

*Sample exam for CSE396, Spring 2021. Content similar but your format will differ.*

Closed books and laptops, one notes sheet allowed, closed neighbors, 75 minutes. Do ALL FOUR problems **on these exam sheets**. Please *show all your work*—this may help for partial credit. The exam totals 80 pts., subdivided as shown.

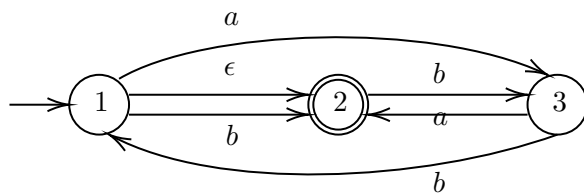
**Notation:** All problems on this exam use alphabet  $\Sigma = \{a, b\}$ .  $\#c(x)$  stands for the number of occurrences of the character  $c$  in the string  $x$ .

**(1) (20 pts.)**

Let  $A = \{x \in \{a, b\}^* : bb \text{ is a substring of } x \text{ and } \#a(x) \text{ is even}\}$ . Design a deterministic finite automaton (DFA)  $M$  such that  $L(M) = A$ . A node-arc diagram that shows the start and final states clearly is good enough—you need not write out tables or “ $M = (Q, \Sigma, \delta, s, F) \dots$ ” etc. For full credit, you must either have comments explaining a design strategy that makes the correctness of your  $M$  clear, or you must use theorems to build  $M$  from smaller machines. Trial-and-error may take too long.

(2) (15 + 12 = 27 pts.)

Let  $N$  be the NFA defined by  $N = (Q, \Sigma, \delta, s, F)$  with  $Q = \{1, 2, 3\}$ ,  $\Sigma = \{a, b\}$ ,  $s = 1$ ,  $F = \{2\}$ , and  $\delta$  given by the arcs  $(1, b, 2)$ ,  $(1, \epsilon, 2)$ ,  $(1, a, 3)$ ,  $(2, b, 3)$ ,  $(3, b, 1)$ , and  $(3, a, 2)$  as shown in the following node-arc diagram:



- (a) Calculate a DFA  $M$  such that  $L(M) = L(N)$ .
- (b) Calculate a regular expression  $r$  such that  $L(N) = L(r)$ .

**(3) (5 × 3 = 15 pts.)** *True/False.*

Please write out the words **true** and/or **false** in full. No justifications are needed.

- (a) If  $A^* = A$ , then the language  $A$  includes the empty string.
- (b) If  $A$  and  $B$  are regular languages recognized by 3-state DFAs, then  $A \cap B$  can be recognized by a 6-state DFA.
- (c) The empty relation on a nonempty set is transitive.
- (d) The intersection of two non-regular languages is always non-regular.
- (e) If there is a string  $w$  such that no string  $x$  in a regular language  $A$  has  $w$  as a substring, then every DFA  $M$  such that  $L(M) = A$  has a dead state.

(a) \_\_\_\_\_ (b) \_\_\_\_\_ (c) \_\_\_\_\_ (d) \_\_\_\_\_ (e) \_\_\_\_\_

**(4) (18 pts.)**

Define  $L = \{x \in \{a, b\}^* : |x| \text{ is even and there is a } b \text{ in the left half of } x\}$ . For instance,  $ba$ ,  $aabaaa$ , and  $baaaaa$  belong to  $L$ , but  $\epsilon$ ,  $ab$ , and  $baaaa$  do not—the last fails because its length is odd. Prove using the Myhill-Nerode technique that  $L$  is not a regular language.