

Lecture 6

CSE 331

Register your project groups

Deadline: Friday, March 3, 11:59pm

CSE 331 Project

Spring 2023

Details and motivations for the project.

Project Overview

Group signup form

Motivation

CSE 331 is primarily concerned with the technical aspects of algorithms: how to design them and then how to analyze their correctness and runtime. However, algorithms are pervasive in our world and are common place in many aspects of society. The main aim of the project is to have you explore in some depth some of the social implications of algorithms.

Just to give some examples for such implications:













Your UB email: [XXX@buffalo.edu](#)

Your UBIT ID is **XXX**

















NOT [XXX @buffalo.edu](#)

NOT your person number

HW 1 gets released Today

Week 2 Mon, Feb 6	Perfect Matchings     F22  F21  S21  S20 x ²	[KT, Sec 1.1] (HW 0 in) Week 2 recitation notes
Wed, Feb 8	Algorithms for stable matching problem     F22  F21  S21  S20 x ²	[KT, Sec 1.1]
Fri, Feb 10	Gale Shapley algorithm  F22  F21  S21  S20 x ²	[KT, Sec 1.1] (HW 1 out) <i>Reading Assignment:</i> Pigeonhole principle <i>Reading Assignment:</i> Asymptotic notation care package

Reading Assignments

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Stable Marriage problem

Set of men M and women W

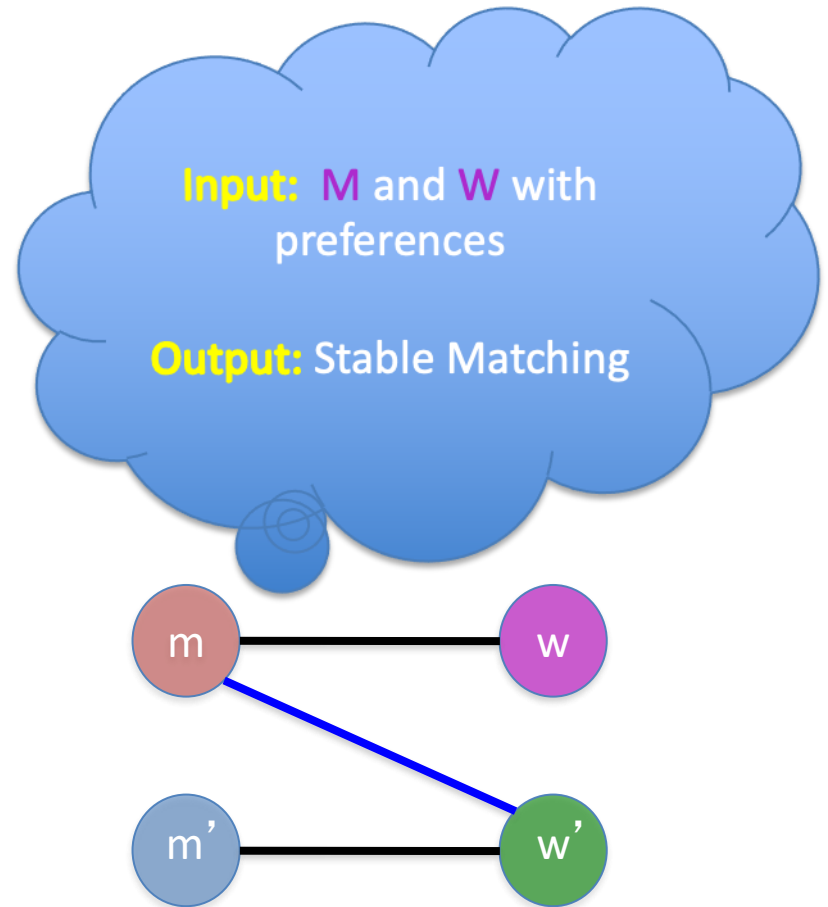
Preferences (ranking of potential spouses)

Matching (no polyandry/gamy in $M \times W$)

Perfect Matching (everyone gets married)

Instability

Stable matching = perfect matching + no instability



Two Questions

Does a stable marriage always exist?

If one exists, how quickly can we compute one?

The naïve algorithm

Go through all possible perfect matchings S

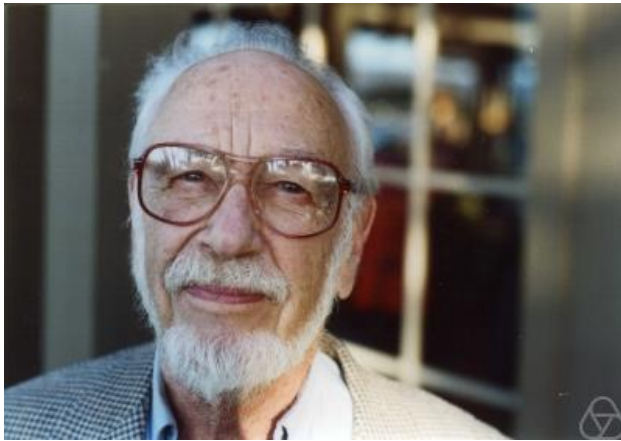
If S is a stable matching

then Stop



Else move to the next perfect matching

Gale-Shapley Algorithm



David Gale



Lloyd Shapley

$O(n^2)$ algorithm

Rest of today's agenda

GS algorithm

Run of GS algorithm on an instance

Prove correctness of the GS algorithm

Gale-Shapley Algorithm

Initially all men and women are **free**

While there exists a free woman who can propose

Let w be such a woman and m be the best man she has not proposed to

w proposes to m

If m is free

(m,w) get **engaged**

Else (m,w') are engaged

If m prefers w' to w

w remains **free**

Else

(m,w) get **engaged** and w' is **free**

Output the engaged pairs as the final output

Rest of today's agenda

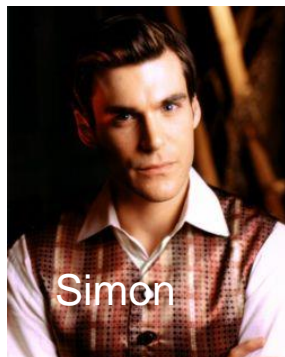
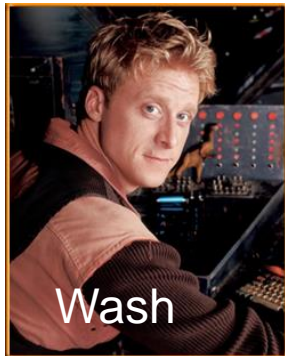
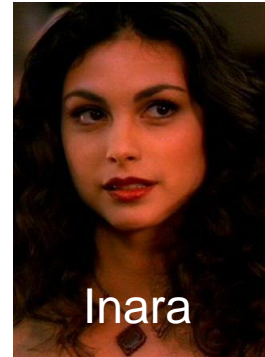
GS algorithm

Run of GS algorithm on an instance

Prove correctness of the GS algorithm

Back to working on paper...

Preferences



GS algorithm: Firefly Edition

