

# Lecture 17

CSE 331

# Runtime analysis of Greedy Algo.

$R$ : set of requests

$O(1)$

Set  $S$  to be the empty set

Repeated at most  $n$  times

While  $R$  is not empty

$O(n)$

$O(1)$

Choose  $i$  in  $R$  with the earliest finish time

Add  $i$  to  $S$

Remove all requests that conflict with  $i$  from  $R$

$O(n)$

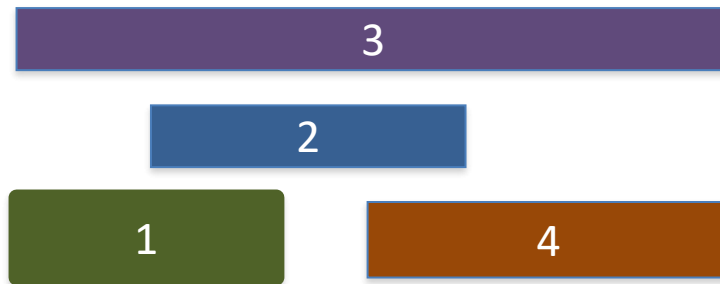
Return  $S^* = S$

$O(n)$

Overall:  
 $O(n) + n \cdot O(n) = O(n^2)$

# Algorithm implementation

Go through the intervals in order of their finish time



Check if  $s[i] < f(1)$

With 1:

In general, if  $j$ th interval is the last one chosen

Pick smallest  $i > j$  such that  $s[i] \geq f(j)$  . . .

$O(n \log n)$  run  
time

# The final algo

$O(n \log n)$  time sort intervals such that  $f(i) \leq f(i+1)$

$O(n)$  time build array  $s[1..n]$  s.t.  $s[i]$  = start time for  $i$

Add 1 to  $A$  and set  $f = f(1)$

For  $i = 2 \dots n$

    If  $s[i] \geq f$

        Add  $i$  to  $A$

        Set  $f = f(i)$

Return  $A^* = A$

# Reading Assignment

Sec 4.1 of [KT]