Stokes Preconditioning on a GPU

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Collaborators

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 - Dept. of Geology and Geophysics, University of Minnesota
- Dr. David May, developer of BFBT (in PETSc)
 - Dept. of Earth Sciences, ETHZ
- Felipe Cruz, developer of FMM-GPU
 - Dept. of Applied Mathematics, University of Bristol
- Prof. Lorena Barba
 - Dept. of Mechanical Engineering, Boston University

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BFBT preconditions the Schur complement using

$$S_b^{-1} = L_p^{-1} G^T K G L_p^{-1} \tag{1}$$

where L_p is the Laplacian in the pressure space.



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Current Problems

The current BFBT code is limited by

- Bandwidth constraints
 - Sparse matrix-vector product
 - Achieves at most 10% of peak performance
- Synchronization
 - GMRES orthogonalization
 - Coarse problem
- Convergence
 - Viscosity variation
 - Mesh dependence



Alternative Proposal

Use a Boundary Element Method, for the Laplace solves in BFBT, accelerated by FMM.

Missing Pieces

BEM discretization and assembly

- Matrix-free operator application using the Fast Multipole Method
- Overcomes bandwidth limit, 480 GF on an NVIDIA 1060C GPU
- Overcomes coarse bottleneck by overlapping direct work

Solver for BEM system

- Same total work as FEM due to well-conditioned operator
- Possibility of multilevel preconditioner (even better)

Interpolation between FEM and BEM

- Boundary interpolation just averages
- Can again use FMM for interior

Direct Fast Method for Variable-Viscosity Stokes

- Complexity not currently precisely quantified
 - We would like a given number of flops/digit of accuracy
- Brute Force
 - Use BEM to compute layers between regions of constant viscosity
 - Better conditioned, but not direct
- Elegant method should be possible
 - The operator is pseudo-differential
 - "Kernel-independent" FMM exists



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Bandwidth

Small bandwidth to main memory can limit performance

- Sparse matrix-vector product
- Operator application
- AMG restriction and interpolation



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STREAM Benchmark

Simple benchmark program measuring sustainable memory bandwidth

- Protoypical operation is Triad (WAXPY): $\mathbf{w} = \mathbf{y} + \alpha \mathbf{x}$
- Measures the memory bandwidth bottleneck (much below peak)
- Datasets outstrip cache

Machine	Peak (MF/s)	Triad (MB/s)	MF/MW	Eq. MF/s
Matt's Laptop	1700	1122.4	12.1	93.5 (5.5%)
Intel Core2 Quad	38400	5312.0	57.8	442.7 (1.2%)
Tesla 1060C	984000	102000.0*	77.2	8500.0 (0.8%)

Table: Bandwidth limited machine performance

http://www.cs.virginia.edu/stream/



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Analysis of Sparse Matvec (SpMV)

Assumptions

- No cache misses
- No waits on memory references

Notation

- m Number of matrix rows
- nz Number of nonzero matrix elements
- V Number of vectors to multiply

We can look at bandwidth needed for peak performance

$$\left(8 + \frac{2}{V}\right) \frac{m}{nz} + \frac{6}{V} \text{ byte/flop}$$
 (2)

or achieveable performance given a bandwith BW

$$\frac{Vnz}{(8V+2)m+6nz}BW \text{ Mflop/s}$$
 (3)

Towards Realistic Performance Bounds for Implicit CFD Codes, Gropp, Kaushik, Keyes, and Smith.

Improving Serial Performance

For a single matvec with 3D FD Poisson, Matt's laptop can achieve at most

$$\frac{1}{(8+2)\frac{1}{7}+6} \text{ bytes/flop(1122.4 MB/s)} = 151 \text{ MFlops/s}, \qquad (4)$$

which is a dismal 8.8% of peak.

Can improve performance by

- Blocking
- Multiple vectors

but operation issue limitations take over.



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Better approaches:

- Unassembled operator application (Spectral elements, FMM)
 - N data, N² computation
- Nonlinear evaluation (Picard, FAS, Exact Polynomial Solvers)
 - N data, N^k computation



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Synchronization

Synchronization penalties can come from

- Reductions
 - GMRES orthogonalization
 - More than 20% penalty for PFLOTRAN on Cray XT5
- Small subproblems
 - Multigrid coarse problem
 - Lower levels of Fast Multipole Method tree

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Convergence

Convergence of the BFBT solve depends on

- Viscosity constrast (slightly)
- Viscosity topology
- Mesh

Convergence of the AMG Poisson solve depends on

Mesh



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- Solver for BEM system
- Interpolation between FEM and BEM

The Poisson problem

$$\Delta u(\mathbf{x}) = f(\mathbf{x}) \quad \text{on } \Omega$$
 (5)

$$u(\mathbf{x})|_{\partial\Omega} = g(\mathbf{x})$$
 (6)



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The Poisson problem (Boundary Integral Equation formulation)

$$C(\mathbf{x})u(\mathbf{x}) = \int_{\partial\Omega} F(\mathbf{x}, \mathbf{y})g(\mathbf{y}) - G(\mathbf{x}, \mathbf{y}) \frac{\partial u(\mathbf{y})}{\partial n} dS(\mathbf{y})$$
(5)

$$G(\mathbf{x}, \mathbf{y}) = -\frac{1}{2\pi} \log r$$
(6)

$$F(\mathbf{x}, \mathbf{y}) = \frac{1}{2\pi r} \frac{\partial r}{\partial n}$$
(7)

$$G(\mathbf{x}, \mathbf{y}) = -\frac{1}{2\pi} \log r \tag{6}$$

$$F(\mathbf{x}, \mathbf{y}) = \frac{1}{2\pi r} \frac{\partial r}{\partial n} \tag{7}$$



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Restricting to the boundary, we see that

$$\frac{1}{2}g(\mathbf{x}) = \int_{\partial\Omega} F(\mathbf{x}, \mathbf{y})g(\mathbf{y}) - G(\mathbf{x}, \mathbf{y}) \frac{\partial u(\mathbf{y})}{\partial n} dS(\mathbf{y})$$
 (5)



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Discretizing, we have

$$-Gq = \left(\frac{1}{2}I - F\right)g\tag{5}$$



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Now we can evaluate u in the interior

$$u(\mathbf{x}) = \int_{\partial\Omega} F(\mathbf{x}, \mathbf{y}) g(\mathbf{y}) - G(\mathbf{x}, \mathbf{y}) \frac{\partial u(\mathbf{y})}{\partial n} dS(\mathbf{y})$$
 (5)



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Or in discrete form

$$u = Fg - Gq \tag{5}$$



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The sources in the interior may be added in using superposition

$$\frac{1}{2}g(\mathbf{x}) = \int_{\partial\Omega} F(\mathbf{x}, \mathbf{y})g(\mathbf{y}) - G(\mathbf{x}, \mathbf{y}) \left(\frac{\partial u(\mathbf{y})}{\partial n} - f\right) dS(\mathbf{y})$$
 (5)



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BEM Solver

The solve has two pieces:

- Operator application
 - Boundary solve
 - Interior evaluation
 - Accomplished using the Fast Multipole Method
- Iterative solver
 - Usually GMRES
 - We use PETSc



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Operator Application

Using the Fast Multiple Method, the Green's functions (*F* and *G*) can be applied:

- in $\mathcal{O}(N)$ time
- using small memory bandwidth
- in the interior and on the boundary
- with much higher serial and parallel performance



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Fast Multipole Method

FMM accelerates the calculation of the function:

$$\Phi(x_i) = \sum_j K(x_i, x_j) q(x_j)$$
 (6)

- Accelerates $\mathcal{O}(N^2)$ to $\mathcal{O}(N)$ time
- The kernel $K(x_i, x_i)$ must decay quickly from (x_i, x_i)
 - Can be singular on the diagonal (Calderón-Zygmund operator)
- Discovered by Leslie Greengard and Vladimir Rohklin in 1987
- Very similar to recent wavelet techniques



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Fast Multipole Method

FMM accelerates the calculation of the function:

$$\Phi(x_i) = \sum_j \frac{q_j}{|x_i - x_j|} \tag{6}$$

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