

Shortest path problem

Input :

Directed graph $G = (V, E)$

Start vertex s

"lengths" $l_e \geq 0 \quad \forall e \in E$
 $\uparrow$ integers

Output :

A shortest s - t path $P \quad \forall t \in V$

$\uparrow$
w.r.t. the
length of P

$$l(P) = \sum_{e \in P} l_e$$

Simpler Version :

Output $d(t) \quad \forall t \in V$

$d(t)$: distance of t from s
 $\uparrow$
in terms of
lengths

$\Rightarrow$ length of the shortest s - t path

Special Case :

$l_e = 1 \quad \forall e \in E \equiv \text{HW3 Q3}$

Q! $l_e = L$ (for some $L \geq 1$) (run BFS from s
and layer # $t = d(t)$)
e.g. $L = 100$

$l_e = 100 \quad \forall e \in E$

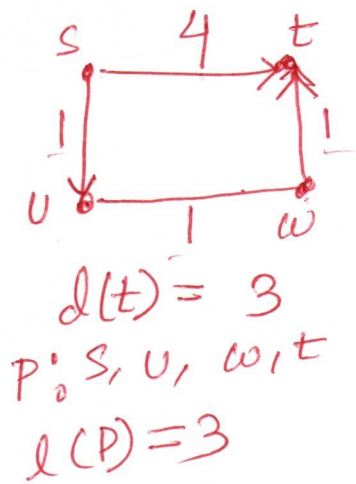
A: Reduce this case to HW3 Q3

How! Ignore L , assume $l_e = 1 \quad \forall e \in E$

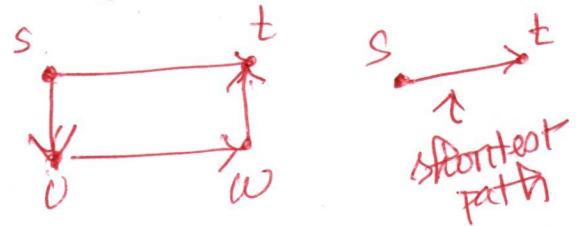
7/18 General case 8 $l_e \geq 0, \forall e \in E$ [l_e can have different lengths]

Algo Idea 0 Reduce the general case to $l'_e = 1$ and run HW3 Q3 algo.

Idea 1 ignore $l_e \Rightarrow$ assume $l_e = 1 \forall e \in E$ and run HW3 Q3 algo.

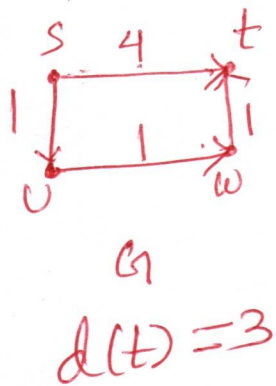


Idea 1
 ignore l_e
 $\Rightarrow$ assume $l_e = 1 \forall e \in E$

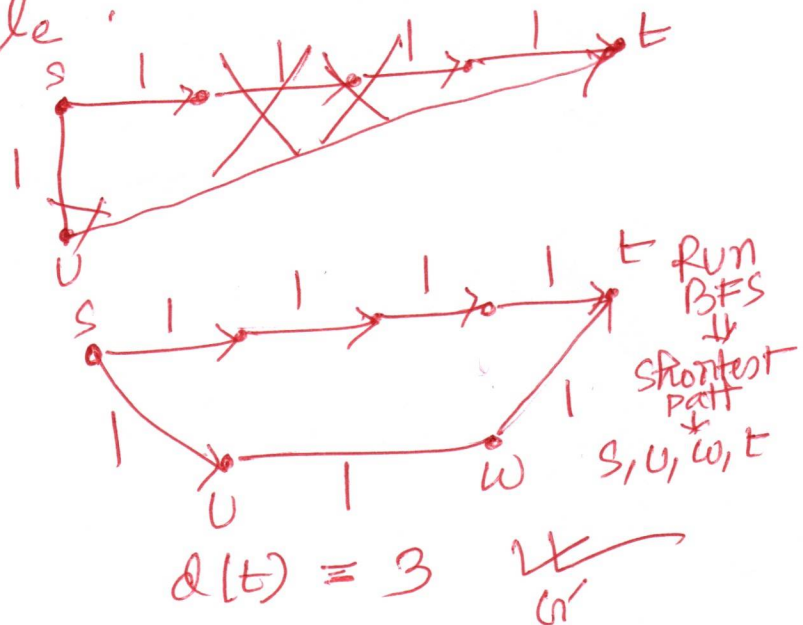


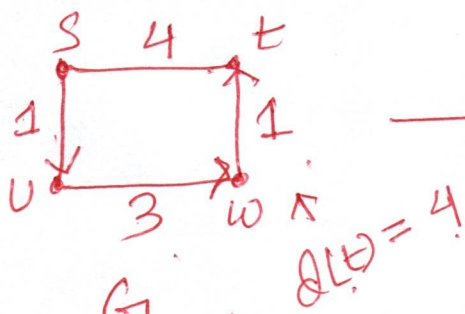
* BFS works only if the edge lengths are the same.

Idea 2 Replace an edge e of length $l_e \geq 1$ by a path of length l_e .

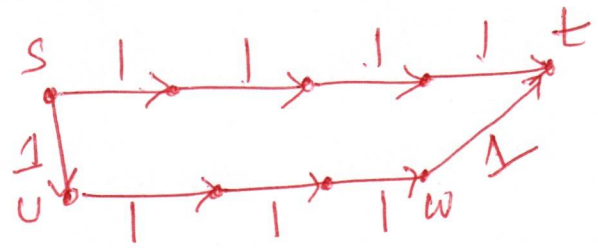


Idea 2





idea 2



$\Rightarrow$ Run BFS
 $\Rightarrow$ shortest path s, t
 $d(t) = 4$

$G = (V, E)$
 Runtime: $O(m+n)$
 $n = |V|$ $m = |E|$

Runtime Analysis of $G' = (V', E')$

Connectness of idea 2

Claim: Shortest path in G .

$\Rightarrow$ equivalent to shortest path in G' .

proof of claim + connectness of HW3 Q3 algo.

$G' = (V', E')$ $|V'| = n'$
 $|E'| = m'$

$l_{\max} = \max_{e \in E} l_e$

$n' \leq l_{\max}(n+m)$

Runtime: $O(n' + m')$

$m' \leq l_{\max}(m+n)$

$= O(l_{\max}(n+m))$

$\Rightarrow O(5(m+n))$

$\Rightarrow O(n^{100}(m+n))$

The graph G is in adj. list format

Space : $O(m+n)$

RAM model : Unit of space is a register.

If you have n input items

$\Rightarrow$ how many bits is a register?

$\Rightarrow \log_2 n$ bits register

w bits

$\Rightarrow 2^w$ items

$2^w \geq n$

$w \geq \log_2 n$

$\Rightarrow$ How many ~~100~~ bit registers
for n^{100} ?

Ans

$2^w \geq n^{100}$

$w \geq 100 (\log n)$

$\Rightarrow 100$ registers

$\Rightarrow O(1)$!!

$\Rightarrow$ runtime : $O(n^{100} (m+n))$