(x,4)$		

Examples (Predicate vs Proposition)

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$P(2,4) \wedge P(x,4)$		

Examples (Predicate vs Proposition)

Let the predicate $P(x,y)$ be defined by:

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$P(2,4) \rightarrow P(2,5)$	Yes	FALSE
$P(2,4) \wedge P(x,4)$	No	—

More Examples

Let the predicate $Q(x,y)$ be defined by:

$Q(x,y): x + y > 4$, where the domain for x and y is **all integers**

Which of the following are predicates? Propositions?

- $Q(1,2)$
- $Q(x,2)$
- $Q(1000,y)$
- $Q(1000,2)$
- $Q(x,y)$

More Examples

Let the predicate $Q(x,y)$ be defined by:

$Q(x,y): x + y > 4$, where the domain for x and y is **all integers**

Which of the following are predicates? Propositions?

- $Q(1,2)$ proposition
- $Q(x,2)$ predicate
- $Q(1000,y)$ predicate
- $Q(1000,2)$ proposition
- $Q(x,y)$ predicate

Outline

Predicates and Quantifiers

- From Propositions to Predicates
- **Quantifiers**

Quantifiers

Quantification expresses the extent to which a predicate is true over a range of elements. For example, in English: all, some, none, many, few, ...

Universal Quantification

Suppose $P(x)$ is a predicate on some domain, D .

The universal quantification of $P(x)$ is the *proposition*:

" $P(x)$ is true for all x in the domain of discourse D ."

Written as: $\forall x, P(x)$

Read as: "For all $x, P(x)$ " or "For every $x, P(x)$ "

$\forall x, P(x)$ is TRUE if $P(x)$ is TRUE for every x in D .

$\forall x, P(x)$ is FALSE if $P(x)$ is FALSE for some x in D .

Universal Quantification Example

$P(x): x + 2 = 5$, where the domain of discourse is $\{1,2,3\}$

$\forall x, P(x)$ means: "for all x in $\{1,2,3\}$, $x + 2 = 5$ "

$$\begin{aligned}\forall x, P(x) &\equiv P(1) \wedge P(2) \wedge P(3) \\ &\equiv (1 + 2 = 5) \wedge (2 + 2 = 5) \wedge (3 + 2 = 5) \\ &\equiv F \wedge F \wedge T\end{aligned}$$

$\therefore \forall x, P(x)$ is FALSE (since $1 + 2 = 5$ and $2 + 2 = 5$ are FALSE)


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 Note that \therefore denotes "therefore"

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An input to that causes a universally quantified statement to evaluate to false is called a **counterexample**

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$\therefore \forall x, P(x)$ is FALSE (since $1 + 2 = 5$ and $2 + 2 = 5$ are FALSE)

Note that \therefore denotes "therefore"

What if the domain of discourse was $\{3\}$?

Universal Quantification Example

Consider:

$A(x)$: x is even

$B(x)$: $x^2 > 0$

$C(x)$: $x < 2$

where the domain of A, B, C is $\{0,1,2,3\}$

True or False?

$\forall x, (C(x) \rightarrow A(x))$

$\forall x, (B(x) \vee C(x))$

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FALSE, counterexample is $x = 1$

$\forall x, (B(x) \vee C(x))$

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True or False?

$\forall x, (C(x) \rightarrow A(x))$

FALSE, counterexample is $x = 1$

$\forall x, (B(x) \vee C(x))$

TRUE

x	$B(x)$	$C(x)$
0	F	T
1	T	T
2	T	F
3	T	F

Universal Quantification Example

Consider:

$S(x)$: x is a student in CSE 191

$T(x)$: x is a CSE major

where the domain for S and T is
{all students enrolled in CSE 191}

True or False?

$$\forall x, S(x)$$

$$\forall x, T(x)$$

$$\forall x, (S(x) \rightarrow T(x))$$

$$\forall x, (T(x) \rightarrow S(x))$$

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True or False?

$\forall x, S(x)$ TRUE

$\forall x, T(x)$ FALSE

$\forall x, (S(x) \rightarrow T(x))$ FALSE

$\forall x, (T(x) \rightarrow S(x))$ TRUE

Existential Quantification

Suppose $P(x)$ is a predicate on some domain, D .

The existential quantification of $P(x)$ is the *proposition*:

" $P(x)$ is true for some x in the domain of discourse D ."

Written as: $\exists x, P(x)$

Read as: "There exists an x such that, $P(x)$ " or "For some x , $P(x)$ "

$\exists x, P(x)$ is TRUE if $P(x)$ is TRUE for some x in D .

$\exists x, P(x)$ is FALSE if $P(x)$ is FALSE for every x in D .

Existential Quantification Example

$P(x): x + 2 = 5$, where the domain of discourse is $\{1,2,3\}$

$\exists x, P(x)$ means: "for some x in $\{1,2,3\}$, $x + 2 = 5$ "

$$\begin{aligned}\exists x, P(x) &\equiv P(1) \vee P(2) \vee P(3) \\ &\equiv (1 + 2 = 5) \vee (2 + 2 = 5) \vee (3 + 2 = 5) \\ &\equiv F \vee F \vee T\end{aligned}$$

$\therefore \exists x, P(x)$ is TRUE (because $3 + 2 = 5$ is TRUE)

Existential Quantification Example

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$\exists x, P(x)$ means: "for some x in $\{1,2,3\}$, $x + 2 = 5$ "

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$\therefore \exists x, P(x)$ is TRUE (because $3 + 2 = 5$ is TRUE)

An input to that causes a predicate to evaluate to true is called a **satisfying assignment**

Existential Quantification Example

Consider:

$$A(x): x = 1$$

$$B(x): x > 5$$

$$C(x): x < 5$$

where the domain of **A**, **B**, **C** is $\{1,2,3\}$

True or False?

$$\exists x, (C(x) \rightarrow A(x))$$

$$\exists x, B(x)$$

Existential Quantification Example

Consider:

$$A(x): x = 1$$

$$B(x): x > 5$$

$$C(x): x < 5$$

where the domain of **A**, **B**, **C** is $\{1,2,3\}$

True or False?

$$\exists x, (C(x) \rightarrow A(x))$$

TRUE, satisfying assignment $x=1$

$$\exists x, B(x)$$

FALSE

$$P(1) \vee P(2) \vee P(3) = F \vee F \vee F \\ = F$$

Note: The existential and universal quantifiers have higher precedence than all logical operators

Existential Quantification Example

Consider:

$S(x)$: x is a student in CSE 191

$T(x)$: x is a CSE major

where the domain for S and T is
{all students enrolled in CSE 191}

True or False?

$\exists x, S(x)$ TRUE

$\exists x, T(x)$ TRUE

$\exists x, (S(x) \rightarrow T(x))$ TRUE

$\exists x, (T(x) \rightarrow S(x))$ TRUE

Quantifier Example

Consider:

$$P(x): x^2 > 9$$

$$Q(x): x^2 \geq 0$$

where the domain is *all integers*

True or False?

$$\forall x, P(x)$$

$$\exists x, P(x)$$

$$\forall x, Q(x)$$

$$\exists x, Q(x)$$

Quantifier Example

Consider:

$$P(x): x^2 > 9$$

$$Q(x): x^2 \geq 0$$

where the domain is *all integers*

True or False?

$\forall x, P(x)$ FALSE, consider: $x = 2$

$\exists x, P(x)$ TRUE, consider: $x = 5$

$\forall x, Q(x)$ TRUE, consider...math

$\exists x, Q(x)$ TRUE

Binding Variables

The occurrence of a variable, x , is said to be **bound** when a quantifier is used on that variable.

The occurrence of a variable, x , is said to be **free** when it is not bound by a quantifier or set to a particular variable.

The part of a logical expression to which a quantifier is applied is called the **scope** of this quantifier.

Note: A variable is free if it is outside of the scope of all quantifiers in the formula that specify this variable.

Binding Example

Consider: $\exists x, (x + y = 1)$

- The variable x is bound by the existential quantifier $\exists x$
- y is free

Consider: $\forall x, (x < 9) \vee (x > 9)$

- The variable x is bound in $(x < 9)$ by $\forall x$
- The variable x is free in $(x > 9)$

Binding Example

Consider: $\exists x, (x + y = 1)$

- The variable x is bound by the existential quantifier $\exists x$
- y is free

Consider: $\forall x, (x < 9) \vee (x > 9)$

- The variable x is bound in $(x < 9)$ by $\forall x$
- The variable x is free in $(x > 9)$

The precedence of \forall and \exists is higher than logical operators

Quantified Statements and English

Consider:

$L(x,y)$: x loves y

where the domain of x is {all students enrolled in CSE 191} and
where the domain of y is {all courses offered by UB CSE}

$$\exists x, (L(x, \text{CSE 191}) \wedge L(x, \text{CSE 250}))$$

Quantified Statements and English

Consider:

$L(x,y)$: x loves y

where the domain of x is {all students enrolled in CSE 191} and
where the domain of y is {all courses offered by UB CSE}

$$\exists x, (L(x, \text{CSE 191}) \wedge L(x, \text{CSE 250}))$$

A CSE 191 student loves both CSE 191 and CSE 250

Quantified Statements and English

Consider:

$L(x,y)$: x loves y

where the domain of x is {all students enrolled in CSE 191} and
where the domain of y is {all courses offered by UB CSE}

$$\exists x \exists y \forall z, ((x \neq y) \wedge (L(x, z) \rightarrow L(y, z)))$$

Quantified Statements and English

Consider:

$L(x,y)$: x loves y

where the domain of x is {all students enrolled in CSE 191} and
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$$\exists x \exists y \forall z, ((x \neq y) \wedge (L(x, z) \rightarrow L(y, z)))$$

There are 2 different students, x and y , in CSE 191 such that if x loves a CSE course, then so does y .

Quantified Statements and English

Consider:

$L(x,y)$: x loves y

where the domain of x is {all students enrolled in CSE 191} and
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Every CSE course is loved by some CSE 191 student

Quantified Statements and English

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where the domain of x is {all students enrolled in CSE 191} and
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Every CSE course is loved by some CSE 191 student

$$\forall y \exists x, L(x, y)$$

Quantified Statements and English

Consider:

$A(x)$: x lives in Amherst

$B(x)$: x is a CSE 191 student

$C(x)$: x has a good GPA

$D(x)$: x majors in CSE

Domain of discourse: all UB students

All UB students are CSE 191 students

Quantified Statements and English

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Domain of discourse: all UB students

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$\forall x, B(x)$

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$$\forall x, (B(x) \rightarrow C(x))$$

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Domain of discourse: all UB students

CSE 191 students not living in Amherst major in CSE

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B(x): x is a CSE 191 student

C(x): x has a good GPA

D(x): x majors in CSE

Domain of discourse: all UB students

CSE 191 students not living in Amherst major in CSE

$$\forall x, ((B(x) \wedge \neg A(x)) \rightarrow D(x))$$

Quantified Statements and English

Consider:

$A(x)$: x lives in Amherst

$B(x)$: x is a CSE 191 student

$C(x)$: x has a good GPA

$D(x)$: x majors in CSE

Domain of discourse: all UB students

No CSE 191 student lives in Amherst

Quantified Statements and English

Consider:

A(x): x lives in Amherst

B(x): x is a CSE 191 student

C(x): x has a good GPA

D(x): x majors in CSE

Domain of discourse: all UB students

No CSE 191 student lives in Amherst

$$\forall x, (B(x) \rightarrow \neg A(x))$$

Exercise

Translate the following theorems to quantified statements

If x is an even number, then $x + 1$ is odd

Every even number is a multiple of 2

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Translate the following theorems to quantified statements

If x is an even number, then $x + 1$ is odd

Every even number is a multiple of 2

First we must define domain/predicates

- Domain: all integers
- $P(x)$: x is an even number
- $Q(x)$: x is an odd number

Exercise

Translate the following theorems to quantified statements

If x is an even number, then $x + 1$ is odd

Every even number is a multiple of 2

First we must define domain/predicates

- Domain: all integers
- $P(x)$: x is an even number
- $Q(x)$: x is an odd number

$$\forall x, (P(x) \rightarrow Q(x + 1))$$

Exercise

Translate the following theorems to quantified statements

If x is an even number, then $x + 1$ is odd

Every even number is a multiple of 2

First we must define domain/predicates

- Domain: all integers
- $P(x)$: x is an even number
- $Q(x)$: x is an odd number

Domain and predicates:

- Domain: all integers
- $R(y)$: y is an even number
- $S(y)$: y is a multiple of 2

$$\forall x, (P(x) \rightarrow Q(x + 1))$$

Exercise

Translate the following theorems to quantified statements

If x is an even number, then $x + 1$ is odd

Every even number is a multiple of 2

First we must define domain/predicates

Domain and predicates:

- Domain: all integers
- $P(x)$: x is an even number
- $Q(x)$: x is an odd number

- Domain: all integers
- $R(y)$: y is an even number
- $S(y)$: y is a multiple of 2

$$\forall x, (P(x) \rightarrow Q(x + 1))$$

$$\forall y, (R(y) \rightarrow S(y))$$

Exercise

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Domain and predicates:

- Domain: all integers
- $P(x)$: x is an even number
- $Q(x)$: x is an odd number

- Domain: all integers
- $R(y)$: y is an even number
- $S(y)$: y is a multiple of 2

$$\forall x, (P(x) \rightarrow Q(x + 1))$$

$$\forall y, (R(y) \rightarrow S(y))$$

How can we do this differently?
With one predicate?

Exercise

Translate the following theorems to quantified statements

If x is an even number, then $x + 1$ is odd

Every even number is a multiple of 2

First we must define domain/predicates

Domain and predicates:

- Domain: all integers
- $P(x)$: x is an even number
- $Q(x)$: x is an odd number

- Domain: all even integers
- $T(z)$: z is a multiple of 2

$$\forall x, (P(x) \rightarrow Q(x + 1))$$

$$\forall z, T(z)$$

Quantifier Negation

Consider:

$D(x)$: x majors in CSE

Domain of discourse:

all UB students

Consider the following:

Not every UB student majors in
computer science.

Some UB students do not major
in computer science.

Quantifier Negation

Consider:

$D(x)$: x majors in CSE

Domain of discourse:

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Consider the following:

Not every UB student majors in
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$$\neg \forall x, D(x)$$

Some UB students do not major
in computer science.

Quantifier Negation

Consider:

$D(x)$: x majors in CSE

Domain of discourse:

all UB students

Consider the following:

Not every UB student majors in
computer science.

$$\neg \forall x, D(x)$$

Some UB students do not major
in computer science.

$$\exists x, \neg D(x)$$

Quantifier Negation Rule

Quantifier Negation

In general we have, for any predicate $P(x)$:

$$\neg \forall x, P(x) \equiv \exists x, \neg P(x) \text{ and } \neg \exists x, P(x) \equiv \forall x, \neg P(x)$$

De Morgan's Law for Quantifiers

Negation	Equivalent statement	Negation is TRUE when...	Negation is FALSE when...
$\neg \exists x, P(x)$	$\forall x, \neg P(x)$	For every x, P(x) is FALSE	There is an x where P(x) is TRUE
$\neg \forall x, P(x)$	$\exists x, \neg P(x)$	There is an x where P(x) is FALSE	P(x) is true for every x

Quantifier Negation Examples

Negate the following statements and simplify

$$\exists x, (P(x) \rightarrow \neg Q(x))$$

$$\forall x, (P(x) \rightarrow \exists y(P(y) \vee Q(y)))$$

Quantifier Negation Examples

Negate the following statements and simplify

$$\exists x, (P(x) \rightarrow \neg Q(x))$$

$$\neg \exists x, (P(x) \rightarrow \neg Q(x))$$

$$\forall x, \neg(P(x) \rightarrow \neg Q(x))$$

$$\forall x, \neg(\neg P(x) \vee \neg Q(x))$$

$$\forall x, (\neg\neg P(x) \wedge \neg\neg Q(x))$$

$$\forall x, (P(x) \wedge Q(x))$$

$$\forall x, (P(x) \rightarrow \exists y(P(y) \vee Q(y)))$$

$$\neg \forall x, (P(x) \rightarrow \exists y(P(y) \vee Q(y)))$$

$$\exists x, \neg(P(x) \rightarrow \exists y(P(y) \vee Q(y)))$$

$$\exists x, \neg(\neg P(x) \vee \exists y(P(y) \vee Q(y)))$$

$$\exists x, (\neg\neg P(x) \wedge \neg \exists y(P(y) \vee Q(y)))$$

$$\exists x (P(x) \wedge \forall y \neg(P(y) \vee Q(y)))$$

Nested Quantifiers

A logical expression with more than one quantifier that bind different variables in the same predicate is said to have **nested quantifiers**.

In order to evaluate them we must consider their *ordering* and *scope*

Nested Quantifiers: Ordering

The order of the quantifiers is important
(unless they are all universal or all existential quantifiers)

Example

$Q(x,y): x + y = 0$, where the domain is all real numbers

- $\exists y \forall x Q(x,y)$: there is a real number y , such that for every real number x , $x + y = 0$.
- $\forall x \exists y Q(x,y)$: for every real number x , there is a real number y such that $x + y = 0$.

Nested Quantifiers: Ordering

The order of the quantifiers is important
(unless they are all universal or all existential quantifiers)

Example

$Q(x,y): x + y = 0$, where the domain is all real numbers

- $\exists y \forall x Q(x,y)$: there is a real number y , such that for every real number x , $x + y = 0$. **FALSE**
- $\forall x \exists y Q(x,y)$: for every real number x , there is a real number y such that $x + y = 0$. **TRUE**

In general, we cannot switch the ordering and guarantee equivalence

Nested Quantifiers: Ordering

Consecutive quantifiers of the same type can be reordered while maintaining equivalence

Consider predicate $Q(i,j,k)$

$$\forall i \forall j \forall k Q(i,j,k) \equiv \forall j \forall k \forall i Q(i,j,k) \equiv \forall j \forall i \forall k Q(i,j,k) \equiv \forall i \forall k \forall j Q(i,j,k)$$

...

$$\exists i \exists j \exists k Q(i,j,k) \equiv \exists j \exists k \exists i Q(i,j,k) \equiv \exists j \exists i \exists k Q(i,j,k) \equiv \exists i \exists k \exists j Q(i,j,k)$$

...

Nested Quantifiers: Ordering

We can simplify consecutive variables with the same quantifier

$$\forall i \forall j \forall k Q(i,j,k) \equiv \forall i,j,k, Q(i,j,k)$$

$$\exists i \exists j \exists k Q(i,j,k) \equiv \exists i, j, k, Q(i,j,k)$$

...but do so carefully and make sure you don't accidentally mix quantifiers

Nested Quantifiers: Scope

The portion of the formula a quantifier covers is called the scope

- The scope of the quantifier is the predicate immediately following
- Precedence is just below parenthesis
- Any variable not covered by any quantifier is a free variable

Consider the following formula:

$$\forall i \exists j, (P(i,j) \rightarrow \forall k, Q(k,j))$$

Nested Quantifiers: Scope

The portion of the formula a quantifier covers is called the scope

- The scope of the quantifier is the predicate immediately following
- Precedence is just below parenthesis
- Any variable not covered by any quantifier is a free variable

Consider the following formula:

$$\forall i \exists j, (P(i,j) \rightarrow \forall k, Q(k,j))$$

The scope of $\forall i$ is the entire formula

Nested Quantifiers: Scope

The portion of the formula a quantifier covers is called the scope

- The scope of the quantifier is the predicate immediately following
- Precedence is just below parenthesis
- Any variable not covered by any quantifier is a free variable

Consider the following formula:

$$\forall i \exists j, (P(i,j) \rightarrow \forall k, Q(k,j))$$

The scope of $\exists j$ is the entire formula (other than $\forall i$)

Nested Quantifiers: Scope

The portion of the formula a quantifier covers is called the scope

- The scope of the quantifier is the predicate immediately following
- Precedence is just below parenthesis
- Any variable not covered by any quantifier is a free variable

Consider the following formula:

$$\forall i \exists j, (P(i,j) \rightarrow \boxed{\forall k, Q(k,j)})$$

The scope of $\forall k$ is limited to $Q(k,j)$

Nested Quantifiers: Scope

The portion of the formula a quantifier covers is called the scope

- The scope of the quantifier is the predicate immediately following
- Precedence is just below parenthesis
- Any variable not covered by any quantifier is a free variable

Consider the following formula:

$$\forall i \exists j, (P(i,j) \rightarrow \forall k, Q(k,j))$$

Nested Quantifiers: Scope

Quantifiers can move as long as their scope doesn't encompass additional quantifiers of a different type

$$\forall i \exists j, (P(i,j) \rightarrow \forall k, Q(k,j))$$

vs

$$\forall i \exists j \forall k, (P(i,j) \rightarrow Q(k,j))$$

Nested Quantifiers: Scope

Quantifiers can move as long as their scope doesn't encompass additional quantifiers of a different type

$$\forall i \exists j (P(i,j) \rightarrow \forall k, Q(k,j))$$

vs

$$\forall i \exists j \forall k (P(i,j) \rightarrow Q(k,j))$$

Nested Quantifiers: Scope

Quantifiers can move as long as their scope doesn't encompass additional quantifiers of a different type

$$\forall i \exists j (P(i,j) \rightarrow \forall k, Q(k,j))$$

VS

$$\forall i \exists j \forall k (P(i,j) \rightarrow Q(k,j))$$

Notice how each qualifier covers the exact same variables in either case

Nested Quantifiers: Scope

Ensure that any reordering doesn't free variables that were originally bound

$$\forall i \exists j, (P(i,j) \rightarrow \forall k, Q(k,j))$$

vs

$$\forall i \exists j, (\forall k, P(i,j) \rightarrow Q(k,j))$$

Nested Quantifiers: Scope

Ensure that any reordering doesn't free variables that were originally bound

$$\forall i \exists j, (P(i,j) \rightarrow \forall k, Q(k,j))$$

vs

$$\forall i \exists j, (\forall k, P(i,j)) \rightarrow Q(k,j)$$

Nested Quantifiers: Scope

Ensure that any reordering doesn't free variables that were originally bound

$$\forall i \exists j, (P(i,j) \rightarrow \forall k, Q(k,j))$$

VS

$$\forall i \exists j, (\forall k, P(i,j)) \rightarrow Q(k,j)$$

Notice how k is no longer bound by a quantifier in the second equation

Nested Quantifiers: Scope

Ensure that any reordering doesn't free variables that were originally bound

$$\forall i \exists j, (P(i,j) \rightarrow \forall k, Q(k,j))$$

VS

$$\forall i \exists j, (\forall k, P(i,j)) \rightarrow Q(k,j)$$

Notice how k is no longer bound by a quantifier in the second equation

These two statements are NOT equivalent. The first is a proposition, and the second is a predicate with a free variable, k