

Algo idea: Replace each edge $e \in E$ by a path of length l_e . (new paths do not share edges/vertices)

\Rightarrow Run HW3 Q3 algo.

Claim: a shortest s-t path in G \Leftrightarrow equiv. a shortest s-t path in G'

Correctness:

\downarrow claim + correctness of HW3 Q3 algo

Runtime:

$$\begin{array}{l} O(|V'| + |E'|) \\ O(m' + n') \end{array} \quad \begin{array}{l} |V'| = n' \\ |E'| = m' \end{array}$$

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$$l_{\max} = \max l_e$$

$$m' \leq l_{\max} \cdot m$$

$$n' \leq l_{\max} \cdot n$$

$$G: O(|V| + |E|)$$

$$O(m+n)$$

$$\begin{array}{l} |V| = m \\ |E| = n \end{array}$$

runtime: $O(l_{\max} \cdot (m+n))$

if $l_{\max} = O(1)$, then fine!

if $l_{\max} = n^{100} \rightarrow$ runtime: $O(n^{100} (m+n))$

Aside/Recap:

RAM model of computation

unit of space is a register.

if you have n items, each register is $\log(m)$ bits.

adj. list representation: $O(n+m)$ registers.

Q. How many registers do you need to represent $l_e \leq n^{100}$

↳ How bits in each register?

↳ $\approx 100 \log n$ bits

≈ 100 registers
= $O(1)$ registers.

Input (# registers): $O(n+m)$

Runtime: $O(n^{100}(n+m))$

IDEALLY: want to have $O(n+m)$ runtime
for the case $l_{\max} \leq n^{O(1)}$