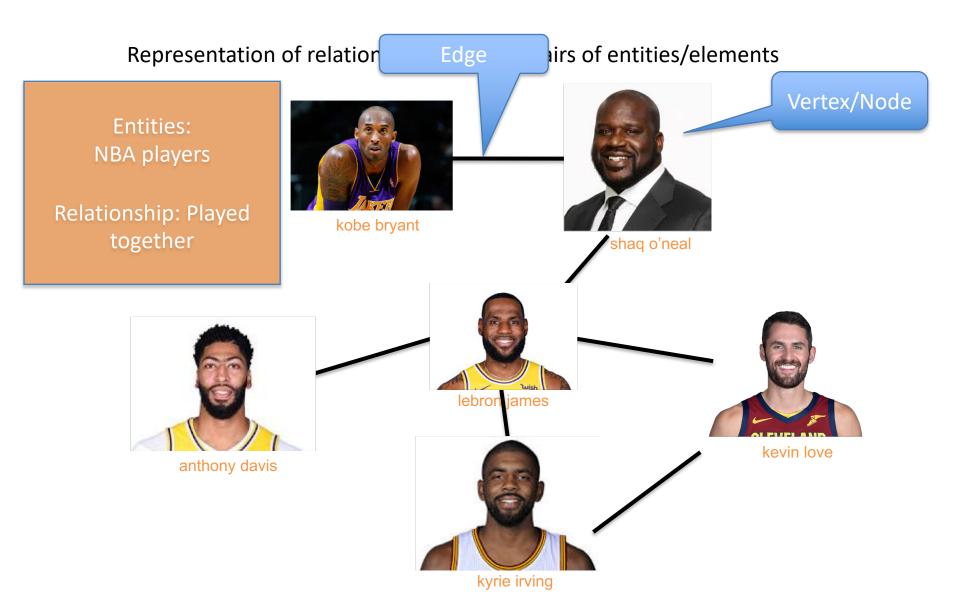
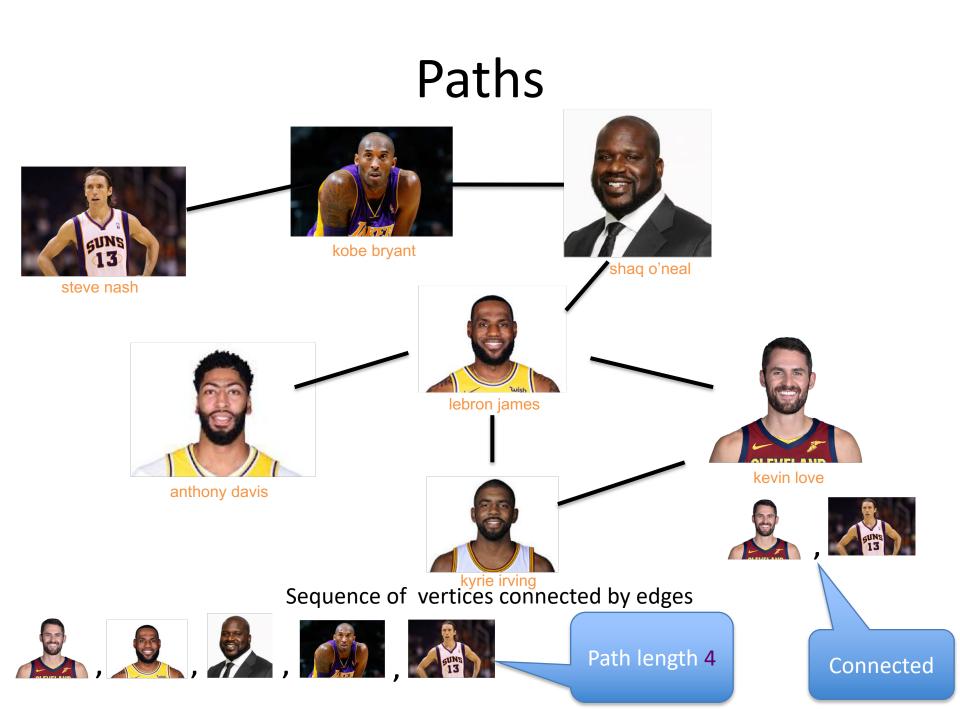
### Lecture 11

CSE 331 Feb 19, 2020

# Graphs

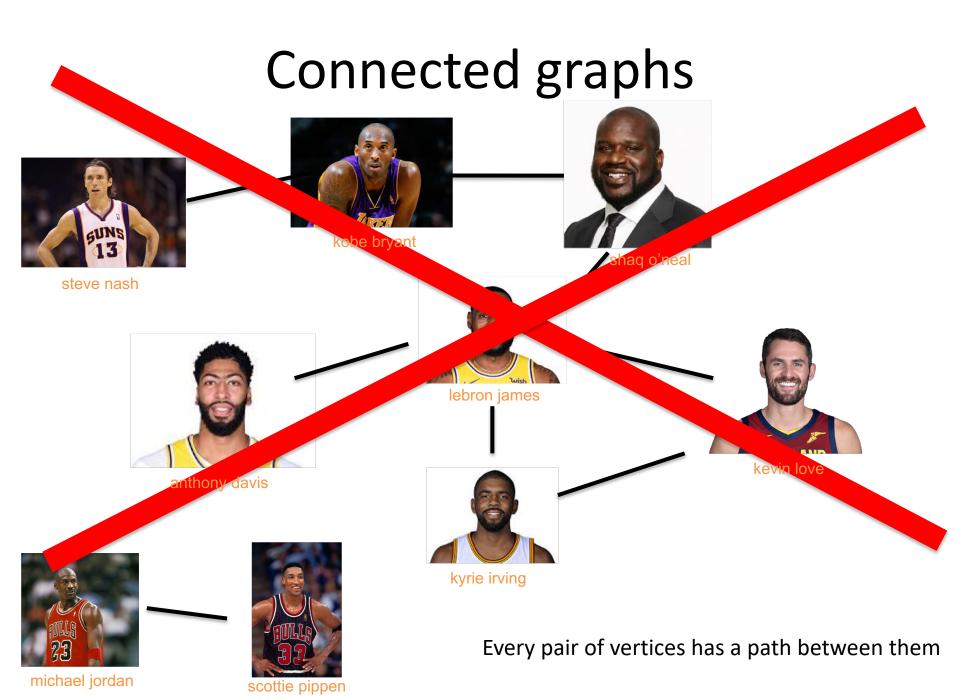


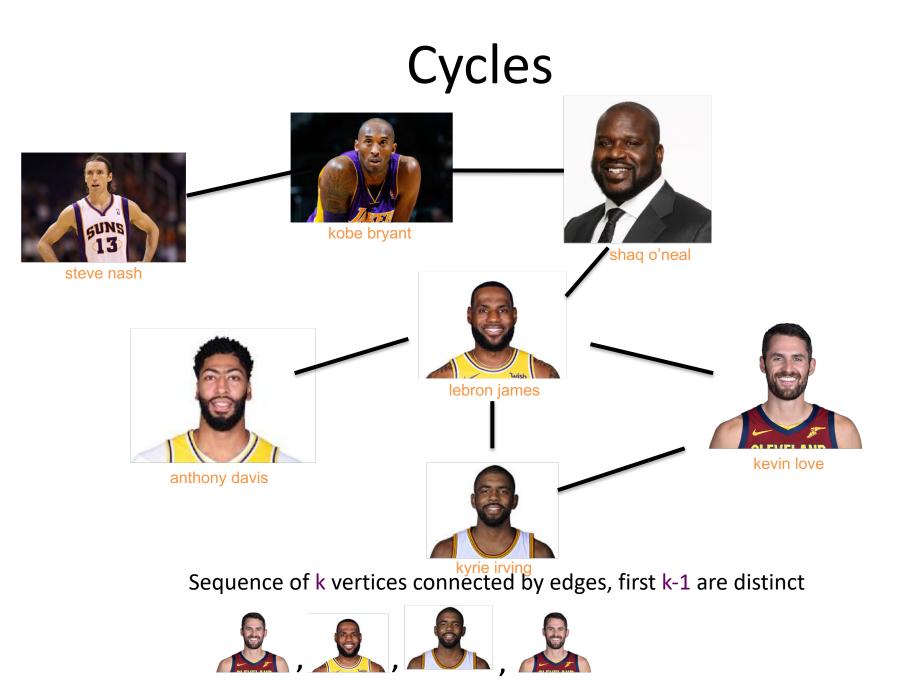


# Connectivity

u and w are connected iff there is a path between them

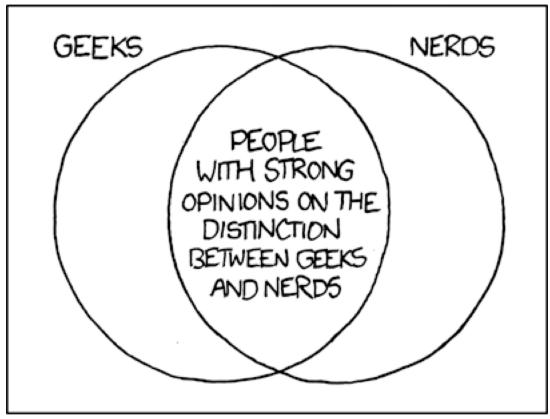
A graph is connected iff all pairs of vertices are connected





## Questions?

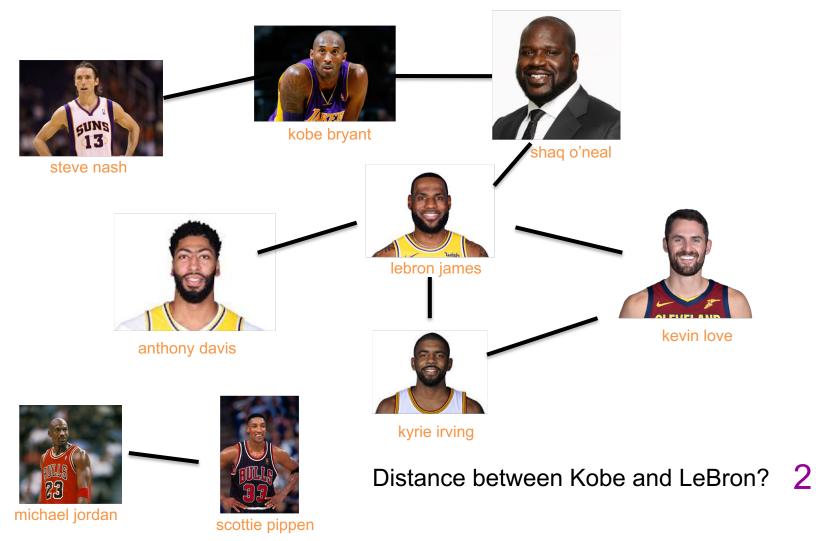
# Formally define everything



http://imgs.xkcd.com/comics/geeks\_and\_nerds.png

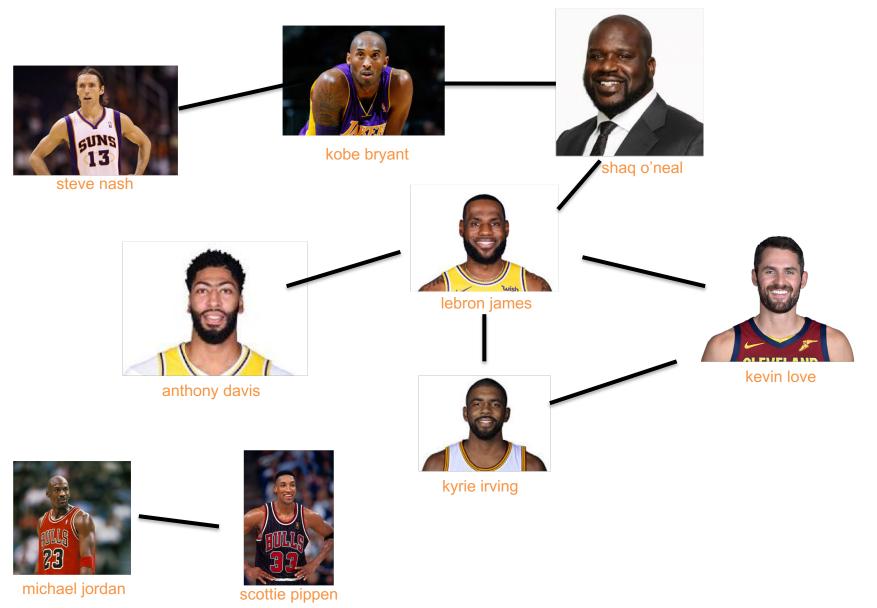
## Distance between u and v

Length of the shortest length path between  $\mathbf{u}$  and  $\mathbf{v}$ 



### Tree

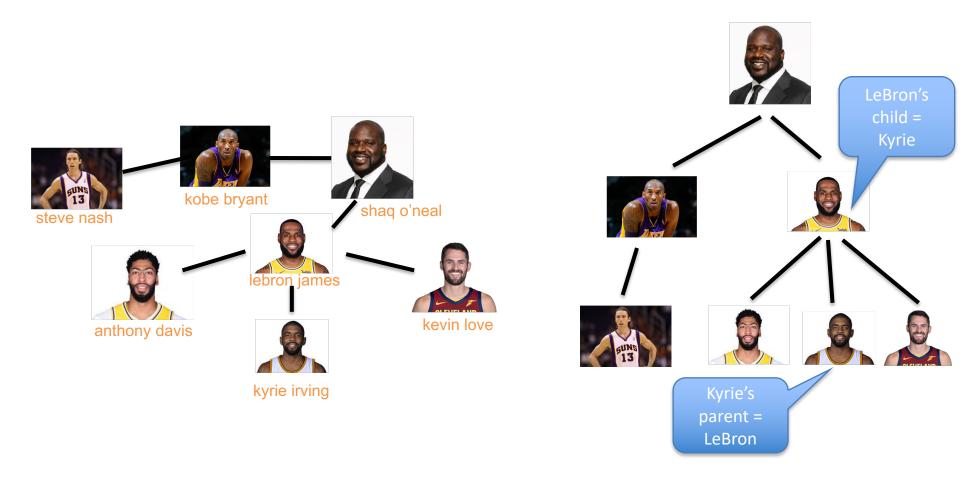
#### Connected undirected graph with no cycles



### **Rooted Tree**



### A rooted tree



Pick any vertex as root

Let the rest of the tree hang under "gravity"

### Every n vertex tree has n-1 edges

### Trees

This page collects material from previous incarnations of CSE 331 on trees, especially the proof that trees with n nodes have exactly n - 1 edges.

### Where does the textbook talk about this?

Section 3.1 in the textbook has the lowdown on trees.

### Fall 2018 material

Here is the lecture video:



### Every n vertex tree has n-1 edges

Let G be an undirected graph on n nodes

Then ANY two of the following implies the third:

T is connected

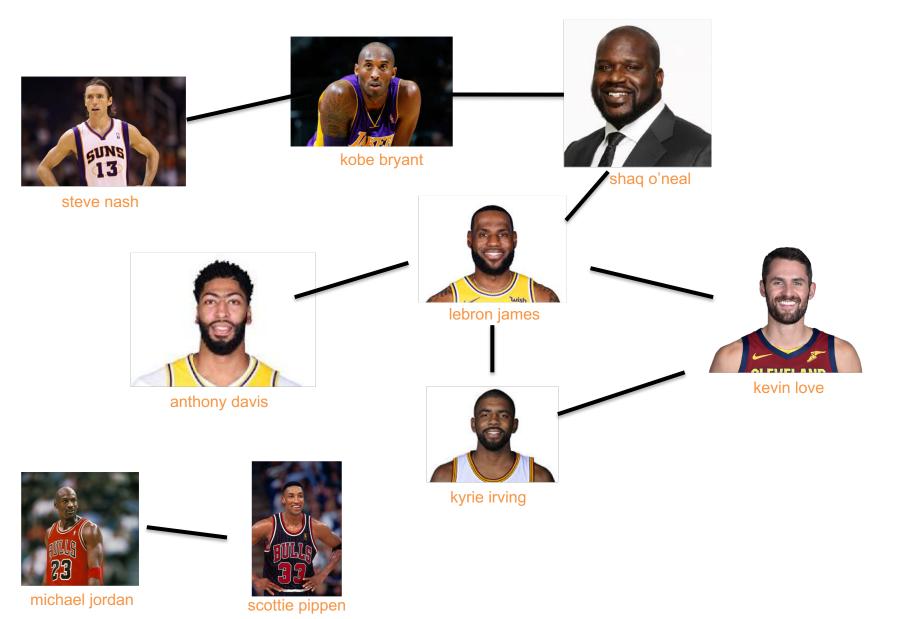
T has no cycles

T has n-1 edges

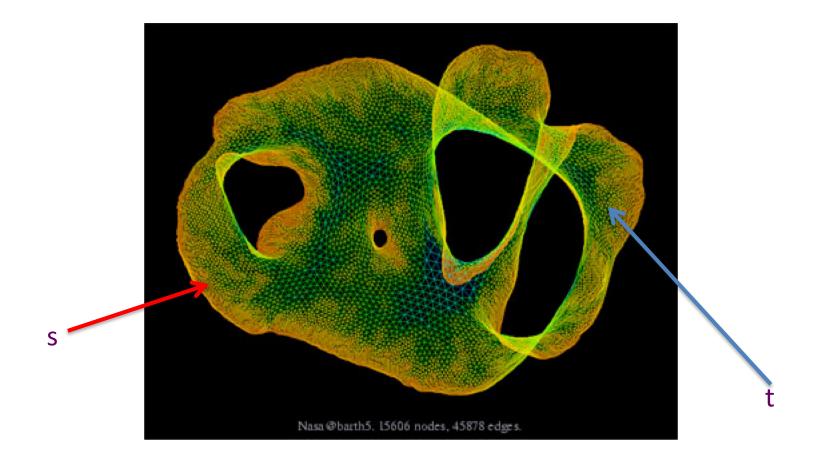
## Rest of Today's agenda

Algorithms for checking connectivity

# Checking by inspection

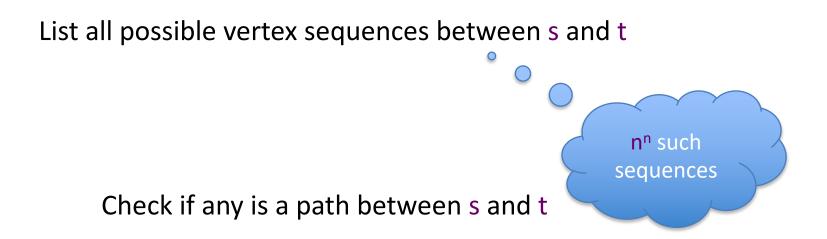


# What about large graphs?

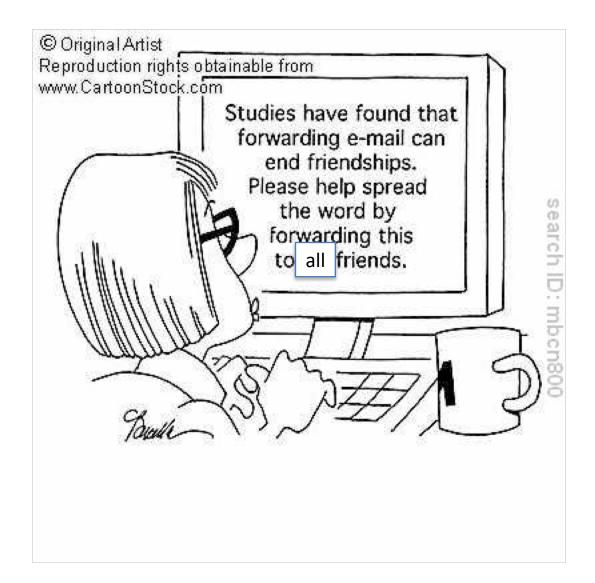


Are s and t connected?

# Brute-force algorithm?



# Algorithm motivation



# Breadth First Search (BFS)

## **BFS via examples**

In which we derive the breadth first search (BFS) algorithm via a sequence of examples.

### Expected background

These notes assume that you are familiar with the following:

- · Graphs and their representation. In particular,
  - Notion of connectivity of nodes and connected components of graphs
  - Adjacency list representation of graphs
  - Notation:
    - G = (V, E)
    - n = |V| and m = |E|
    - CC(s) denotes the connected component of s
- Trees and their basic properties

### The problem

In these notes we will solve the following problem:

### **Connectivity Problem**

*Input:* Graph G = (V,E) and s in V

Output: All t connected to s in G

## Breadth First Search (BFS)

Build layers of vertices connected to s

 $L_0 = \{s\}$ 

Assume  $L_0,...,L_i$  have been constructed

 $L_{i+1}$  is the set of vertices not chosen yet but are connected by an edge to  $L_i$ 

Stop when new layer is empty