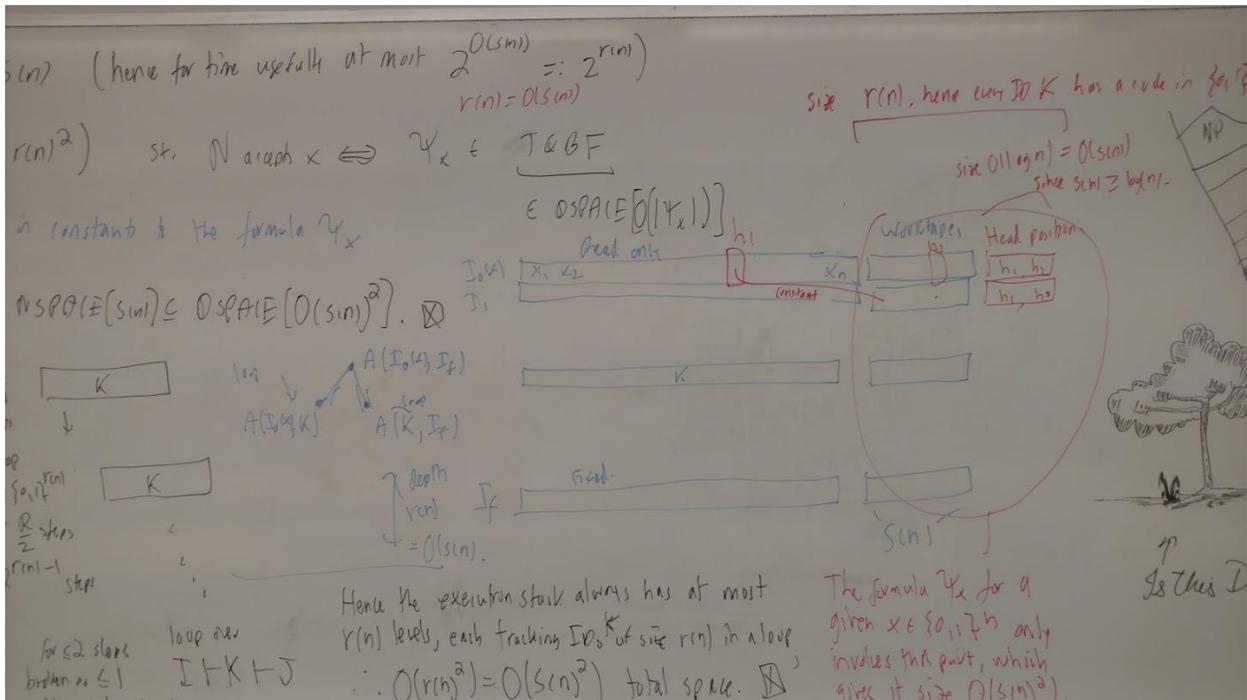


The missing word at upper left is "Recap:"



And for the far-right part of the whiteboard:

a) \equiv BFS Keeping colors of each node.

b) \equiv G3C in NP because

$$G \in G3C \Leftrightarrow (\exists f: V(G) \rightarrow \{0,1,2\}) \bigwedge_{(u,v) \in E} f(u) \neq f(v)$$

encoded by a string of length n over $\{0,1,2\}$.

So this is an NP dec.

in order

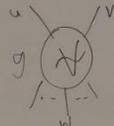
twice

Turing machine)

ignoring $(\log n)$ factor.

at $poly(m,n)$.

3. Suppose we use NOR instead of NAND in the Cook-Levin proof.



$$\phi'_g = (\bar{u} \vee \bar{w}) \wedge (\bar{v} \vee \bar{w})$$

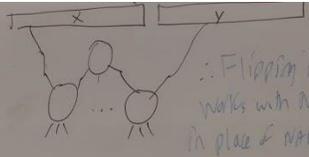
$$\wedge (u \vee v \vee w)$$

\therefore An asst a satisfies $\phi'_g \Leftrightarrow$ its negation a satisfies ϕ_g

$$u=1 \rightarrow w=0 \quad v=0 \wedge v=0 \Rightarrow w=1 \text{ for NAND.}$$

$$\text{SO } u=0 \vee v=0 \rightarrow (u=0 \wedge v=0) \vee w=1$$

$$\therefore u=1 \vee v=1 \vee w=1$$



\therefore Flipping All works with NOR in place of NAND.

Two

Q3 was fully asking: Can we map X to ϕ'_x

So that an assignment a satisfies $\phi'_x \Leftrightarrow$

its flip a' satisfies the original ϕ_x except that

a' does not flip the value of w_0 (and a' still has the same value of x_1, \dots, x_n)?

Key

1 Code the verifier $V(x,y)$ to reject when $R(x,y)$ holds rather than accept. Then the original construction will have $w_0=0$, so flipping all gives $w_0=1$.

2 Code $V(x,y)$ to read y as \bar{y} , i.e. to flip bits of y .

But we can't declare X to be x' instead because X is given.

3 If $x_i=1$, instead of using $\{x_i\}$ as a singleton clause, use $\{\bar{x}_i\}$.

