

## Knapsack Problem

**Input:** an integer bound  $W > 0$

a set of  $n$  items, each with an integer weight  $w_i > 0$

a value  $v_i > 0$  for each item  $i$

**Output:** a subset  $S$  of items that

$$\text{maximizes } \sum_{i \in S} v_i \quad \text{s.t. } \sum_{i \in S} w_i \leq W.$$

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- Motivation: you have budget  $W$ , and want to buy a subset of items of maximum total value

# DP for Knapsack Problem

- $opt[i, W']$ : the optimum value when budget is  $W'$  and items are  $\{1, 2, 3, \dots, i\}$ .
- If  $i = 0$ ,  $opt[i, W'] = 0$  for every  $W' = 0, 1, 2, \dots, W$ .

$$opt[i, W'] = \begin{cases} & i = 0 \\ & i > 0, w_i > W' \\ & i > 0, w_i \leq W' \end{cases}$$

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$$opt[i, W'] = \begin{cases} 0 & i = 0 \\ opt[i - 1, W'] & i > 0, w_i > W' \\ \max \left\{ \begin{array}{l} opt[i - 1, W'] \\ opt[i - 1, W' - w_i] + v_i \end{array} \right\} & i > 0, w_i \leq W' \end{cases}$$

## Exercise: Items with 3 Parameters

**Input:** integer bounds  $W > 0$ ,  $Z > 0$ ,  
a set of  $n$  items, each with an integer weight  $w_i > 0$   
a size  $z_i > 0$  for each item  $i$   
a value  $v_i > 0$  for each item  $i$

**Output:** a subset  $S$  of items that

$$\begin{aligned} &\text{maximizes } \sum_{i \in S} v_i && \text{s.t.} \\ &\sum_{i \in S} w_i \leq W \text{ and } \sum_{i \in S} z_i \leq Z \end{aligned}$$

# Outline

- 1 Weighted Interval Scheduling
- 2 Subset Sum Problem
- 3 Knapsack Problem
- 4 Longest Common Subsequence**
  - Longest Common Subsequence in Linear Space
- 5 Shortest Paths in Directed Acyclic Graphs
- 6 Matrix Chain Multiplication
- 7 Optimum Binary Search Tree
- 8 Summary



# Subsequence

- $A = bacdca$
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**Def.** Given two sequences  $A[1 .. n]$  and  $C[1 .. t]$  of letters,  $C$  is called a **subsequence** of  $A$  if there exists integers  $1 \leq i_1 < i_2 < i_3 < \dots < i_t \leq n$  such that  $A[i_j] = C[j]$  for every  $j = 1, 2, 3, \dots, t$ .

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- Exercise: how to check if sequence  $C$  is a subsequence of  $A$ ?

# Common subsequence

**Def.** Given two sequences  $A[1 .. n]$  and  $B[1 .. m]$  of letters,  $C$  is called a **common subsequence** of  $A$  and  $B$  if  $C$  is a subsequence of  $A$  and also a subsequence of  $B$ .

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- Example:  $A = \text{adecadf}$  and  $B = \text{caefcad}$

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- Common subsequence:  $C = \text{adcaf}$  ?

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- Example:  $A = \text{adecadf}$  and  $B = \text{caefcad}$
- Common subsequence:  $C = \text{adcaf}$  ?
- Common subsequence:  $C = \text{aead}$  ?



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- Example:  $A = \text{adecadf}$  and  $B = \text{caefcad}$
- Common subsequence:  $C = \text{adcaf}$  ?
- Common subsequence:  $C = \text{aead}$  ?
- Common subsequence:  $C = \text{acad}$  ?

# Edit distance with two operations (insertions and deletions)

**Def.** Given two sequences  $A[1 .. n]$  and  $B[1 .. m]$  of letters,  $d(A, B)$  is called a **edit distance with insert and delete operations** of  $A$  and  $B$  if  $d(A, B)$  is the minimum number of edit operations needed to transform  $A$  into  $B$ , where possible operations are:

- insert a character
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- Distance  $d(A, B) = 5$ : delete  $b$ , delete  $c$ , insert  $d$ , insert  $e$ , and insert  $f$ .

# Edit distance with three operations (insertions, deletions and replacing)

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- insert a character
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- Example:  $A = abc$  and  $B = adef$
  - Distance  $d(A, B) = 3$ : replace  $b$  to  $d$ , replace  $c$  to  $e$ , and insert character  $f$ .

## Quiz 5 on Ublearns

- Questions about subsequence, common subsequence, and edit distance
- Deadline: 25 Wed 2023, 11:59PM