Lecture 8

CSE 331

Gale-Shapley Algorithm

Intially all men and women are free

While there exists a free woman who can propose

Let w be such a woman and m be the best man she has not proposed to w proposes to m If m is free (m,w) get engaged Else (m,w') are engaged If m prefers w' to w w remains free Else (m,w) get engaged and w' is free

Output the engaged pairs as the final output

The Lemmas

Lemma 1: The GS algorithm has at most n² iterations

Lemma 2: S is a perfect matching

Lemma 3: S has no instability

Proof Details of Lemma 1

Gale Shapley algorithm terminates

This page collects material from Fall 17 incarnation of CSE 331, where we proof details for the claim that the Gale-Shapley algorithm terminates in $O(n^2)$ iterations.

Where does the textbook talk about this?

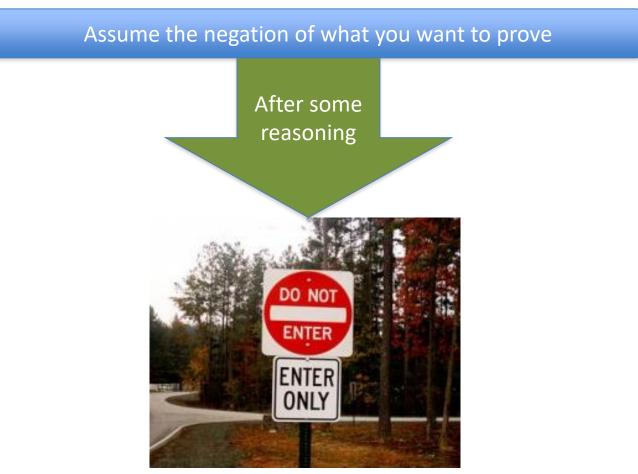
Section 1.1 in the textbook has the argument (though not in as much detail as below).

Fall 2017 material

Here is the lecture video (it starts from the part where we d the proof details):



Proof by contradiction



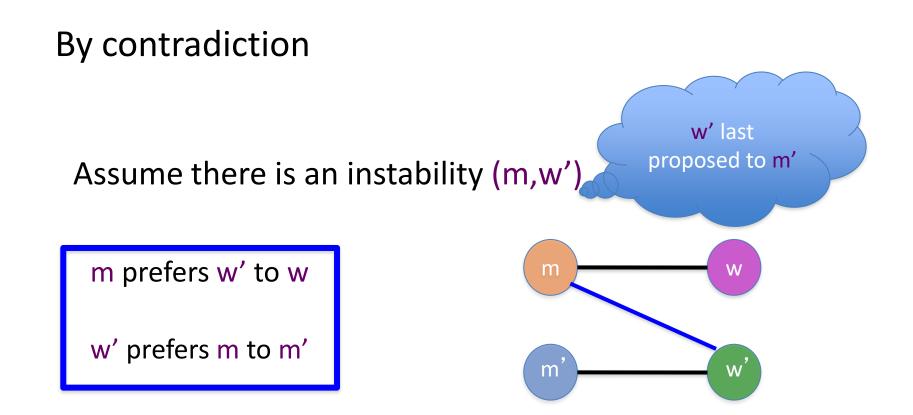
Source: 4simpsons.wordpress.com

Two obervations

Obs 1: Once m is engaged he keeps getting engaged to "better" women

Obs 2: If w proposes to m' first and then to m (or never proposes to m) then she prefers m' to m

Proof of Lemma 3



Contradiction by Case Analysis

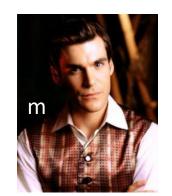
Depending on whether w' had proposed to m or not

Case 1: w' never proposed to m

w' prefers m' to m_

By Obs 2

Assumed w' prefers m to m'







Source: 4simpsons.wordpress.com



Case 2: w' had proposed to m

Case 2.1: m had accepted w' proposal

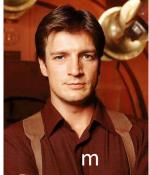
m is finally engaged to w

Thus, m prefers w to w'

By Obs 1



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Case 2.2: m had rejected w' proposal

m was engaged to w'' (prefers w'' to w') By Algo def

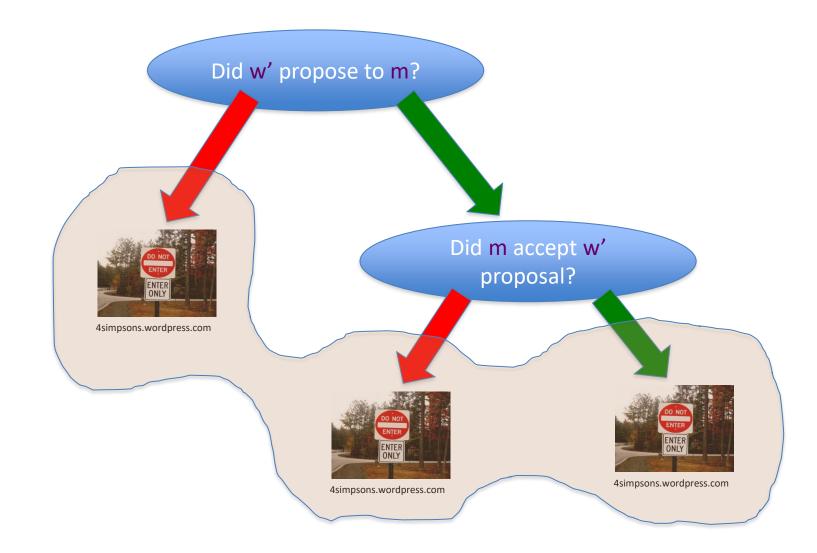
m is finally engaged to w (prefers w to w") By Obs 1

m prefers w to w'



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Overall structure of case analysis



Questions?



Fairness of the GS algorithm

Different executions of the GS algorithm