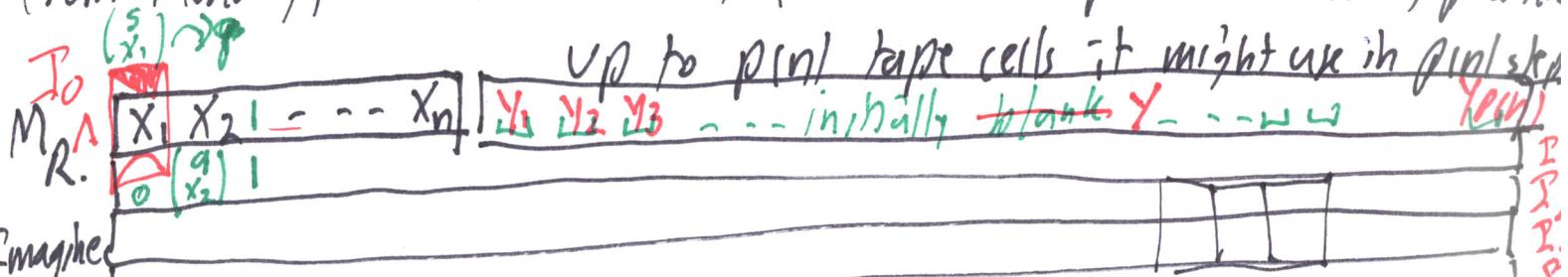


Theorem: [Steve Cook 1970-71; Leonid Levin 1970-1973]
 SAT is NP-complete, i.e. $\text{SAT} \in \text{NP}$ (already seen) and for all $A \in \text{NP}$, $A \leq_m^P \text{SAT}$. [We will prove a reduction to a subcase of SAT]

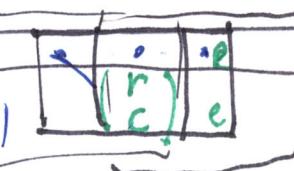
Proof: By $A \in \text{NP}$, there is a polynomial $p(n)$ and a single-tape TM $M_R(x,y)$ that computes a witness predicate $R(x,y)$ with $|y| \leq p(|x|)$.

From Monday, we saw that $R(x,y)$ can be the computation checking predicate.



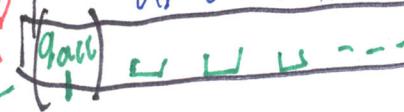
Imagine
after TMs
and out
now to

- The possible effects in any 6-cell
- "Lozenge" can be written as a function depending only on S of M_R .
- It determines the middle bottom cell as a function of the three above.



At the previous step,
the head could have
been in any of the
top 3 cells.

$\vdash p(n)$



$t \leq p(n)$

present
FMA
halfs

before
 $p(n)$ steps.
 $x \in A \equiv$
the win
is 1.

Each lozenge has the same function Δ of six characters (3 inputs, 1 output) over $\{0, 1, X\}$.

We can re-code those over binary and write the finite Δ function as a Boolean function, using NAND gates only. $\therefore M_R$ can be simulated by $p(n) \times p(n)$ sized circuits. (Can get $O(p(n))$ log $p(n)$.)

We have abstracted this to a circuit C_n s.t. $\forall x \in \{0,1\}^n \exists y \text{ M.R. all } (x_i, y) \Leftrightarrow \exists y C_n(x_i, y) = 1$

$$x_1 - x_i = \bar{x}_i \quad y_1 - \dots - y_j = \bar{y}_{p(n)}$$

We can stamp out "

(u and/or v
can be wires
from input x_i
or guesses y_j)



C_n
in $O(p(n)^3)$
time.

Gate g functions correctly
on bits u, v giving output
if and only if ϕ_g is made
true, where

$$\begin{aligned} \phi_g = & (u \vee w) \wedge (\bar{v} \vee w) \\ & \wedge (\bar{u} \vee \bar{v} \vee \bar{w}). \end{aligned}$$

$\vdash (w_0 \quad \text{Final } \phi_x \text{ is:}$

$$(\underset{\substack{\text{NAND gates } g \\ \text{in } C}}{\wedge} \phi_g) \wedge (w_0) \wedge$$

Only part that depends
not just $n = |x|$

Singleton clauses
 (x_i) or (\bar{x}_i) to set
those vars to the
actual bit of x . E.g

$$x = 011 : (\bar{x}_1) \wedge (x_2) \wedge (x_3)$$

wires, u, v, w.

Thus ϕ_x is computed by an $\tilde{O}(p(n)^2)$
time function of X , and for all $x \in \mathbb{Z}^n$:

$$\begin{aligned} X \in A &\Leftrightarrow (\exists y: |y| \leq p(n)) R(x, y) \\ &\Leftrightarrow (\exists y) C_n(x, y) = 1 \Leftrightarrow \exists (y, \bar{w}) \phi_x(x, y, \bar{w}) = 1 \\ &\Leftrightarrow \phi_x \in \text{SAT.} \quad \square \end{aligned}$$

Observe: Every clause has at most 3 literals,
all of the same sign,
and ϕ_x is a conjunction of those clauses. 3SAT.