• If the input tapes of both machine are right-only as well as read-only, then there is no problem: the output \( y = f(x) \) of \( M \) is streamed to \( M' \) computing \( g(y) = z \) and never has to be written down.

• If each machine is allowed \( r(n) \) left-to-right streaming passes over its input and \( y \) is a stream, then the tandem can operate with \( r(n)^2 \) passes on \( x \).

• But if \( M' \) can demand to back up to a previous input bit \( y_{i-1} \) at any time, then we need to allow \( M \) to be restarted arbitrarily many times. This can be implemented by storing the current demand-bit \( i \) on another log-sized tape.
• All the NP-completeness results we've shown have been valid under $\leq_m^{\log}$. 
• **GAP** is complete for NL under $\leq_m^{\log}$. 
• The language **CVP** of the **Circuit Value Problem**: given a Boolean circuit $C_n$ and an input $x \in \{0, 1\}^n$, is $C_n(x) = 1$? is complete for P under $\leq_m^{\log}$. 
• The language **TQBF** of true **quantified Boolean formulas** is complete for **PSPACE** under $\leq_m^{\log}$.