(1) Let $L_{3}$ be the language of binary strings that represent positive multiples of 3. Define $L$ to be the language of binary strings that do not have a substring in $L_{3}$. That is, $L$ is the complement of $(0+1)^{*} \cdot L_{3} \cdot(0+1)^{*}$. Follow these steps to characterize the language $L$.
(a) Design a DFA $M_{3}$ such that $L(M)=L_{3}$. It is possible that you might find this just in the course of innocently reading online notes, and that's OK, actually. (Just 3 pts.)
(b) Now make some small additions to your DFA $M_{3}$ to create an NFA $N_{3}$ such that $L\left(N_{3}\right)=$ $(0+1)^{*} \cdot L_{3} \cdot(0+1)^{*}$. (3 pts.)
(c) Convert $N_{3}$ into an equivalent DFA $M^{\prime}$ such that $L\left(M^{\prime}\right)=(0+1)^{*} \cdot L_{3} \cdot(0+1)^{*}$. (12 pts.)
(d) Complement the final states of $M^{\prime}$ to get the needed $M$. ( 3 pts.)
(e) The string 11 sends $M$ to a dead state wherever you start from because it is 3 in binary, so it makes a substring that belongs to $L_{3}$. Find a similar "dead substring" that does not have two consecutive 1 s in it. ( 3 pts ., for 24 total)
(2) Design both nondeterministic finite automata $N_{a}, N_{b}, N_{c}$ and regular expressions $r_{a}, r_{b}, r_{c}$ that denote the following three languages described in prose. (It is OK for your NFAs to have $\epsilon$-arcs, and it is fine if one or more are DFAs since a DFA "Is-A" NFA.) All use the alphabet $\Sigma=\{0,1\}$. $(3 \times(6+6)=36$ points total $)$
(a) $L_{a}=$ the set of binary strings in which the substring 10 occurs an odd number of times.
(b) $L_{b}=$ the set of binary strings of the form $x=y 00 z$ where $|z|$ is odd.
(c) $L_{c}=$ the set of binary strings having an occurrence of the substring 10 that do not have an occurrence of 11 after it.
(3) Convert the folowing NFA $N$ into an equivalent DFA $M$ (18 pts.):


Also answer (for $3 \times 3=9$ more pts., making 87 total on the set):
(a) Is there a string $v$ that $N$ cannot process from start to any state? Again give a shortest such $v$.
(b) Is there a string $w$ such that no matter what string $y$ follows it, the string $w y$ is accepted? Again use your $M$.
(c) Stronger than (a) and counter to (b), is there a string $z$ that $N$ cannot process at all, not from any of its states $p$ to any state $q$ ? Again use your $M$ to explain your answer.

