

Mar 6

Runtime analysis:

$$l_{\max} = \max_{e \in E} l_e$$

Let n' & m'
be the #nodes,
#edges in G'

$$O(n' + m')$$

$$= O(l_{\max}(n+m))$$

$$m' \leq l_{\max} \cdot (n+m)$$

$$n' \leq l_{\max} \cdot (n+m)$$

Recap/Aside: RAM model: unit space is a register

If you have n items, each register has $O(\log n)$ bits

→ All basic ops on $O(1)$ registers is $O(1)$ time.

let $l_{\max} = n^{100} \Rightarrow 100 \cdot \lg n$ bits $\Rightarrow 100$ registers $\hookrightarrow O(1)$

→ Input size (#registers) = $O(m+n)$

runtime $O(n^{100}(n+m))$

Assume $l_{\max} = n^{o(1)} \Rightarrow O(1)$ registers for each l_e
 \Rightarrow length = $O(n+m)$

WANT: Ideally runtime $O(m+n)$