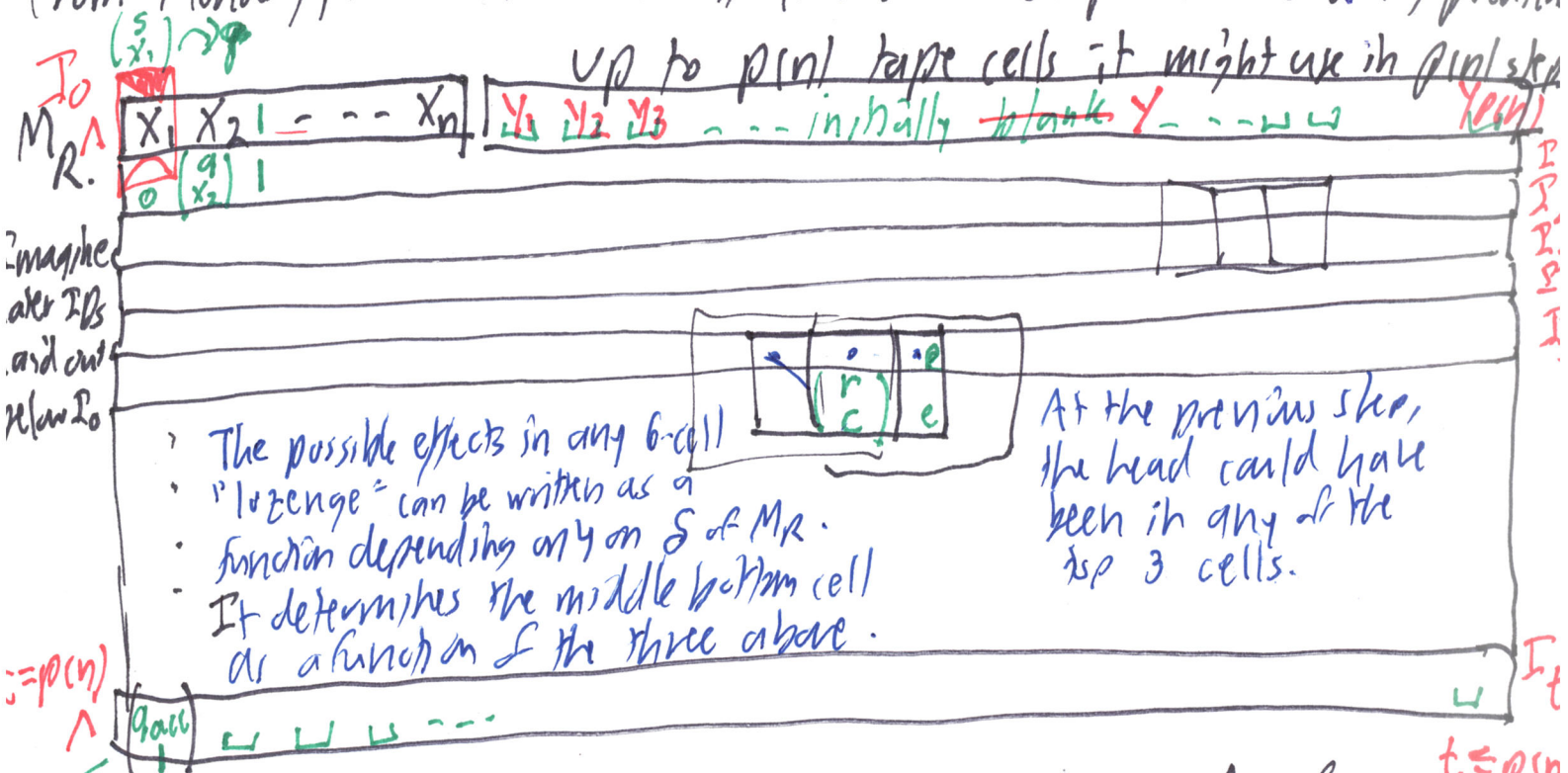


Theorem: [Steve Cook 1970-71, Leonid Levin 1970-1973]

SAT is NP-complete, i.e. $SAT \in NP$ (already seen) and for all $A \in NP$, $A \leq_m^P SAT$. [We will prove a reduction to a subclass of SAT.]

Proof: By $A \in NP$, there is a polynomial $p(n)$ and a single-tape TM $M_R(x, \gamma)$ that computes a witness predicate $R(x, \gamma)$ with $|\gamma| \leq p(|x|)$.
From Monday, we saw that $R(x, \gamma)$ can be the computation checking predicate



Each lozenge has the same function Δ of six characters (3 inputs, 1 output) over $\Sigma \cup \{ \square \}$.

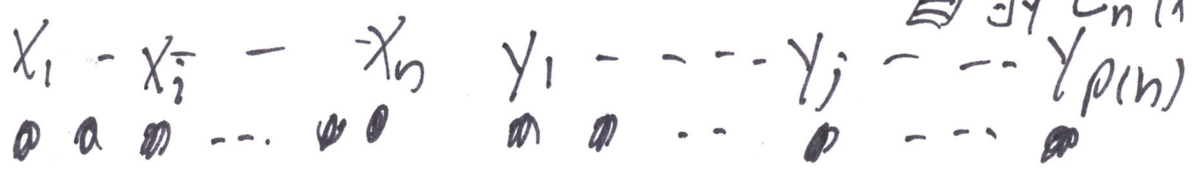
We can re-code those over binary and write the finite Δ function as a Boolean function, using NAND gates only. $\therefore M_R$ can be simulated by $p(n) \times p(n)$ sized circuits. [Can get $O(p(n))$ by $p(n)$.]

output write w_0

$x \in A \iff$ the witness is 1.

observe M_R paths before $p(n)$ steps.

We have abstracted this to a circuit C_n st. $x \in A \Leftrightarrow \exists y \bigwedge R(x, y)$
 $\Leftrightarrow \exists y C_n(x, y) = 1$



We can stamp out =

(u and/or v can be wires from inputs x_i or guesses y_j)



Gate g functions correctly on bits u, v giving output w if and only if ϕ_g is made true, where

$$\phi_g = (u \vee w) \wedge (v \vee w) \wedge (\bar{u} \vee \bar{v} \vee \bar{w})$$

C_n in $O(p(n)^2)$ time.

Final ϕ_x is:

Only part that depends on not just $n = |X|$

$$\left(\bigwedge_{\text{NAND gates } g \text{ in } C} \phi_g \right) \wedge (W_0)$$

Singleton clauses (x_i) or (\bar{x}_i) to set those vars to the actual bits of x . E.g.

$$x = 011 : (\bar{x}_1) \wedge (x_2) \wedge (x_3)$$

Thus ϕ_x is computed by an $\tilde{O}(p(n)^2)$ time function of X , and for all $x \in \Sigma^n$:

$$x \in A \Leftrightarrow (\exists y : |y| \leq p(n)) R(x, y) \Leftrightarrow \exists(y, \vec{w}) \phi_x(x, y, w) = 1$$

wires u, v, w .

$$\Leftrightarrow (\exists y) C_n(x, y) = 1 \Leftrightarrow \phi_x \in \text{SAT. } \square$$

Observe: Every clause has at most 3 literals, all of the same sign, and ϕ_x is a conjunction of those clauses. 3SAT.