

# Lecture 11

CSE 331

Feb 24, 2021

# Graphs

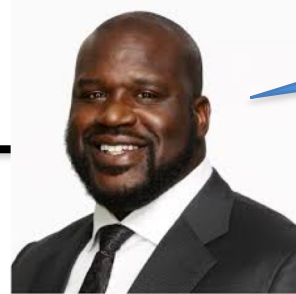
Representation of relationship pairs of entities/elements

Entities:  
NBA players

Relationship: Played together



kobe bryant



shaq o'neal



lebron james



kevin love



anthony davis



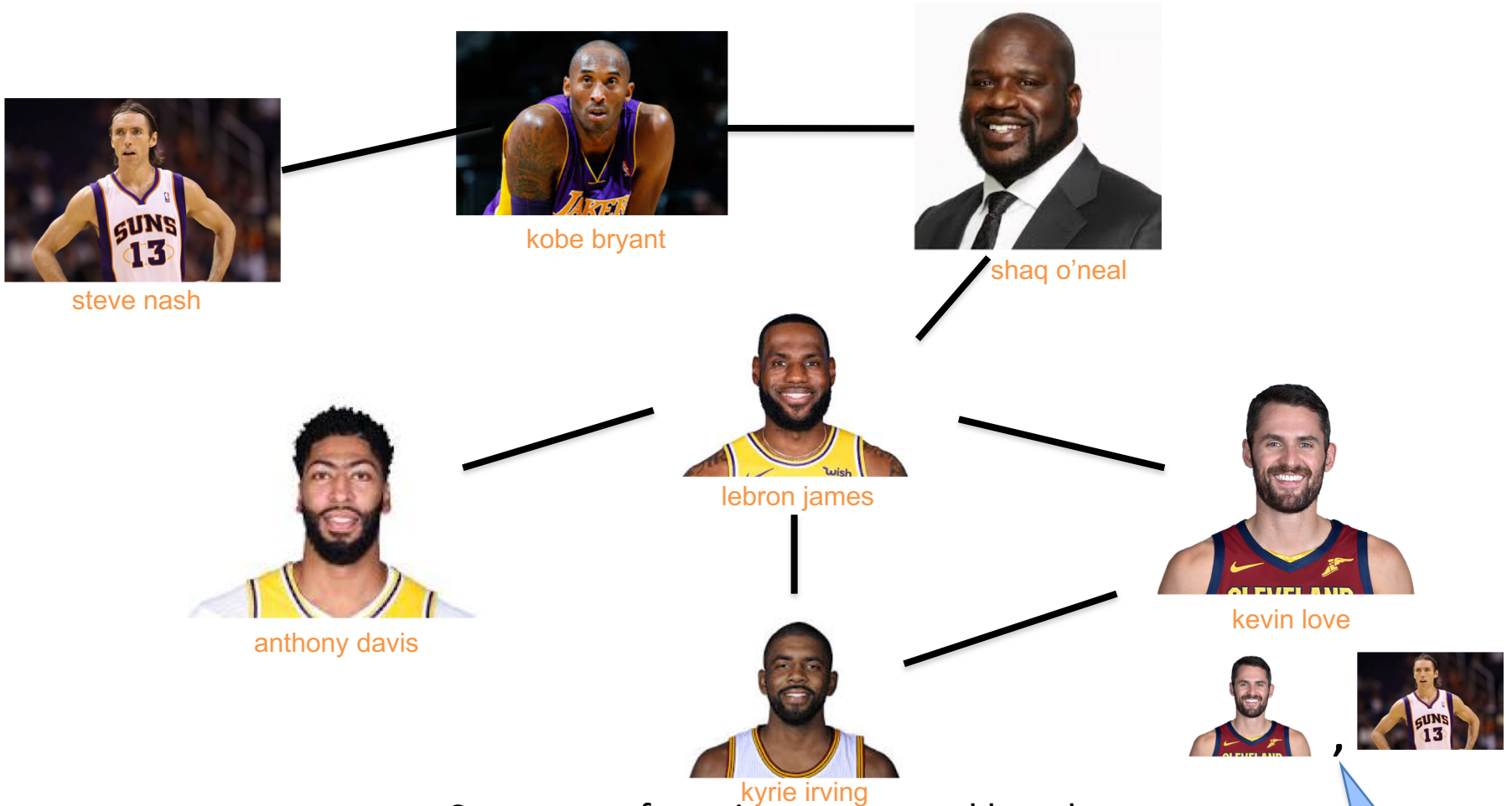
kyrie irving

Edge

Vertex/Node



# Paths



Sequence of vertices connected by edges



Path length 4

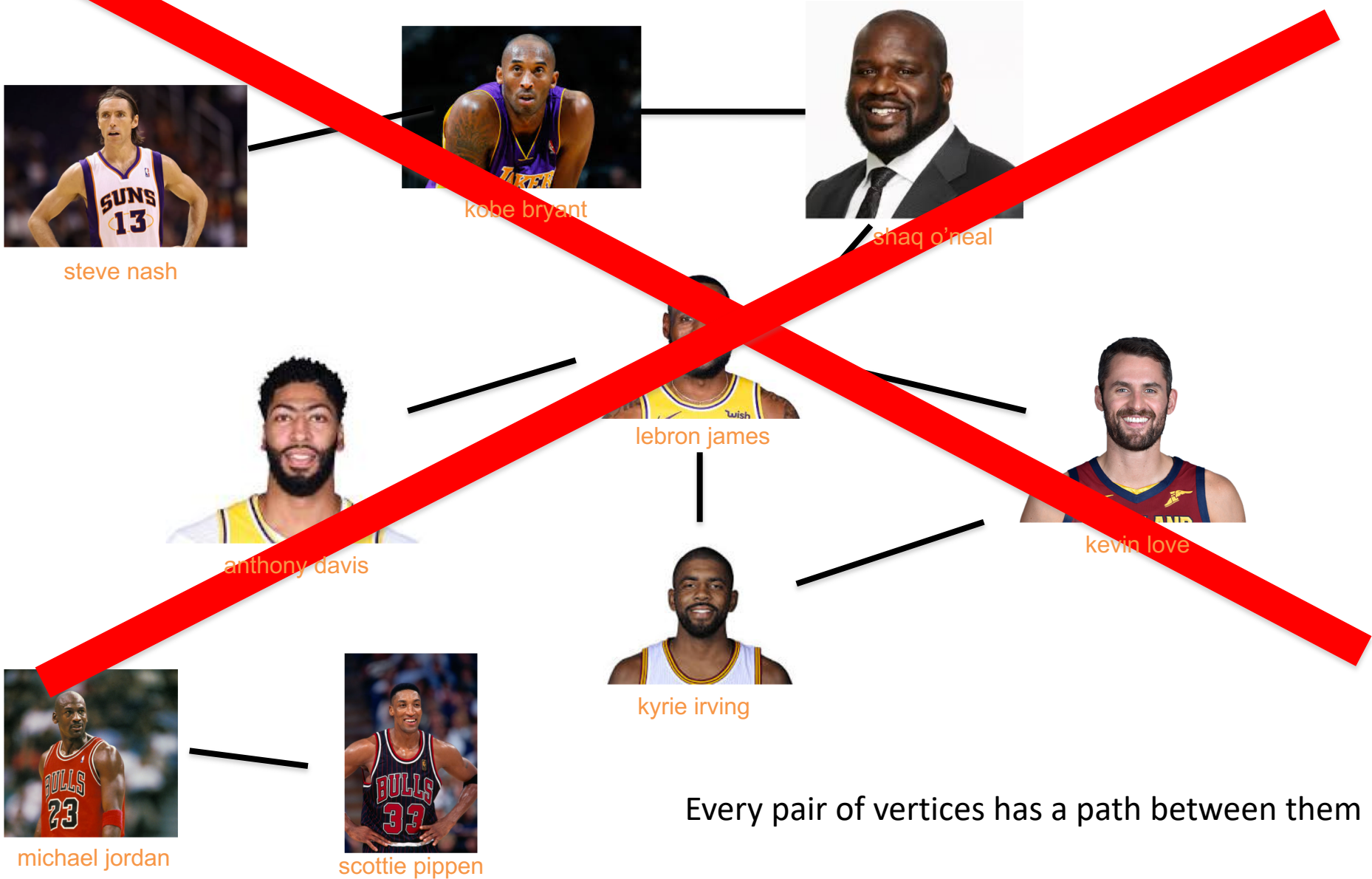
Connected

# Connectivity

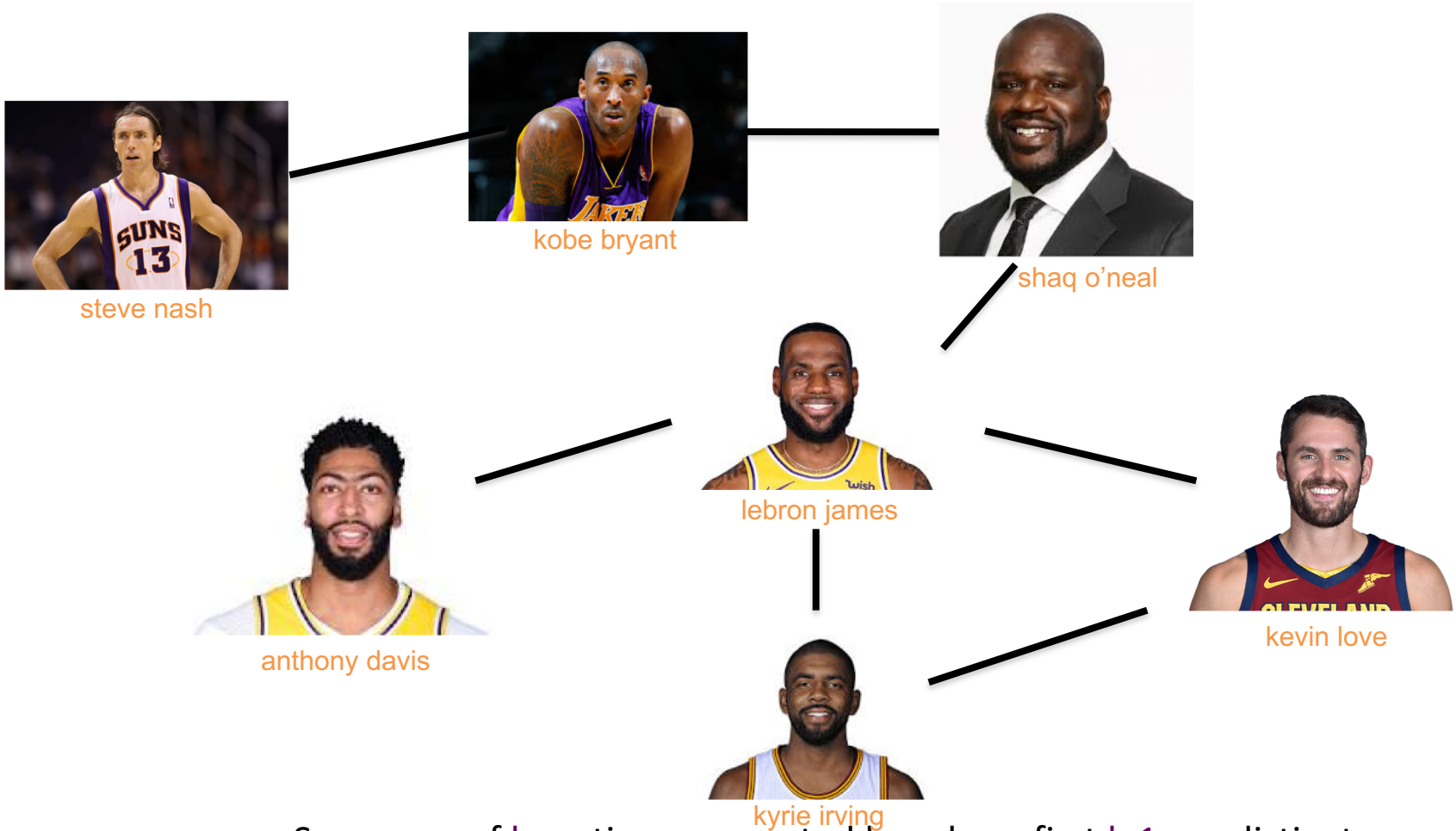
$u$  and  $w$  are connected iff there is a path between them

A graph is connected iff all pairs of vertices are connected

# Connected graphs



# Cycles

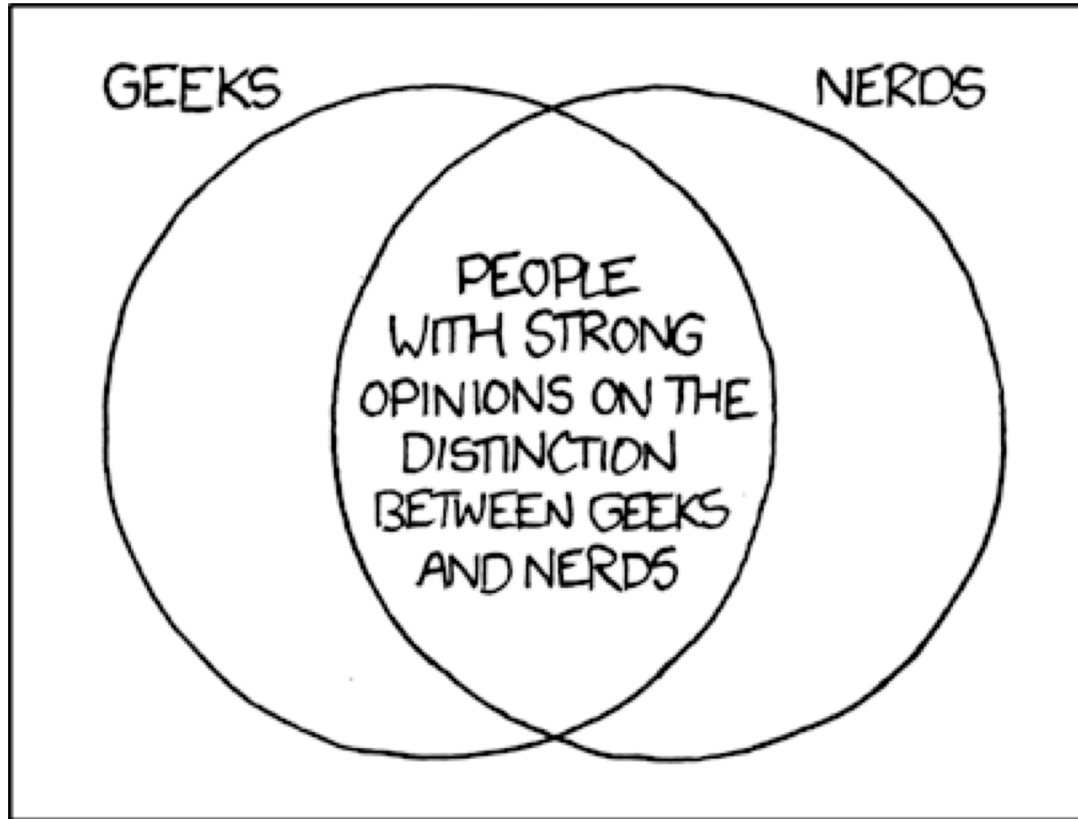


Sequence of  $k$  vertices connected by edges, first  $k-1$  are distinct



Questions?

# Formally define everything

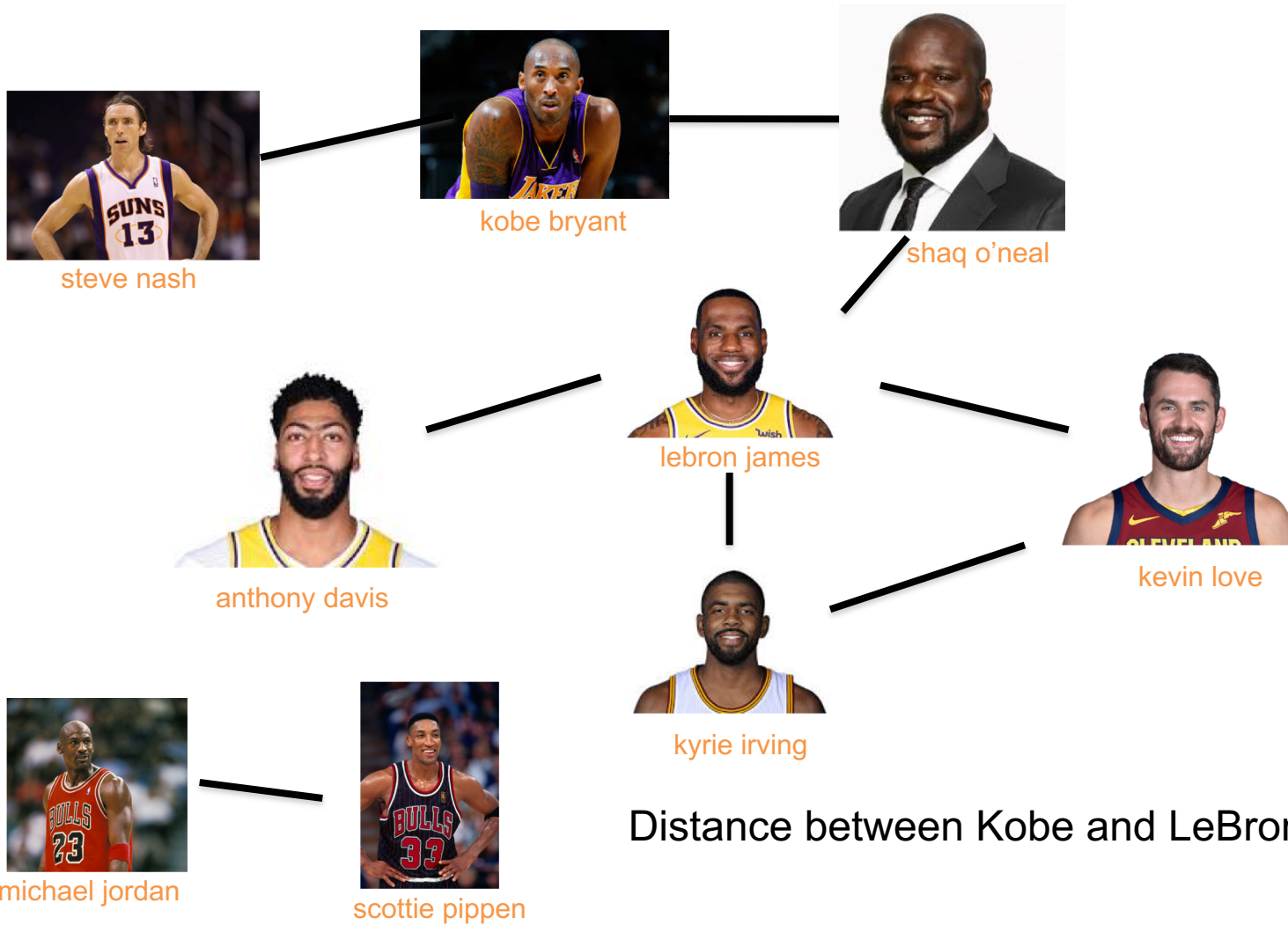


[http://imgs.xkcd.com/comics/geeks\\_and\\_nerds.png](http://imgs.xkcd.com/comics/geeks_and_nerds.png)



# Distance between $u$ and $v$

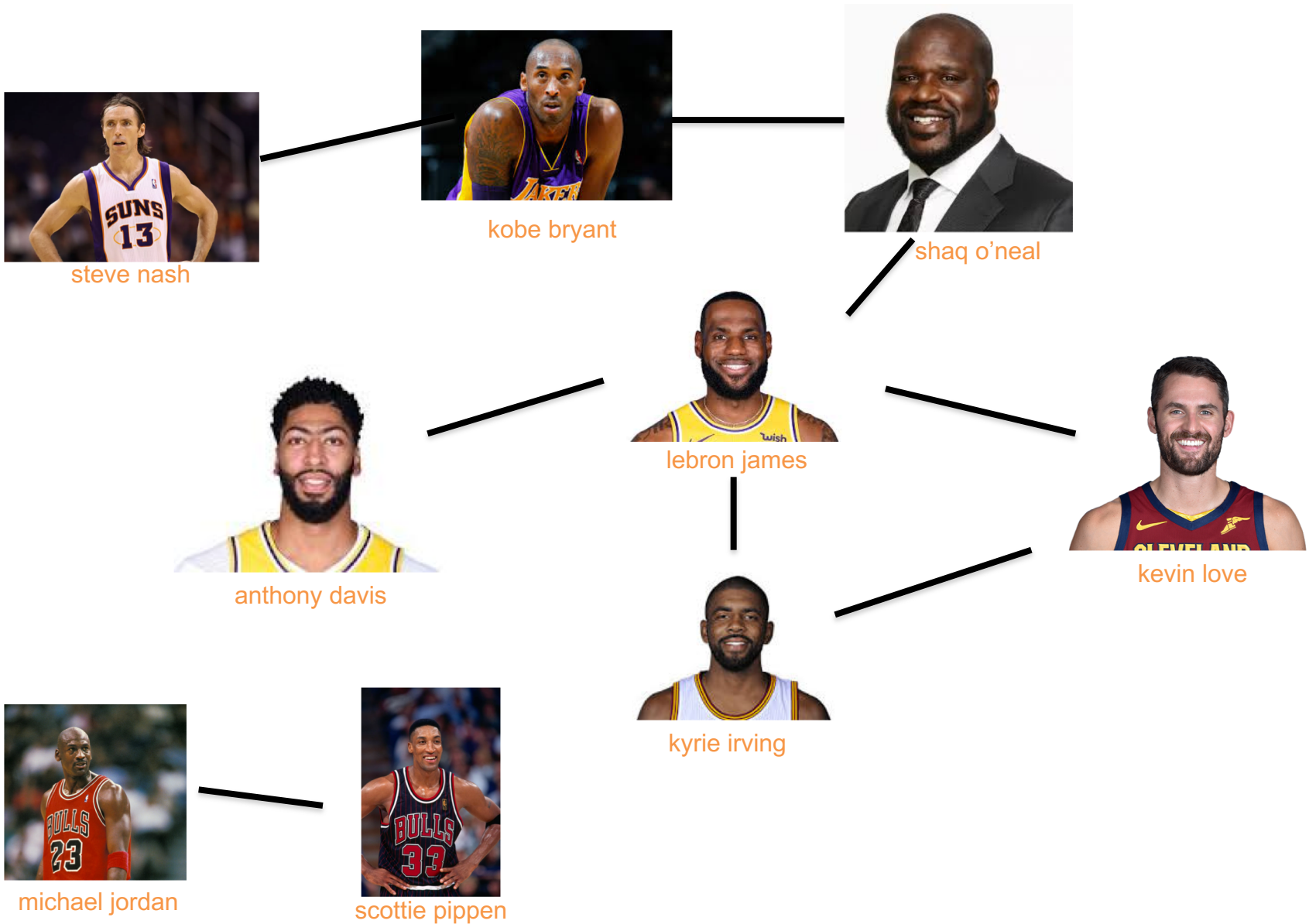
Length of the shortest length path between  $u$  and  $v$



Distance between Kobe and LeBron? 2

# Tree

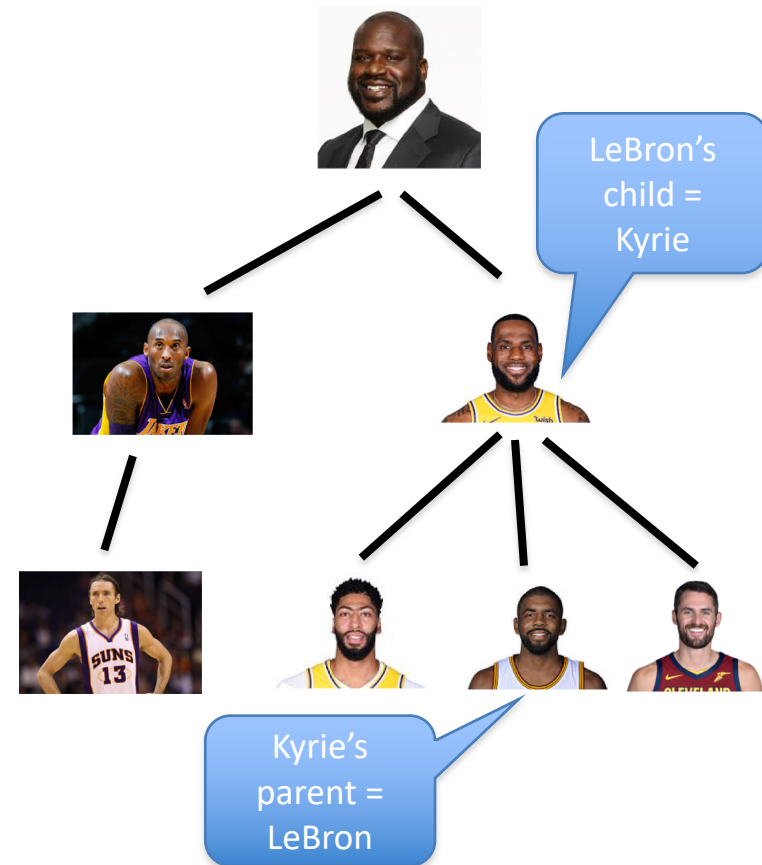
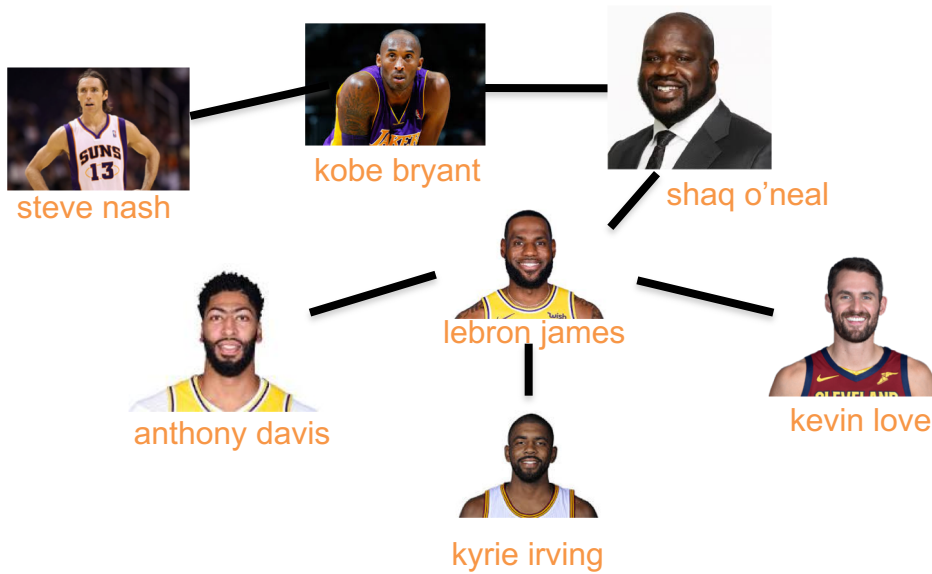
Connected undirected graph with no cycles



# Rooted Tree



# A rooted tree



Pick any vertex as root

Let the rest of the tree hang under “gravity”

# Every $n$ vertex tree has $n-1$ edges

## Trees

This page collects material from previous incarnations of CSE 331 on trees, especially the proof that trees with  $n$  nodes have exactly  $n - 1$  edges.

### Where does the textbook talk about this?

Section 3.1 in the textbook has the lowdown on trees.

### Fall 2018 material

Here is the lecture video:

CSE331 on 9/21/2018 (Fri)



Every  $n$  vertex tree has  $n-1$  edges

Let  $G$  be an undirected graph on  $n$  nodes

Then ANY two of the following implies the third:

$T$  is connected

$T$  has no cycles

$T$  has  $n-1$  edges

# Rest of Today's agenda

Algorithms for checking connectivity

# Checking by inspection



steve nash



kobe bryant



shaq o'neal



anthony davis



lebron james



kevin love



kyrie irving



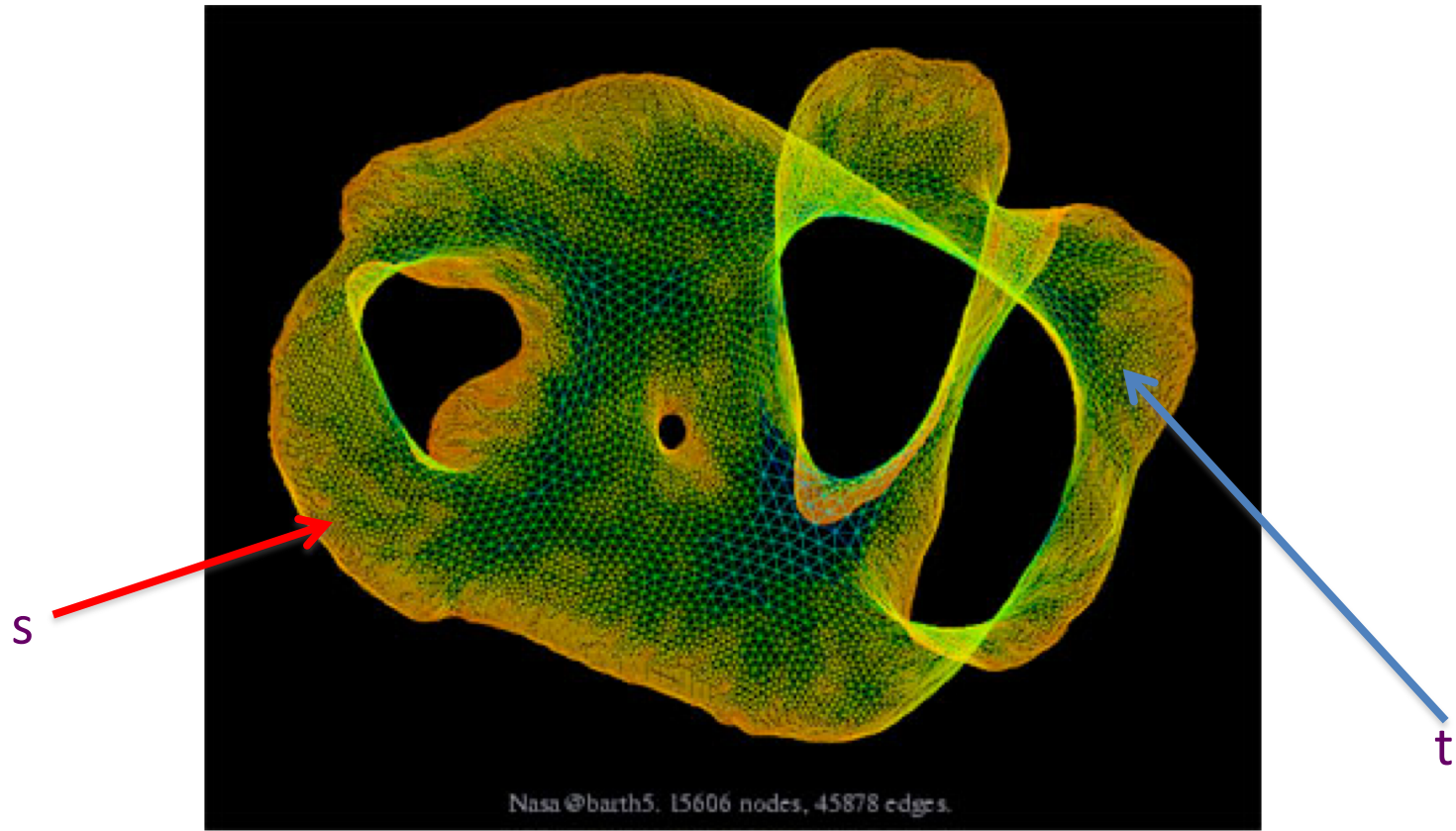
michael jordan



scottie pippen



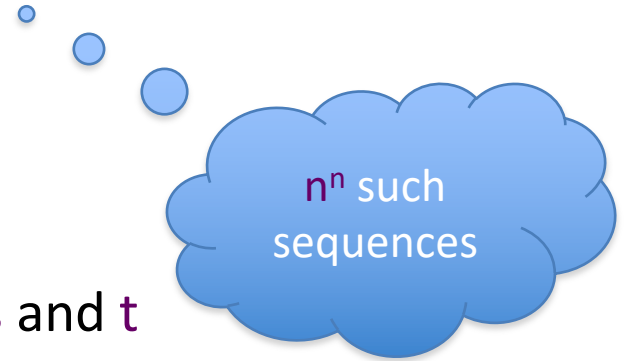
# What about large graphs?



Are  $s$  and  $t$  connected?

# Brute-force algorithm?

List all possible vertex sequences between  $s$  and  $t$



Check if any is a path between  $s$  and  $t$

# Algorithm motivation

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search ID: mbcn800

# Breadth First Search (BFS)

## BFS via examples

In which we derive the breadth first search (BFS) algorithm via a sequence of examples.

### Expected background

These notes assume that you are familiar with the following:

- Graphs and their representation. In particular,
  - Notion of connectivity of nodes and connected components of graphs
  - Adjacency list representation of graphs
  - Notation:
    - $G = (V, E)$
    - $n = |V|$  and  $m = |E|$
    - $CC(s)$  denotes the connected component of  $s$
- Trees and their basic properties

### The problem

In these notes we will solve the following problem:

# Connectivity Problem

*Input:* Graph  $G = (V, E)$  and  $s$  in  $V$

*Output:* All  $t$  connected to  $s$  in  $G$

# Breadth First Search (BFS)

Build layers of vertices connected to  $s$

$$L_0 = \{s\}$$

Assume  $L_0, \dots, L_j$  have been constructed

$L_{j+1}$  is the set of vertices not chosen yet but are connected by an edge to  $L_j$

Stop when new layer is empty