

Lecture 19

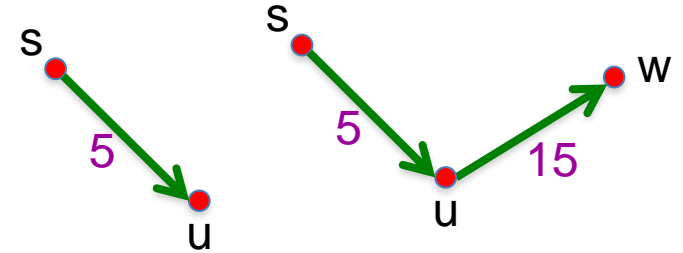
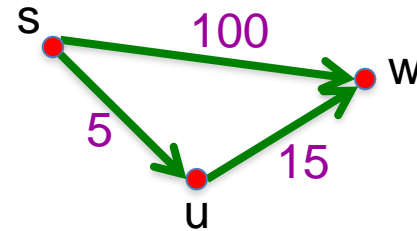
CSE 331

Shortest Path problem

Input: *Directed* graph $G=(V,E)$

Edge lengths, l_e for e in E

“start” vertex s in V

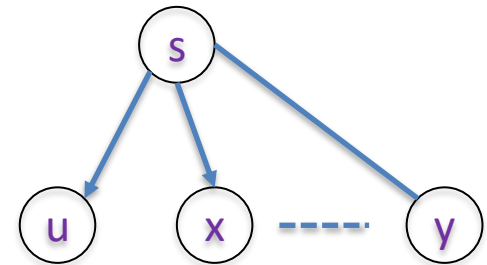


Output: Length of shortest paths from s to all nodes in V

Towards Dijkstra's algo: part one

Determine $d(t)$ one by one

$$d(s) = 0$$



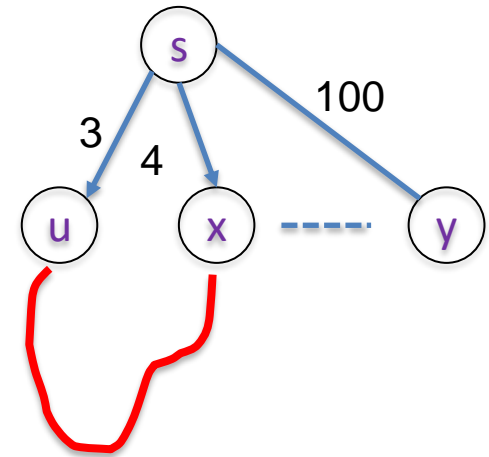
Towards Dijkstra's algo: part two

Determine $d(t)$ one by one

Let u be a neighbor of s with smallest $l_{(s,u)}$

$$d(u) = l_{(s,u)}$$

Not making any claim
on other vertices



Length of  is ≥ 0

Towards Dijkstra's algo: part three

Determine $d(t)$ one by one

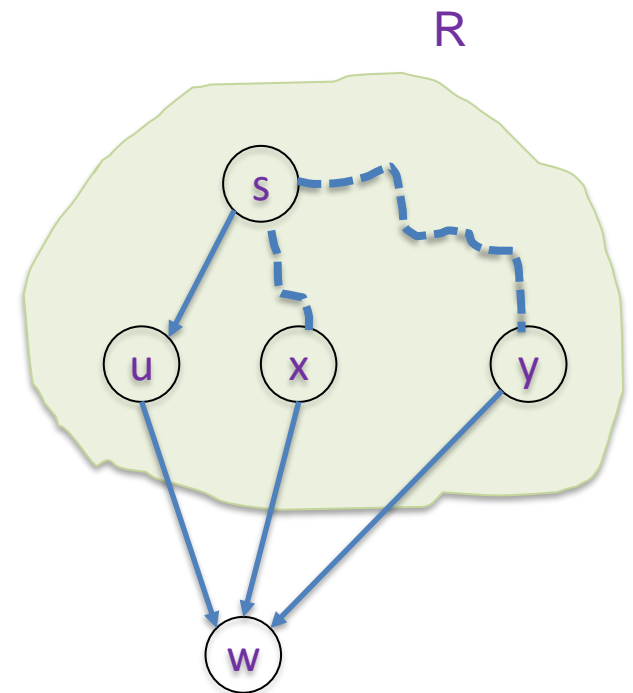
Assume we know $d(v)$ for every v in R

Compute an upper bound $d'(w)$ for every w not in R

$$d(w) \leq d(u) + l_{(u,w)}$$

$$d(w) \leq d(x) + l_{(x,w)}$$

$$d(w) \leq d(y) + l_{(y,w)}$$



$$d'(w) = \min_{e=(u,w) \in E, u \in R} d(u) + l_e$$

Reading Assignment

Sec 4.4 of [KT]